



Return of CFA: Call-Site Sensitivity Can Be Superior to Object Sensitivity Even for Object-Oriented Programs

Minseok Jeon and Hakjoo Oh



SW재난연구센터 workshop @ Jeju, Korea



Two major camps

Call-Site Sensitivity

Object Sensitivity

Can

Even for

Object-Oriented Programs

Minseok Jeon and Hakjoo Oh



KOREA
UNIVERSITY

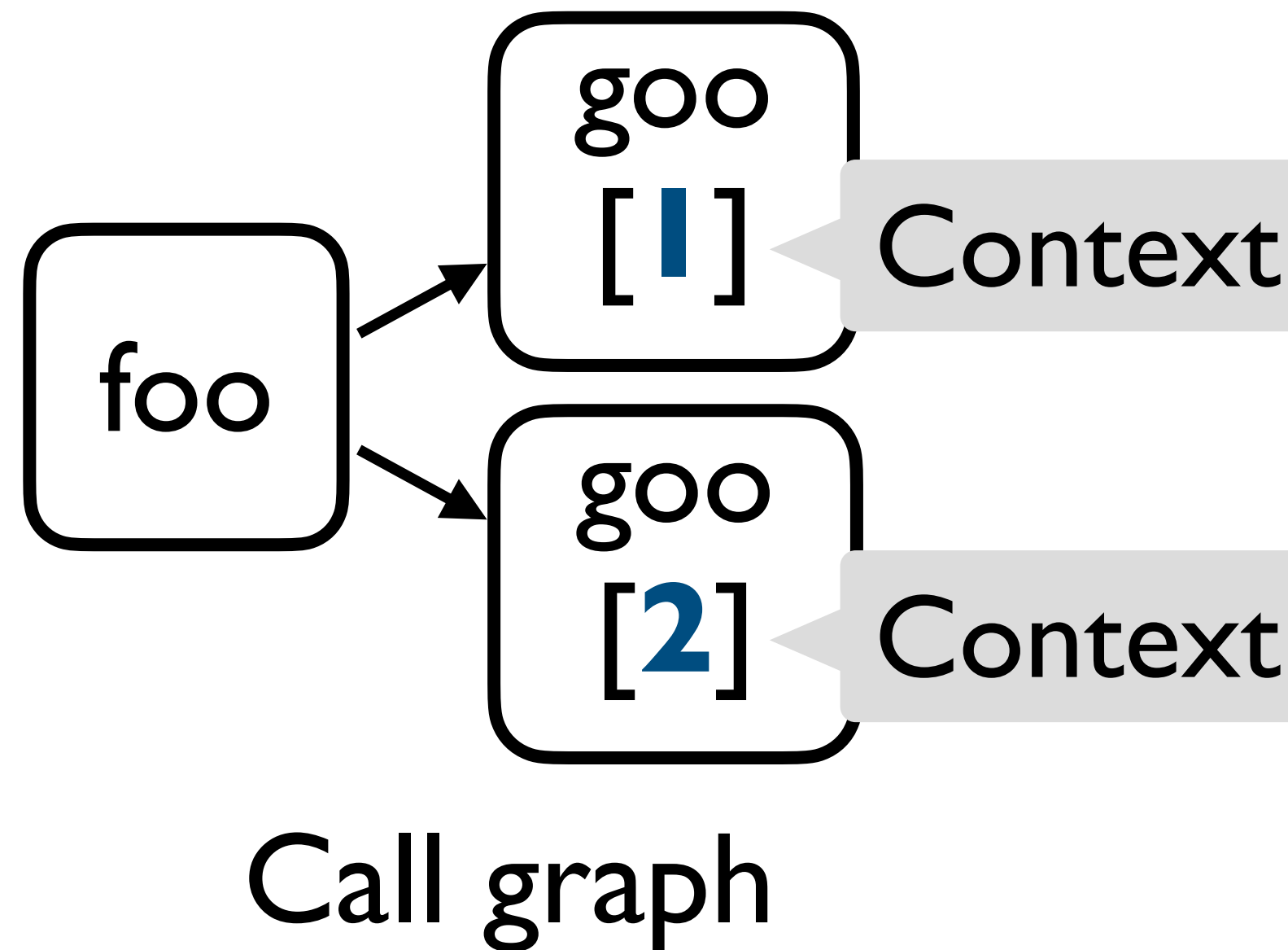
SW재난연구센터 workshop @ Jeju, Korea

Call-site Sensitivity vs Object Sensitivity

Call-site sensitivity was born in 1981

- Considers “**Where**”

```
0: foo(){  
1:   goo();  
2:   goo();  
3: }
```



Call-site sensitivity

1981

2002

2010

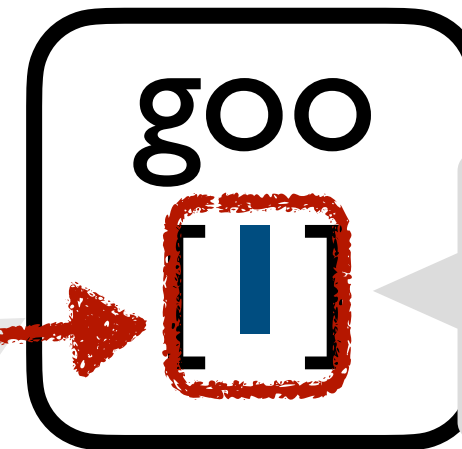
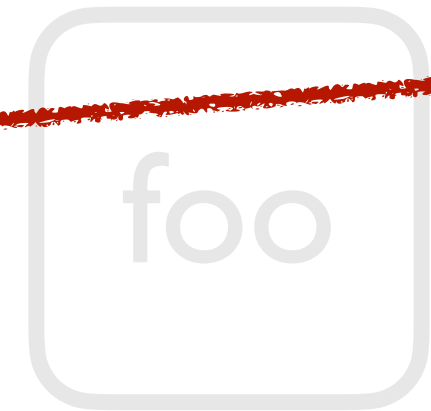
2022

Call-site Sensitivity vs Object Sensitivity

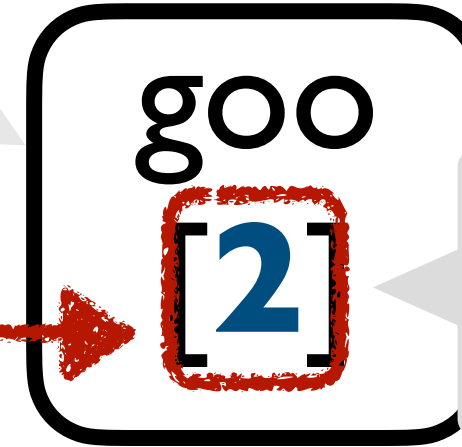
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```
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```



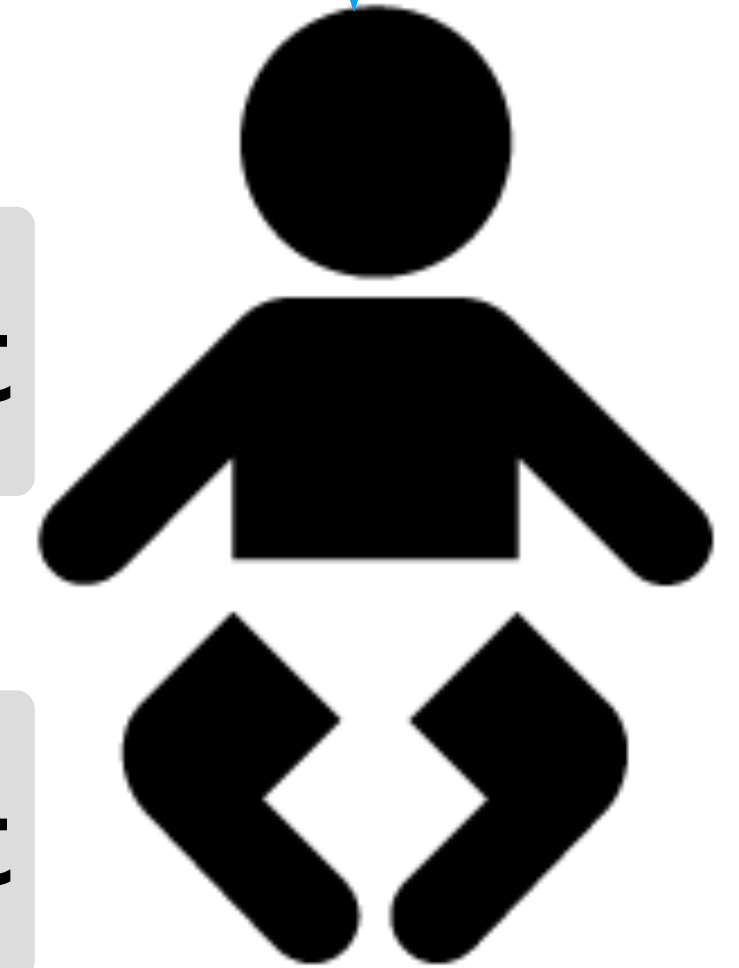
Call-site is context



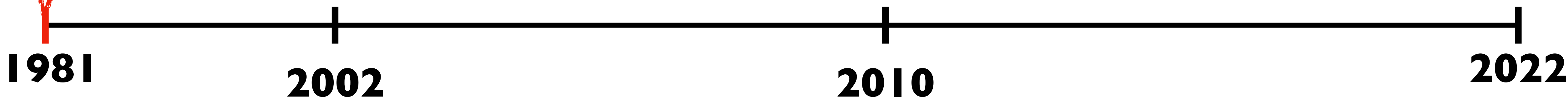
Call-site is context

Call graph

Where is it called from?



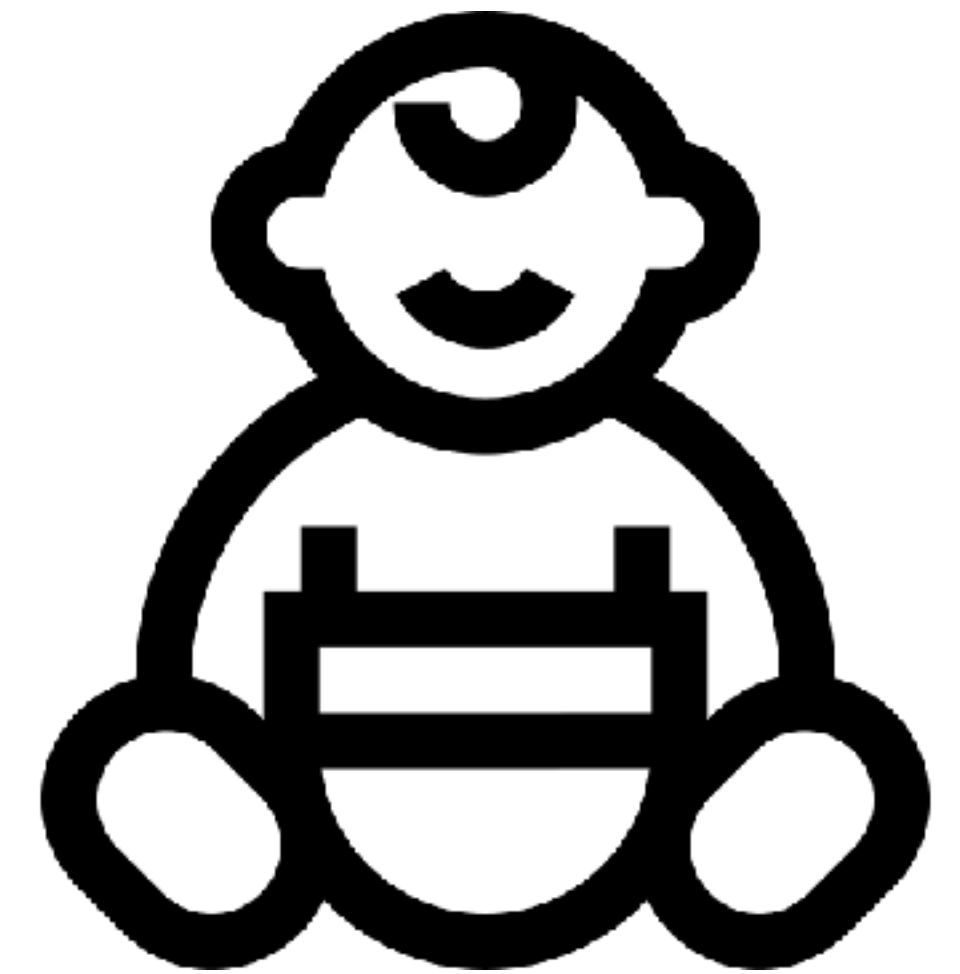
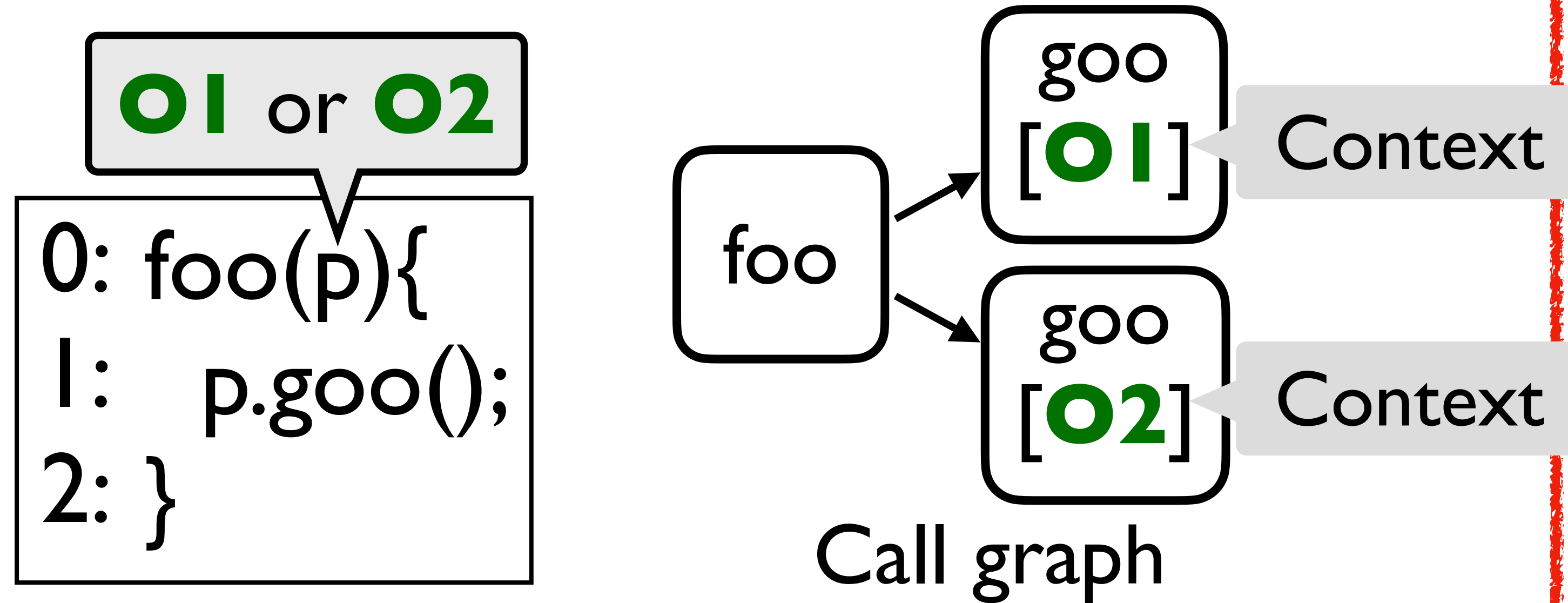
Call-site sensitivity



Call-site Sensitivity vs Object Sensitivity

Object sensitivity appeared in 2002

- Considers “**What**”



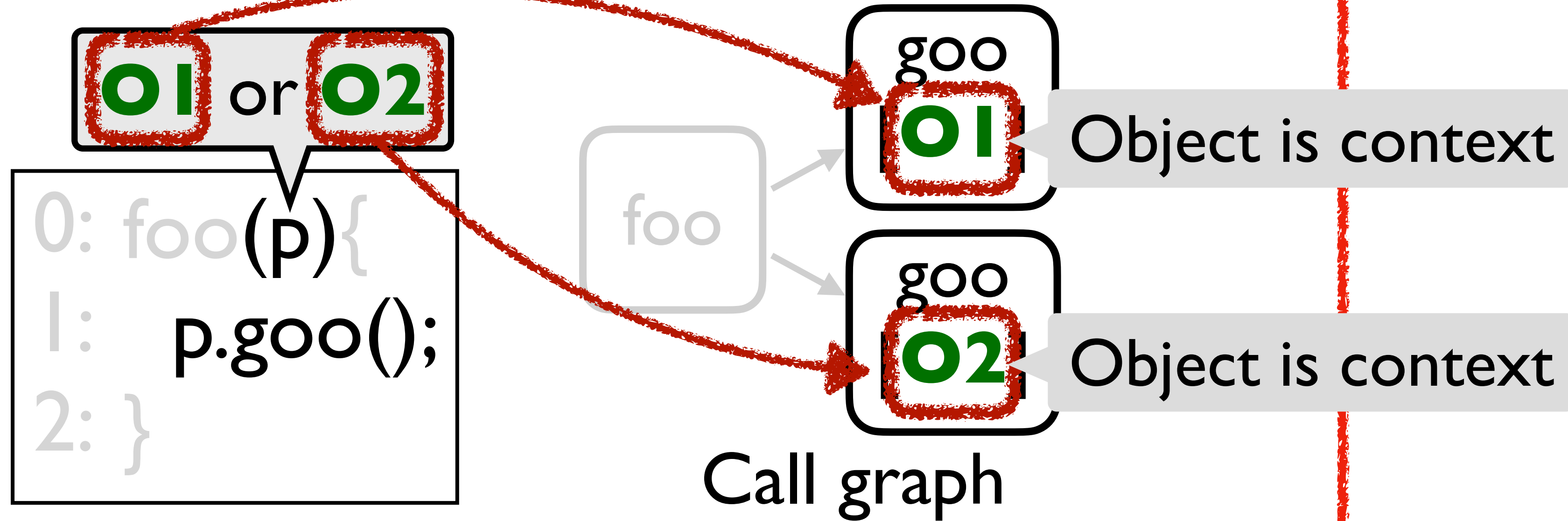
Object sensitivity



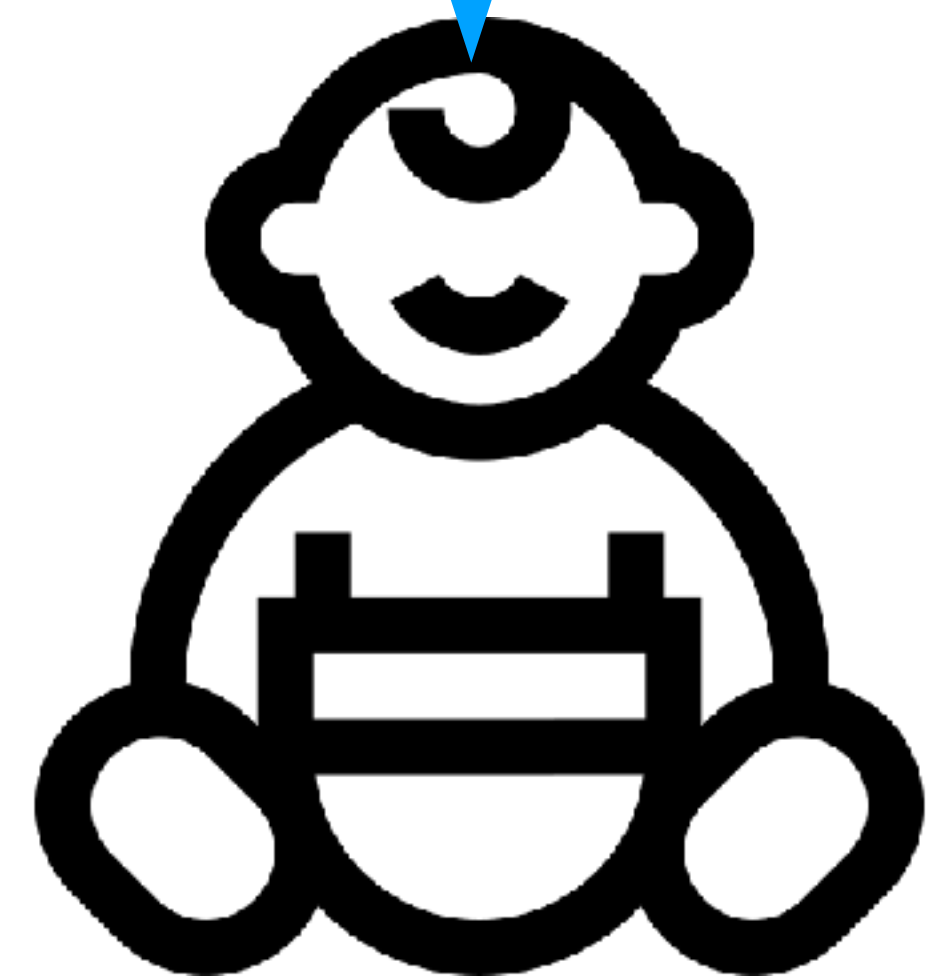
Call-site Sensitivity vs Object Sensitivity

Object sensitivity appeared in 2002

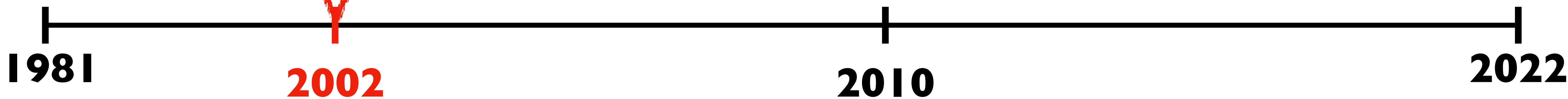
- Considers “**What**”



What is it called with?



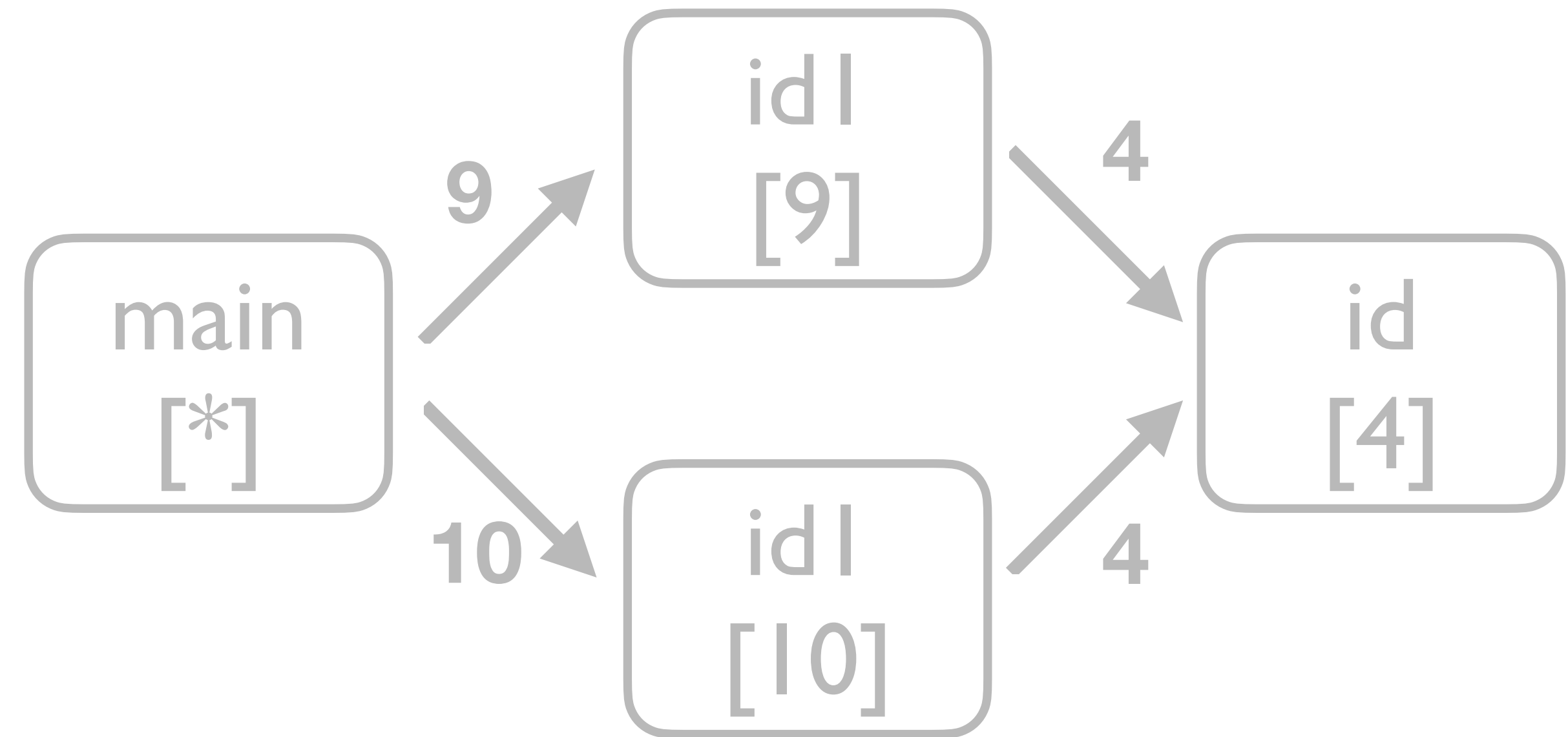
Object sensitivity



Call-site Sensitivity vs Object Sensitivity

- An example shows the **limitation** of CFA and **strength** of object sensitivity

```
0: class C{
1:   id(v){
2:     return v;}
3:   idl(v){
4:     return this.id(v);}
5: }
6: main(){
7:   c1 = new C();//C1
8:   c2 = new C();//C2
9:   a = (A) c1.idl(new A());//query1
10:  b = (B) c2.idl(new B());//query2
11: }
```



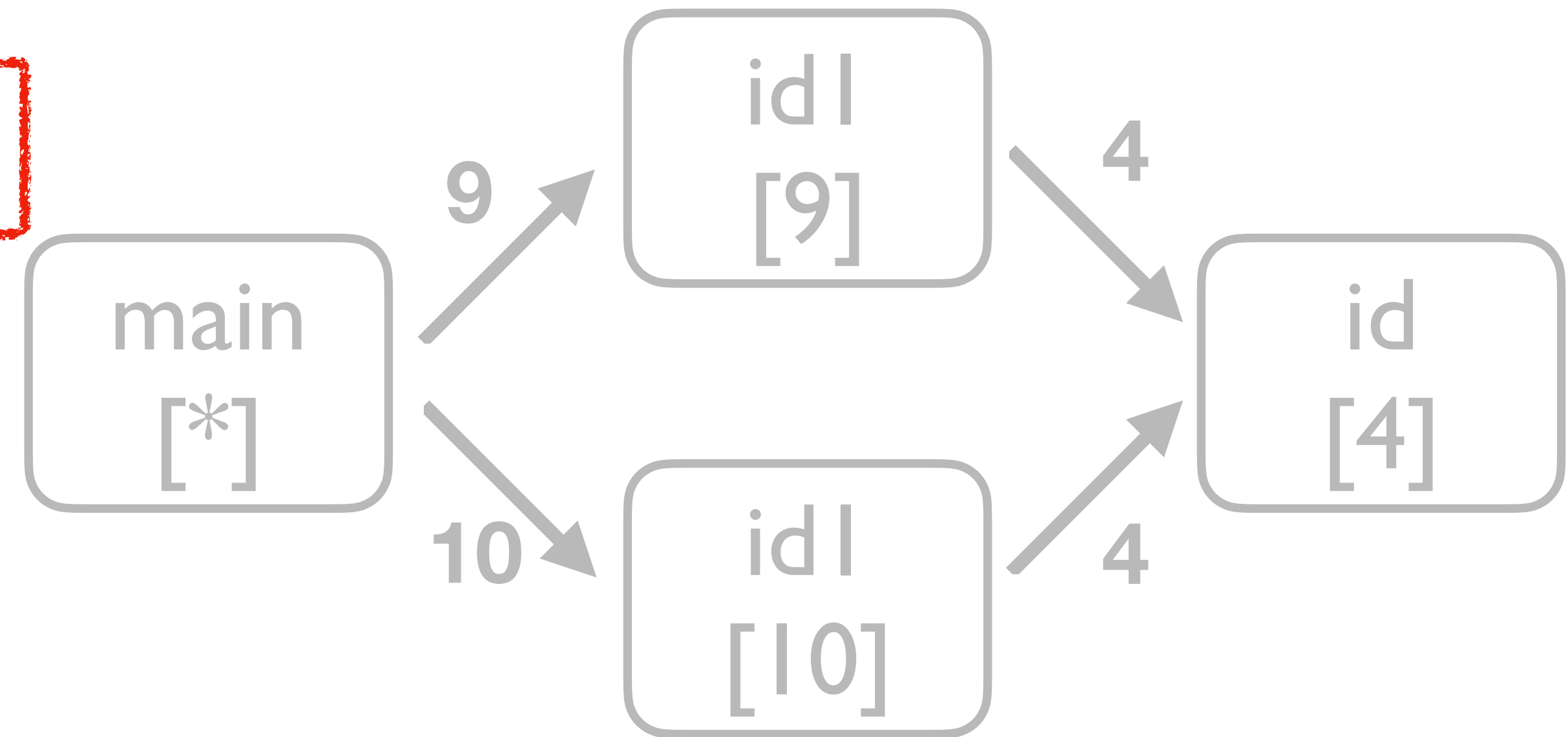
Call-graph of I-CFA

Call-site Sensitivity vs Object Sensitivity

- An example shows the **limitation** of CFA and **strength** of object sensitivity

```
0: class C{
1:   id(v){
2:     return v;}
3:   idI(v){
4:     return this.id(v);}
5: }
6: main(){
7:   c1 = new C();//C1
8:   c2 = new C();//C2
9:   a = (A) c1.idI(new A());//query1
10:  b = (B) c2.idI(new B());//query2
11: }
```

Identity function



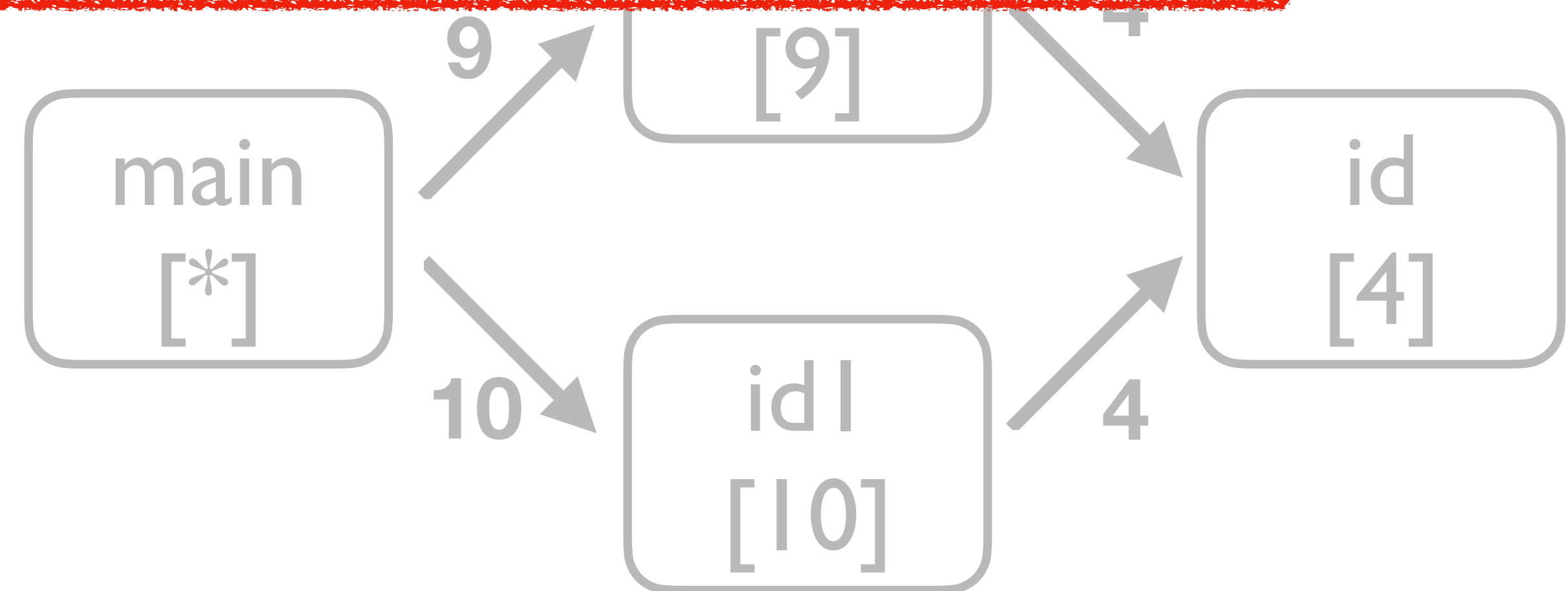
Call-graph of I-CFA

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5: }
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8:   c2 = new C();//C2
9:   a = (A) c1.idl(new A());//query1
10:  b = (B) c2.idl(new B());//query2
11: }
```

Also an identity function implemented with id

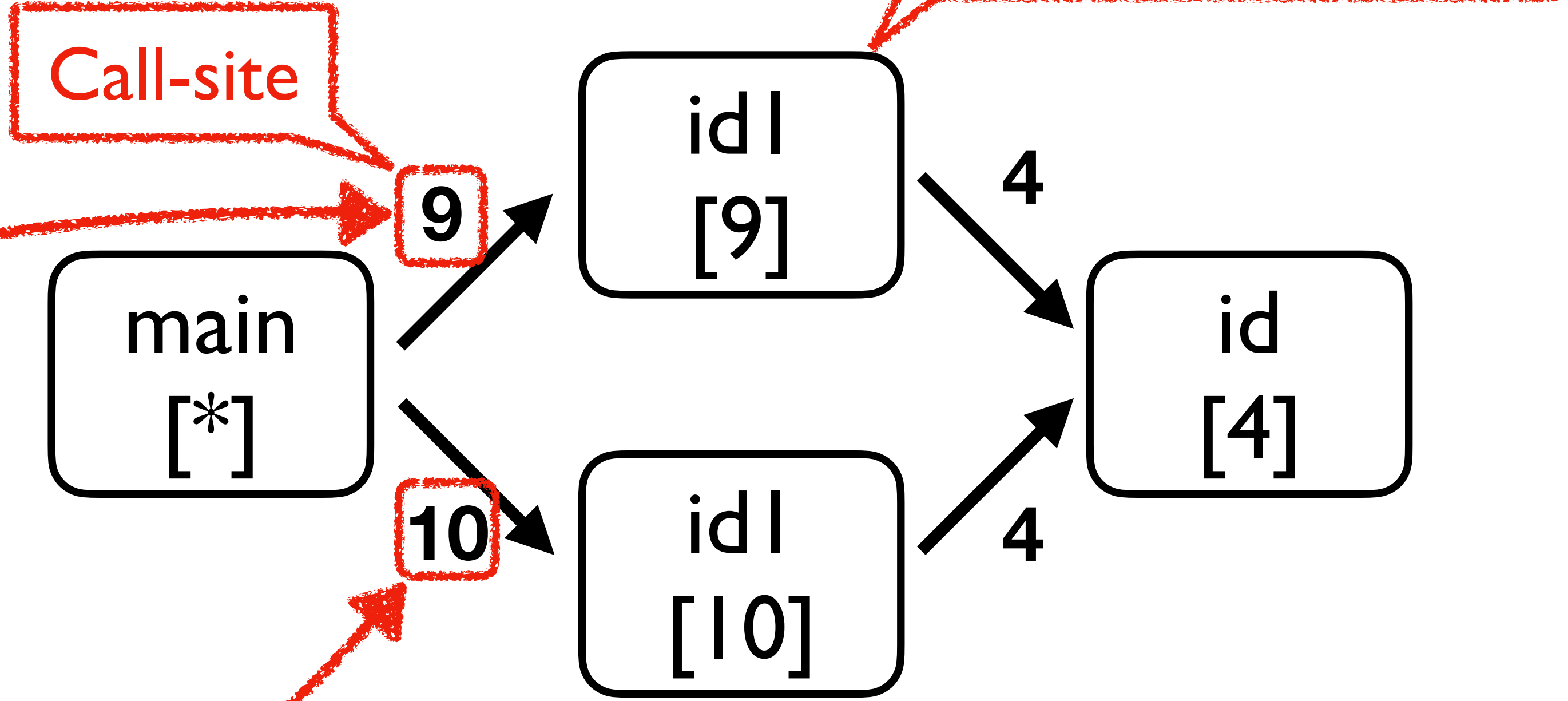


Call-graph of I-CFA

Call-site Sensitivity vs Object Sensitivity

- An example shows the **limitation** of CFA and **strength** of **Method & Context**

```
0: class C{
1:   id(v){
2:     return v;}
3:   idl(v){
4:     return this.id(v);}
5: }
6: main(){
7:   c1 = new C();//C1
8:   c2 = new C();//C2
9:   a = (A) c1.idl(new A());//query1
10:  b = (B) c2.idl(new B());//query2
11: }
```



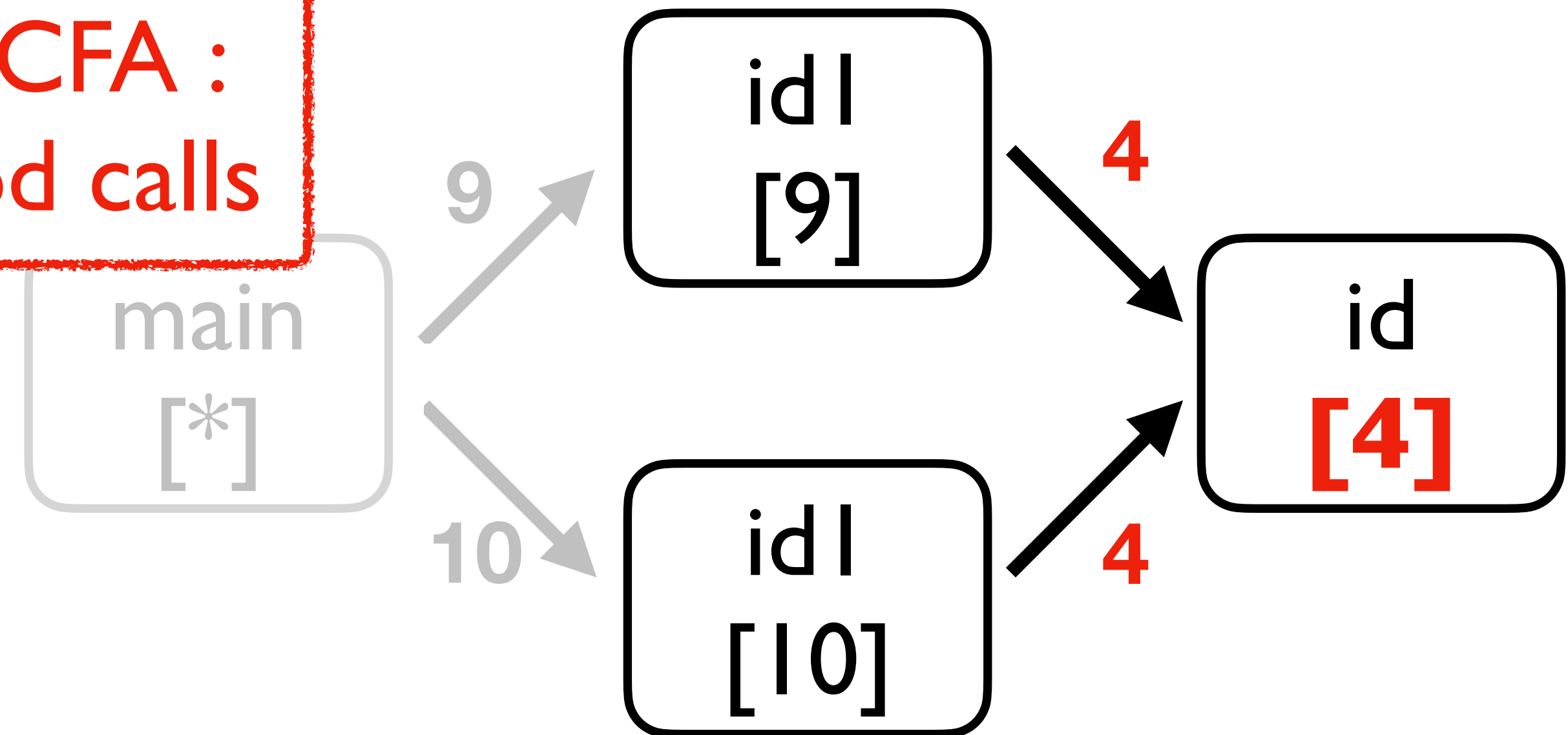
Call-graph of I-CFA

Call-site Sensitivity vs Object Sensitivity

- An example shows the **limitation** of CFA and strength of object sensitivity

```
0: class C{
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3:   idI(v){
4:     return this.id(v);}
5: }
6: main(){
7:   c1 = new C();//C1
8:   c2 = new C();//C2
9:   a = (A) c1.idI(new A());//query1
10:  b = (B) c2.idI(new B());//query2
11: }
```

Limitation of CFA :
Nested method calls

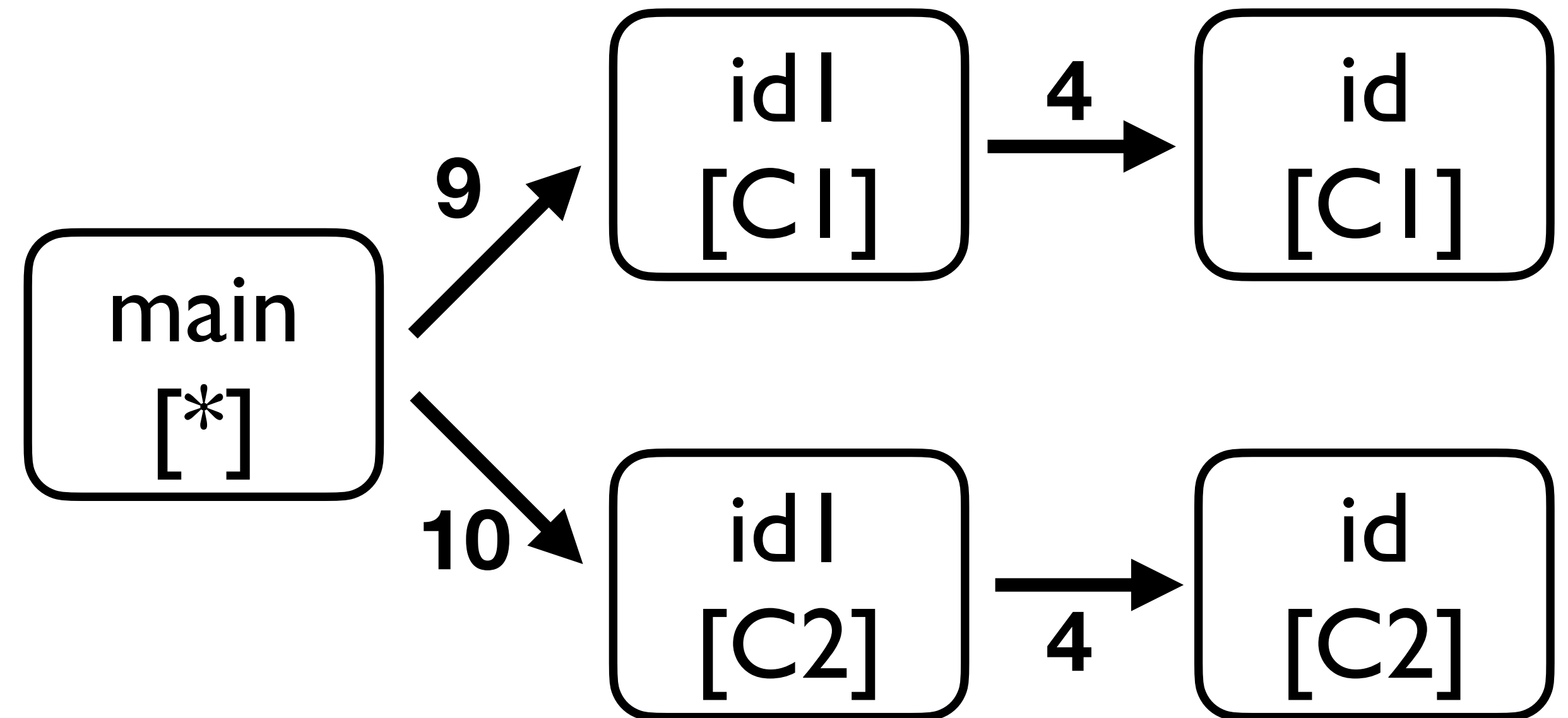


Call-graph of I-CFA

Call-site Sensitivity vs Object Sensitivity

- An example shows the **limitation** of CFA and **strength** of object sensitivity

```
0: class C{
1:   id(v){
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3:   idl(v){
4:     return this.id(v);}
5: }
6: main(){
7:   c1 = new C();//C1
8:   c2 = new C();//C2
9:   a = (A) c1.idl(new A());//query1
10:  b = (B) c2.idl(new B());//query2
11: }
```



Call-graph of l-Obj

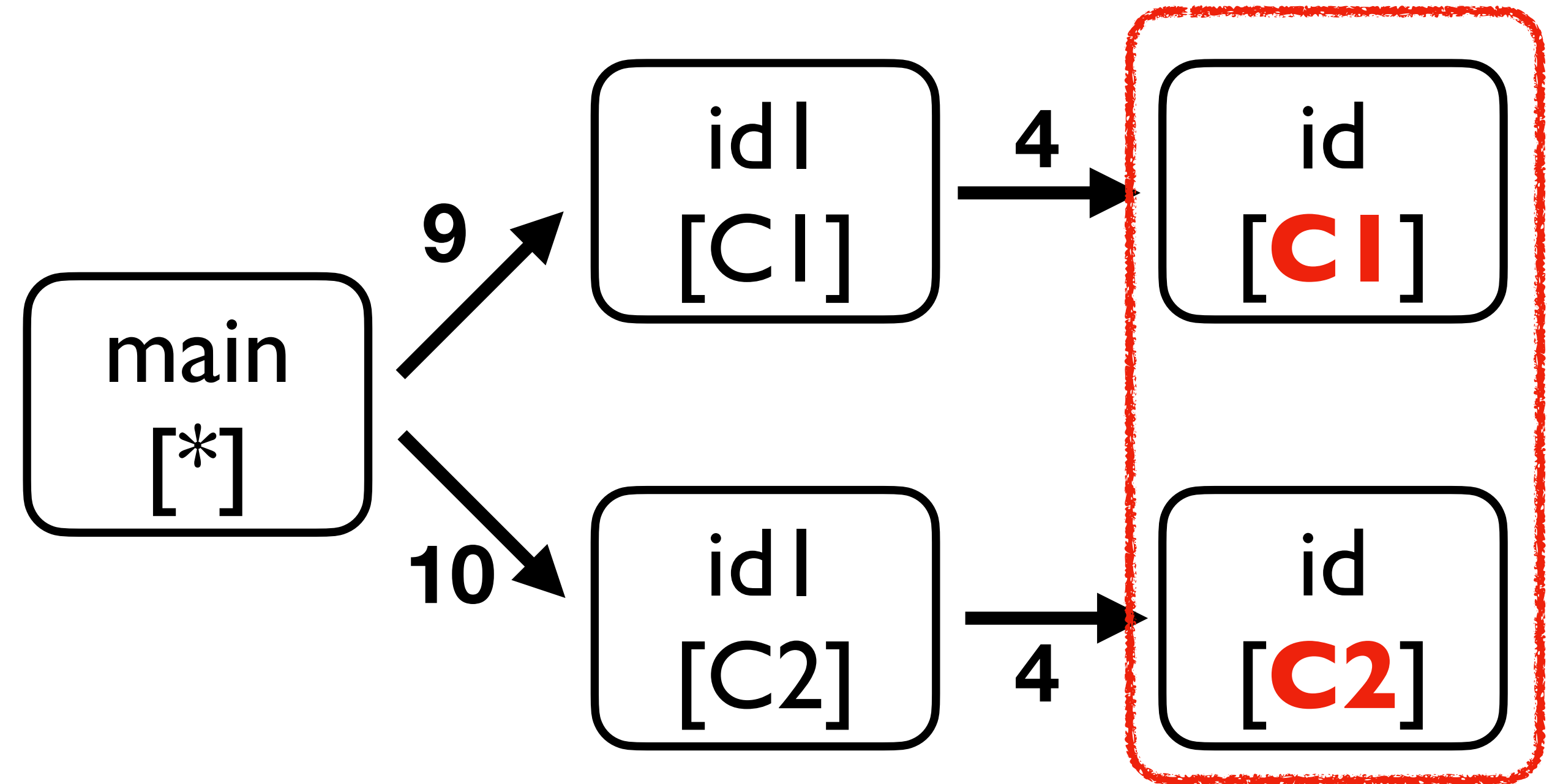
Call-site Sensitivity vs Object Sensitivity

- An example shows the limitation of CFA and **strength of object sensitivity**

```
0: class C{
1:   id(v){
2:     return v;}
3:   idl(v){
4:     return this.id(v);}
5: }
6: main(){
```

C1 or C2

```
7: c1 = new C();//C1
8: c2 = new C();//C2
9: a = (A) c1.idl(new A());//query1
10: b = (B) c2.idl(new B());//query2
11: }
```

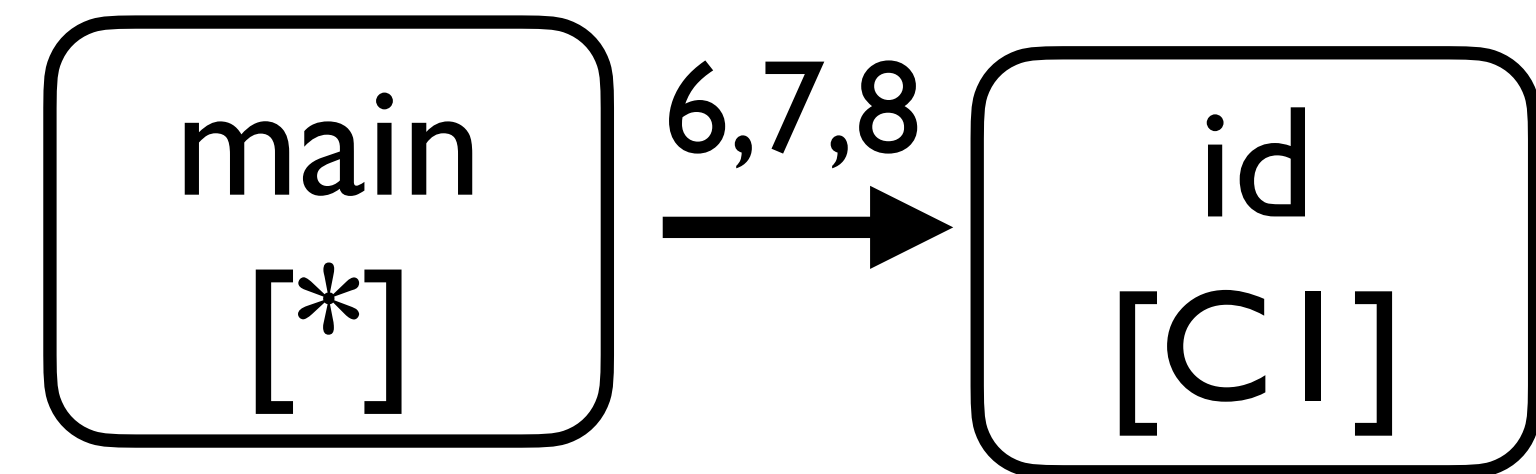


Call-graph of I-Obj

Call-site Sensitivity vs Object Sensitivity

- An example shows the **limitation** of **object sensitivity** and **strength** of **CFA**

```
0: class C{
1:   id(v){
2:     return v;}
3: }
4: main(){
5:   c1 = new C();//CI
6:   a = (A) c1.id(new A());//query1
7:   b = (B) c1.id(new B());//query2
8:   c = (B) c1.id(new C());//query3
9: }
```

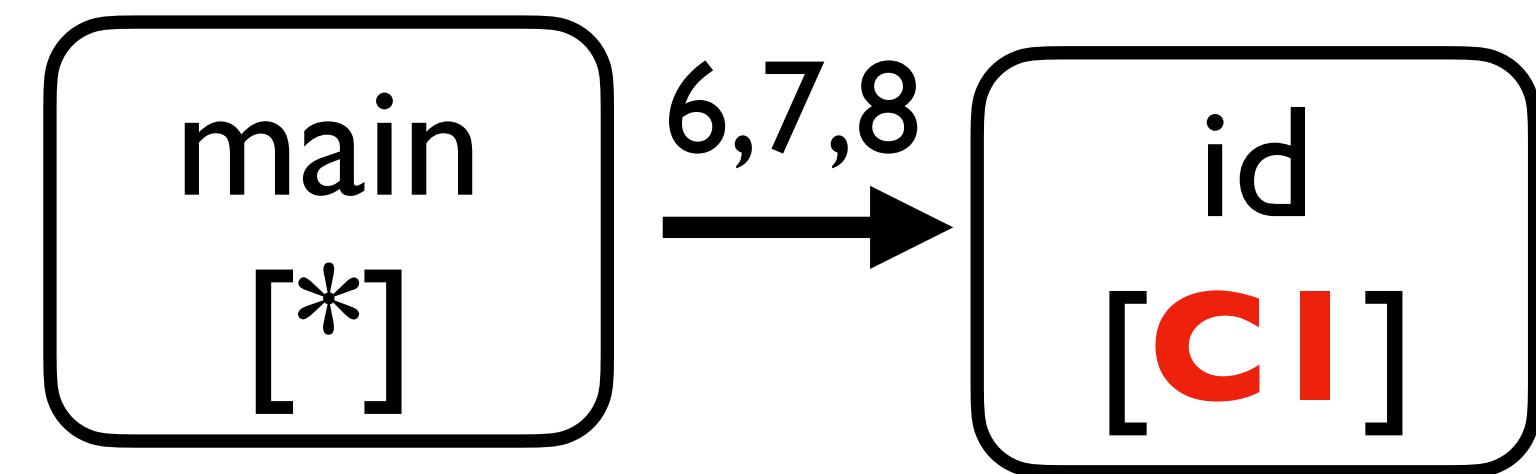


Call-graph of I-Obj

Call-site Sensitivity vs Object Sensitivity

- An example shows the **limitation** of **object sensitivity** and strength of CFA

```
0: class C{
1:   id(v){
2:     return v;}
3: }
4: main(){
5:   cI = new C();//CI
6:   a = (A) cI.id(new A());//query1
7:   b = (B) cI.id(new B());//query2
8:   c = (B) cI.id(new C());//query3
9: }
```



Call-graph of I-Obj

The three method calls share the same receiver object CI

Call-site Sensitivity vs Object Sensitivity

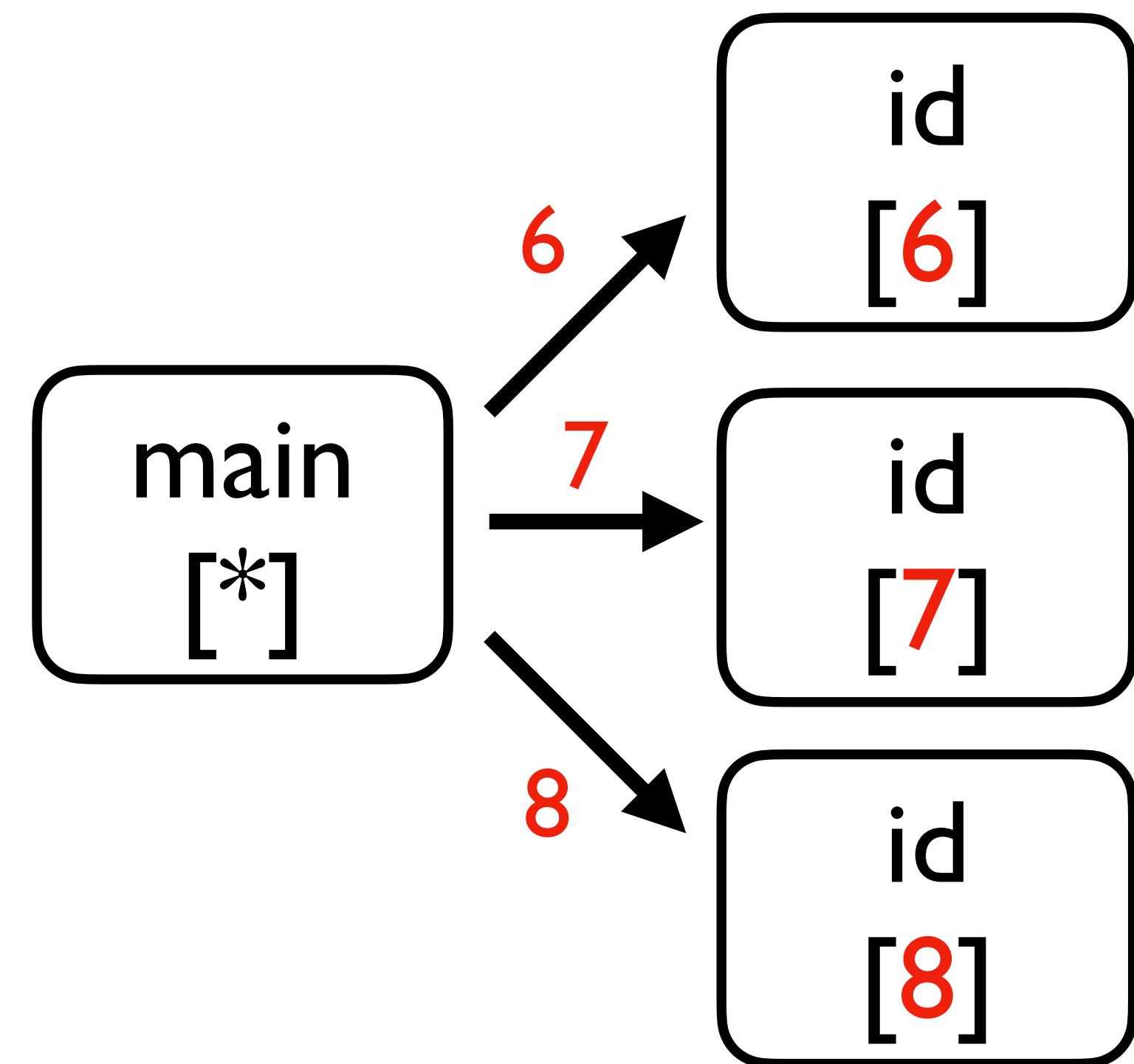
- An example shows the **limitation** of object sensitivity and **strength** of CFA

```
0: class C{  
1:   id(v){  
2:     return v;}  
3: }
```

```
4: main(){  
5:   c1 = new C();//C1
```

```
6:   a = (A) c1.id(new A());//query1  
7:   b = (B) c1.id(new B());//query2  
8:   c = (C) c1.id(new C());//query3
```

```
9: }
```



Call-graph of I-CFA

Call-site sensitivity easily separates the three method calls

Call-site Sensitivity vs Object Sensitivity

- Call-site Sensitivity and Object Sensitivity had been **actively compared**

Call-site Sensitivity vs Object Sensitivity



Parameterized Object Sensitivity for Points-to Analysis for Java

ANA MILANOVA
Researcher Polytechnic Institute
ATANAS BOUNTEV
Catic State University
and
BARBARA G. FISHER
Rutgers University

The goal of points-to analysis for Java is to determine the set of objects pointed to by a reference variable at a reference site, via field. We present a novel sensitivity analysis for Java that is based on the idea of sensitivity analysis for Java. The key idea of this approach is to analyze sensitivity at the level of the type system, not the concrete values of variables. This allows us to analyze programs that use dynamic dispatch, which is not possible in traditional points-to analysis.

Our other contributions are the sensitivity analysis that may be modified by the selection of program parameters. This analysis is based on the idea of sensitivity analysis for Java, and is implemented as a tool. The tool is implemented as a set of Java classes and is implemented as a set of Java classes and is implemented as a set of Java classes.

We have implemented and evaluated our parameterized object sensitivity points-to analysis. Our results show that our analysis is more precise than the state-of-the-art analysis for Java. Our analysis is more precise than the state-of-the-art analysis for Java. Our analysis is more precise than the state-of-the-art analysis for Java.

This work was supported by the National Science Foundation (NSF) grant CCF-0500000. Authors' addresses: A. Milanova, Department of Computer Science, Simon Fraser University, 8888 University Drive, Burnaby, BC V5A 1S8, Canada; B. G. Fisher, Department of Computer Science and Engineering, Ohio State University, 2035 Neil Avenue, Columbus, OH 43210, email: barbara.fisher@osu.edu; B. G. Fisher, Department of Computer Science, Rutgers University, 117 Frelinghuysen Road, Piscataway, NJ 08854, email: barbara.fisher@rutgers.edu.

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ACM Transactions on Software Engineering and Methodology, Vol. 11, No. 1, January 2002, Pages 1-18.

Context-sensitive points-to analysis: Is it worth it?*

Chang Liu¹ and Laurie Hendren²
1 School of Computer Science, University of Waterloo, Waterloo, ON, Canada
2 School of Computer Science, McGill University, Montreal, QC, Canada

Abstract. We present the results of an empirical study evaluating the precision of context-sensitive points-to analysis with several variations of context sensitivity, or how much context information is used. We compare the use of call-site strings in the context-sensitive points-to analysis, and the use of call-site strings in the context-sensitive points-to analysis, and the use of call-site strings in the context-sensitive points-to analysis.

1 Introduction

Context-sensitive points-to analysis is a form of interprocedural analysis of object-oriented programs. It is often argued that it is worth it, but lack of available implementations has hindered the widespread adoption of this technique.

Of the many context-sensitive points-to analyses that have been proposed [2, 4, 8, 11, 17, 18, 28, 31], which represent a wide variety of techniques, we are likely to find context-sensitive points-to analysis more precise than the state-of-the-art implementation of a specific analysis. We should have empirical evidence to back up these claims.

This study aims to provide these answers. Recent advances in the use of Binary Decision Diagrams (BDDs) in program analysis [3, 12, 29, 31] have made context-sensitive points-to analysis efficient enough to perform empirical studies. Our analysis is based on the idea of context-sensitive points-to analysis, and we have implemented three different variations of context-sensitive points-to analysis, and we have measured their precision in terms of several client analyses. Specifically, we compared the use of call-site strings in the context-sensitive points-to analysis [17, 28], and the algorithm proposed by Zhu and Calman [31].

*This work was supported, in part, by NSERC and an IBM FELLOW.

Evaluating the Benefits of Context-Sensitive Points-to Analysis Using a BDD-Based Implementation

ONDREJ LHOTAK
University of Waterloo
and
LAURIE HENDREN
McGill University

We present Points, a framework of BDD-based context-sensitive points-to analysis for Java, as well as other analyses that use these results. Points is a novel, context-sensitive points-to analysis for Java, as well as other analyses that use these results. Points is a novel, context-sensitive points-to analysis for Java, as well as other analyses that use these results.

ACM Reference Format:
Lhotak, O. and Hendren, L. 2002. Evaluating the benefits of context-sensitive points-to analysis using a BDD-based implementation. *ACM Trans. Softw. Eng. Methodol.* 11, 1. Article 11 (December 2002), 20 pages. DOI = 10.1145/581951.1001951 <http://doi.acm.org/10.1145/581951.1001951>

This is a revised and extended version of an article which appeared in Proceedings of the 11th International Conference on Compiler Construction, Lecture Notes in Computer Science, vol. 2623, Springer, 44.

Authors' addresses: O. Lhotak, D. E. Condon, School of Computer Science, University of Waterloo, 200 University Avenue West, Waterloo, ON, N2L 2G1, Canada; L. Hendren, School of Computer Science, McGill University, 3480 University Street, Room 304, Montreal, QC, H3A 2K7, Canada.

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ACM Transactions on Software Engineering and Methodology, Vol. 11, No. 1, December 2002, Pages 9803.

Strictly Declarative Specification of Sophisticated Points-to Analyses

Markus Bruegger, Yannis Smaragdakis
Department of Computer Science
University of Massachusetts, Amherst
Amherst, MA 01003, USA
markus.bruegger@comlab.org yannis@cs.umass.edu

Abstract. We present the first framework for writing analysis of Java programs. This is done by using a declarative specification language for defining point-to analysis. We use the declarative approach to define point-to analysis for Java, as well as other analyses that use these results. Points is a novel, context-sensitive points-to analysis for Java, as well as other analyses that use these results.

1. Introduction

Points is a declarative specification language for writing analysis of Java programs. This is done by using a declarative specification language for defining point-to analysis. We use the declarative approach to define point-to analysis for Java, as well as other analyses that use these results.

1981

2002

2010

2022

Call-site Sensitivity vs Object Sensitivity

- Object Sensitivity outperformed call-site sensitivity

Call-site Sensitivity vs Object Sensitivity

Parameterized Object Sensitivity for Points-to Analysis for Java
ANA MILANOVA, PANOS KONTOPOULOS, ANTONIOS ROUNTIS, BARBARA G. RYDER
Abstract: We present the results of an empirical study evaluating the precision of abstract-based points-to analysis with several variations of context sensitivity, or lack thereof, in the context of Java. We compare the use of call-site strings in the context of abstract object sensitivity, and the 30 abstract context-sensitive strategies proposed by Zhu and Collins, and by Wang and Lam. Our study includes analysis that automatically synthesizes only pointer variables, as well as some that specify the heap abstraction. We measure both the precision of the points-to sets themselves, as well as their effect on the precision of client analyses. To guide development of efficient analysis implementations, we measure the number of nodes, the number of abstract constants, and the number of abstract pointers set that arise with each context-sensitive variation. To evaluate precision, we compare the results of the call-site graphs to those of nodes and edges, the number of abstract constants, and the number of abstract pointers set.

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Chang Liao† and Laurie Hendren‡
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Evaluating the Benefits of Context-Sensitive Points-to Analysis Using a BDD-Based Implementation
ONDREJ LHOTAK
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and
LAURIE HENDREN
McGill University
Abstract: We present Points, a framework of BDD-based context-sensitive points-to analysis for Java, as well as other analyses that use these models. Points supports several variations of context-sensitive points-to analysis, including call-site sensitivity and context-sensitive pointer sets for both pointers and non-pointers. We empirically evaluate the precision of these context-sensitive analyses on significant Java programs. We find that context-sensitive analyses are more precise than comparable variations of the other approaches, and that specializing the heap abstraction improves precision more than extending the length of context strings.

Strictly Declarative Specification of Sophisticated Points-to Analyses
Markus Bruegger, Yannis Smaragdakis
Department of Computer Science
University of Massachusetts, Amherst
Abstract: We present the Doo framework for automatic analysis of Java programs. Doo is based on the idea of specifying points-to analysis algorithms declaratively, using Doo, a logic-based language for defining abstract relations. We carry the declarative approach further than past work by allowing the full expression of new inlining and unwinding abstractions using a novel technique specifically targeting high-resolution Doo programs. As a result, Doo achieves several benefits, including full control over the implementation of points-to analysis, as well as the ability to specify the state of the heap in context-sensitive analyses. For the most sophisticated points-to analyses (and, consequently, identical points-to analyses), Doo is faster than other tools. In a declarative analysis of the Doo framework, we show that Doo is more precise than other points-to analyses, and that its declarative approach is more precise than other points-to analyses. Additionally, Doo enables more precise analysis than is possible with other tools, and it is the only context-sensitive analysis that is both declarative and precise. Finally, context-sensitive analysis is made to be easy to integrate with a wide range of characteristics, largely due to its declarative nature.



Call-site Sensitivity vs Object Sensitivity

- Lectures have taught the **superiority** of object sensitivity



Obj

Object-Sensitivity

- The dominant flavor of context-sensitivity for object languages.
- It uses object abstractions (i.e. allocation sites) as qualifying a method's local variables with the allocation receiver object of the method call.

```
class A { void m() { return; } }
...
b = new B();
b.m();
```

The context of `m` is the allocation site of `b`.

Object-Sensitivity (vs. call-site sensitivity)

```
class S {
  Object id(Object a) { return a; }
  Object id2(Object a) { return id(a); }
}
class C extends S {
  void fun1() {
    Object a1 = new A1();
    Object b1 = id2(a1);
  }
}
class D extends S {
  void fun2() {
    Object a2 = new A2();
    Object b2 = id2(a2);
  }
}
```

Object-sensitive pointer

- Milanova, Rountev, and Ryder. *Parameterized sensitivity for points-to analysis for Java*. *ACM Eng. Methodol.*, 2005.
- Context-sensitive interprocedural pointer analysis
- For context, use stack of receiver objects
- (More next week?)
- Lhotak and Hendren. *Context-sensitive pointer worth it?* *CC 06*
- Object-sensitive pointer analysis more precise than for Java
- Likely to scale better

Lecture Notes: Pointer Analysis

15-819C: Program Analysis
Jonathan Aldrich
jonathan.aldrich@cs.cmu.edu
Lecture 9

1 Motivation for Pointer Analysis

In programs with pointers, program analysis can become more difficult. Consider constant-propagation analysis of the following program:

```
1: x = 1;
2: p := &x;
3: *p = 2;
4: print x;
```

In order to analyze this program correctly we must be able to determine what instruction `p` points to. If this information is available we can flow function as follows:

$$f_{CP}[*p := y](e) = [x \mapsto y]e \quad \text{where } \text{must_point}(p);$$

When we know exactly what a variable `x` points to, we say it has *exact-point-to* information, and we can perform a strong update of variable `x`, because we know with confidence that assigning `y` to `x` is a technicality in the rule is quantifying over all `x` such that `x` points to `a`. How is this possible? It is not possible in C or Java, a language with pass-by-reference, for example C++, it is possible for the same location are in scope. Of course, it is also possible that we are uncertain to which distinct locations `p` points. For example:

Pointer Analysis

Yannis Smaragdakis
University of Athens
smaragd@cc.uoi.gr

George Balacouras
University of Athens
gbalacou@cc.uoi.gr

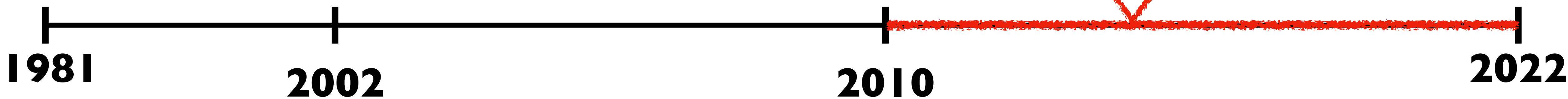
now
the essence of knowledge
Boston — Delhi



National and Kapodistrian University of Athens



Carnegie Mellon University



Call-site Sensitivity vs Object Sensitivity

- Lectures have taught the **superiority** of object sensitivity

Object-Sensitivity

- The dominant flavor of context-sensitivity for object-oriented languages.
- It uses object abstractions (i.e. allocation sites) as contexts, qualifying a method's local variables with the allocation site of the receiver object of the method call.

```
class A { void m() { return; } }  
...  
b = new B();  
b.m();  
The context of m is the allocation site of b.
```

Object-sensitive pointer analysis

- Milanova, Rountev, and Ryder. *Parameterized object sensitivity for points-to analysis for java*. ACM Trans. Softw. Eng. Comput. 2004.

Lecture Notes: Pointer Analysis

15-819C: Program Analysis
Jonathan Aldrich
jonathan.aldrich@cs.cmu.edu

Pointer Analysis

Yannis Smaragdakis
University of Athens

I was also taught like that



Obj



Call-site Sensitivity vs Object Sensitivity

- Researches focused on improving Object Sensitivity

Researches on Object Sensitivity

Pick Your Call
Yannis Smaragdakis
Abstract: Call-site sensitivity is a key property of pointer analysis. It is the ability to distinguish between different call sites of a method. This is important for many applications, such as garbage collection, security, and program optimization. In this paper, we present a new algorithm for call-site sensitivity analysis. Our algorithm is based on a novel abstraction of the call graph. It is able to handle a large number of call sites, and it is more precise than previous algorithms. We evaluate our algorithm on a set of benchmarks, and we show that it is more precise and faster than previous algorithms.

Hybrid Context-Sensitive Pointer Analysis
Tian Tan, Yao Li
Abstract: Object sensitivity is a key property of pointer analysis. It is the ability to distinguish between different call sites of a method. This is important for many applications, such as garbage collection, security, and program optimization. In this paper, we present a new algorithm for call-site sensitivity analysis. Our algorithm is based on a novel abstraction of the call graph. It is able to handle a large number of call sites, and it is more precise than previous algorithms. We evaluate our algorithm on a set of benchmarks, and we show that it is more precise and faster than previous algorithms.

Making k-Object-Sensitive Pointer Analysis More Precise with Still k-Limiting
Tian Tan, Yao Li
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Efficient Modeling the Heaps
Yue Li, Tian Tan
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Precision-Guided Context-Sensitive Pointer Analysis
Yue Li, Tian Tan
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Scalability-First Self-Tuning Context-Sensitive Pointer Analysis
Yue Li, Tian Tan
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Data-Driven Context-Sensitive Pointer Analysis
Sehun Jeong, Minseok Jeon, MyungHo Lee, and Haejoon Choo
Abstract: Object sensitivity is a key property of pointer analysis. It is the ability to distinguish between different call sites of a method. This is important for many applications, such as garbage collection, security, and program optimization. In this paper, we present a new algorithm for call-site sensitivity analysis. Our algorithm is based on a novel abstraction of the call graph. It is able to handle a large number of call sites, and it is more precise than previous algorithms. We evaluate our algorithm on a set of benchmarks, and we show that it is more precise and faster than previous algorithms.

Learning Graph-based Heuristics without Handcrafting Applications
Minseok Jeon, MyungHo Lee, and Haejoon Choo
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Precision-Preserving Analysis with Partial Context Sensitivity
Jingbo Lu, UNSW Sydney, Australia
Abstract: Object sensitivity is a key property of pointer analysis. It is the ability to distinguish between different call sites of a method. This is important for many applications, such as garbage collection, security, and program optimization. In this paper, we present a new algorithm for call-site sensitivity analysis. Our algorithm is based on a novel abstraction of the call graph. It is able to handle a large number of call sites, and it is more precise than previous algorithms. We evaluate our algorithm on a set of benchmarks, and we show that it is more precise and faster than previous algorithms.

Making Pointer Analysis More Precise by Unleashing the Power of Selective Context Sensitivity
Tian Tan, Yao Li, Xiaoting Na, Chaozuo Xu, and Yannis Smaragdakis
Abstract: Object sensitivity is a key property of pointer analysis. It is the ability to distinguish between different call sites of a method. This is important for many applications, such as garbage collection, security, and program optimization. In this paper, we present a new algorithm for call-site sensitivity analysis. Our algorithm is based on a novel abstraction of the call graph. It is able to handle a large number of call sites, and it is more precise than previous algorithms. We evaluate our algorithm on a set of benchmarks, and we show that it is more precise and faster than previous algorithms.



Obj

1981 | 2002 | 2010 | 2022

Call-site Sensitivity vs Object Sensitivity

- Call-site Sensitivity has been ignored

“We do not consider call-site sensitive analyses ...”
- Li et al. [2018]



CFA

The collage features several academic paper abstracts:

- A Machine-Learning Algorithm with Disjunctive M-Data-Driven Program Analysis** by Yuhao Chen, Jun Sun, Zhenjun Jiang, Junhui Liu, and Anupam Joshi.
- Making k-Object-Sensitive Pointer Analysis More Precise with Still-Limiting** by Tian Tan, Yue Li, and Jingling Xue.
- Scalability-First Pointer Analysis Self-Tuning Context Sensitivity** by Yuhao Chen, Jun Sun, and Anupam Joshi.
- Pick Your Contexts Well: Understanding Object-Sensitive Pointer Analysis** by Yuhao Chen, Jun Sun, and Anupam Joshi.
- Hybrid Context-Sensitivity for Pointers-to-Arrays** by Yuhao Chen, Jun Sun, and Anupam Joshi.
- Precision-Guided Context Sensitivity for Pointers-to-Arrays** by Yuhao Chen, Jun Sun, and Anupam Joshi.
- Introspective Analysis: Context-Sensitivity, Across the Board** by Yuhao Chen, Jun Sun, and Anupam Joshi.

1981

2002

2010

2022

Call-site Sensitivity vs Object Sensitivity

- Call-site Sensitivity has been ignored

“We have included $2cs+h$ to demonstrate the superiority of object sensitivity over call-site sensitivity”

- Tan et al. [2016]

The collage features several research papers:

- A Machine Learning Algorithm with Disjunctive M...** (Data-Driven Program Analysis)
- Making k-Object-Sensitive Pointer Ana...** (More Precise with Still k-Limiting)
- Scalability-First: Pointer Ana...** (Self-Tuning Context-Sen)
- Pick Your Contexts Well: Understanding Object-Sen...** (The Making of a Precise and Scalable Pointer Analysis)
- Hybrid Context-Sensitivity for Points-To A...**
- Precision-Guided Context Sensitivity for Point...**
- Introspective Analysis: Context-Sensitivity, Across the Board**



CFA

1981

2002

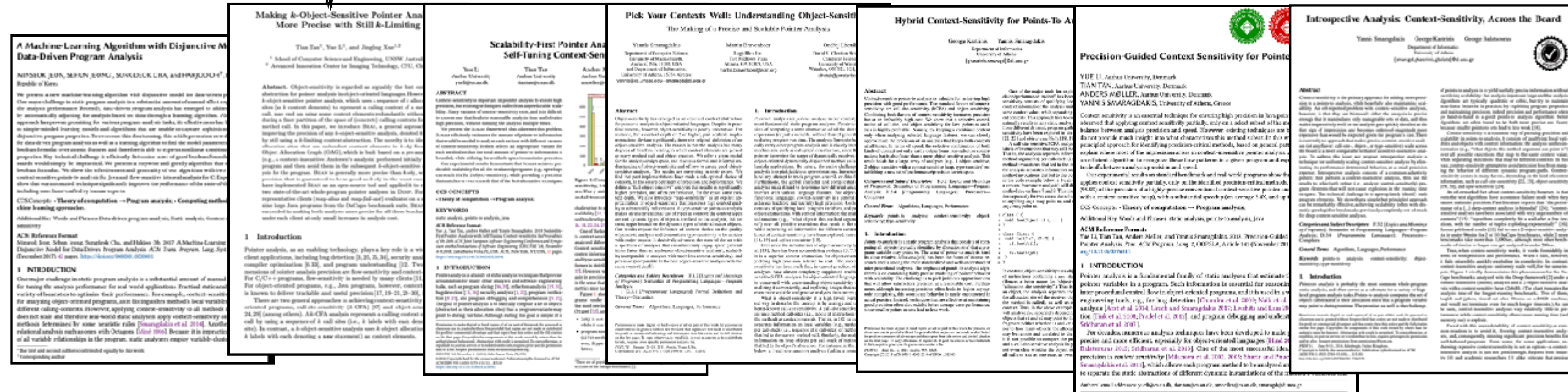
2010

2022

Call-site Sensitivity vs Object Sensitivity

- Call-site Sensitivity has been ignored

“... we do not discuss our approach for call-site sensitivity” - Jeon et al. [2019]



CFA

1981

2002

2010

2022

Call-site Sensitivity vs Object Sensitivity

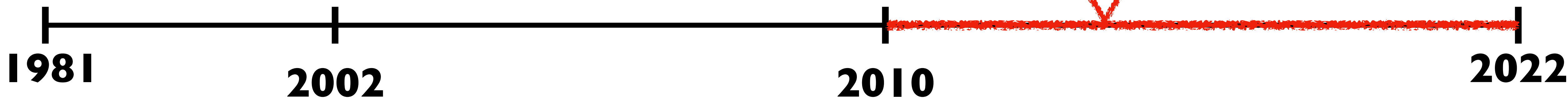
- Call-site Sensitivity has been ignored

“... we do not discuss our approach for call-site sensitivity”
- Jeon et al. [2019]

I also strongly dismissed call-site sensitivity

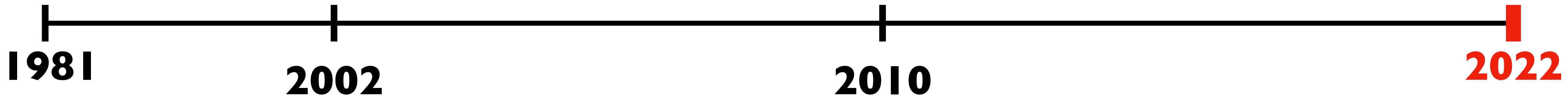


CFA



Call-site Sensitivity vs Object Sensitivity

Currently, call-site sensitivity is known as a bad context



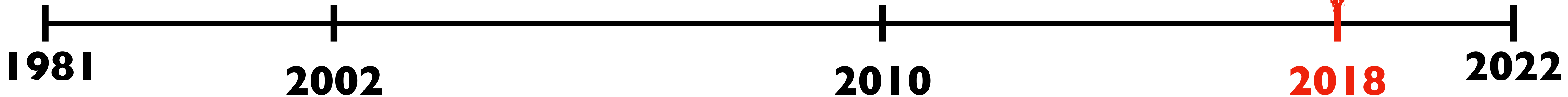
Call-site Sensitivity vs Object Sensitivity

A technique **context tunneling** is proposed



Context tunneling can improve both **call-site sensitivity** and **object sensitivity**

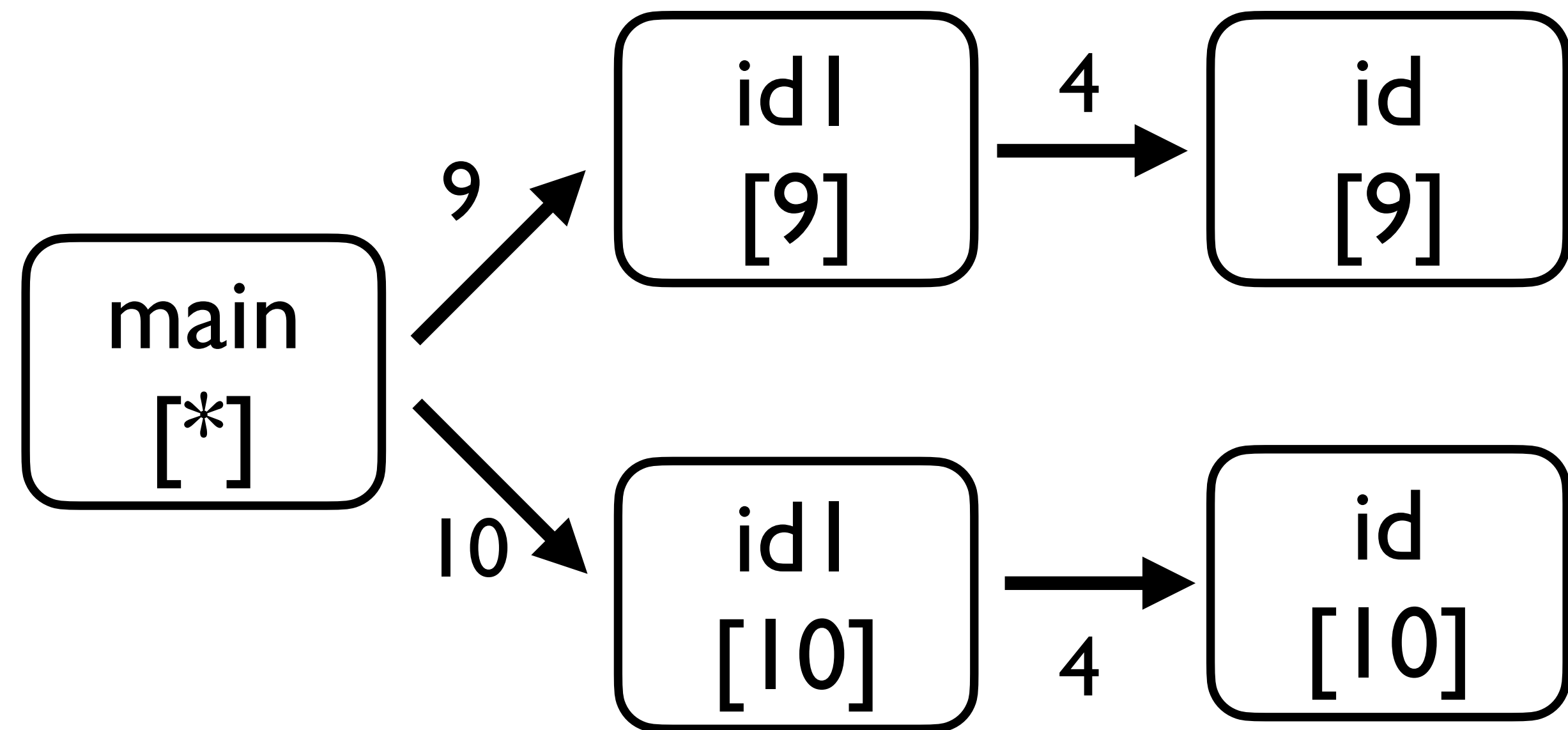
Jeon et al. [2018]



Call-site Sensitivity vs Object Sensitivity

- **Context tunneling** can remove the limitation of call-site sensitivity

```
0: class C{
1:   id(v){
2:     return v;}
3:   idl(v){
4:     return id0(v);}
5: }
6: main(){
7:   c1 = new C();//C1
8:   c2 = new C();//C2
9:   a = (A) c1.idl(new A());//query1
10:  b = (B) c2.idl(new B());//query2
11: }
```

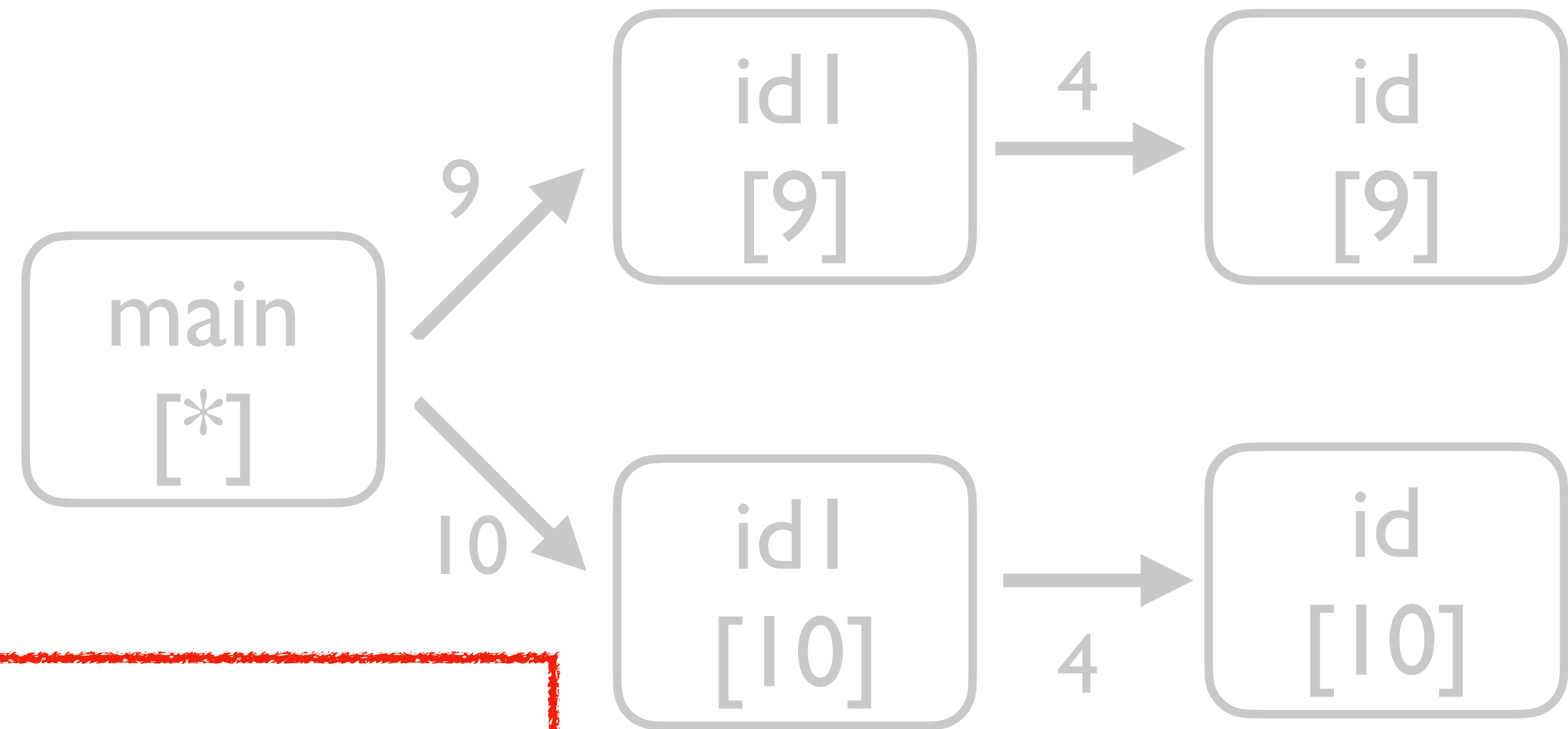


I-CFA with context tunneling
($T = \{4\}$)

Call-site Sensitivity vs Object Sensitivity

- Context tunneling can remove the limitation of call-site sensitivity

```
0: class C{
1:   id(v){
2:     return v;}
3:   idl(v){
4:     return id0(v);}
5: }
6: main(){
7:   c1 = new C();//C1
```



Tunneling abstraction:

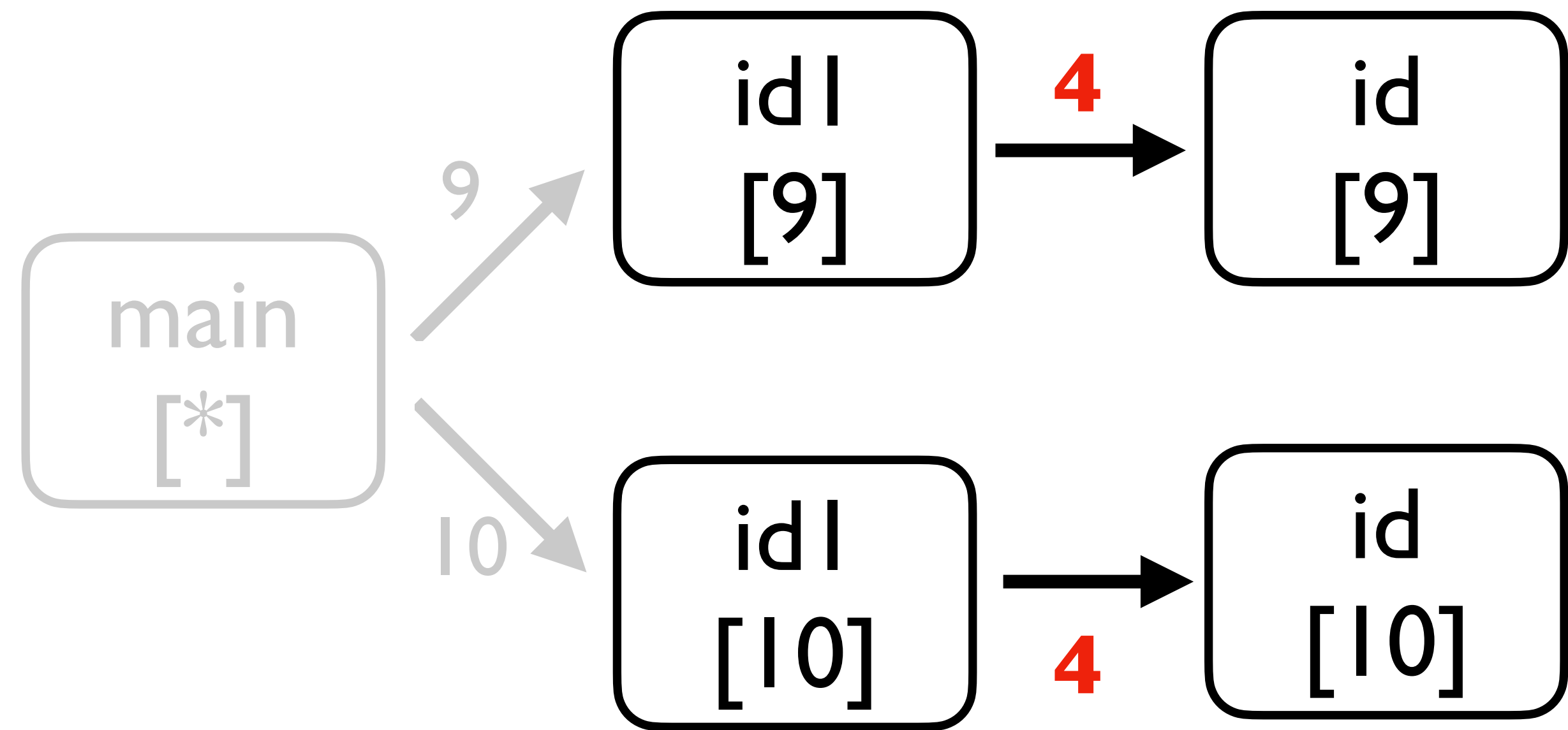
Determines where to apply context tunneling

with context tunneling
($T = \{4\}$)

Call-site Sensitivity vs Object Sensitivity

- Context tunneling can remove the limitation of call-site sensitivity

```
0: class C{
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2:     return v;}
3:   idl(v){
4:     return id0(v);}
5: }
6: main(){
7:   c1 = new C();//C1
8:   c2 = new C();//C2
9:   a = (A) c1.idl(new A());//query1
10:  b = (B) c2.idl(new B());//query2
```



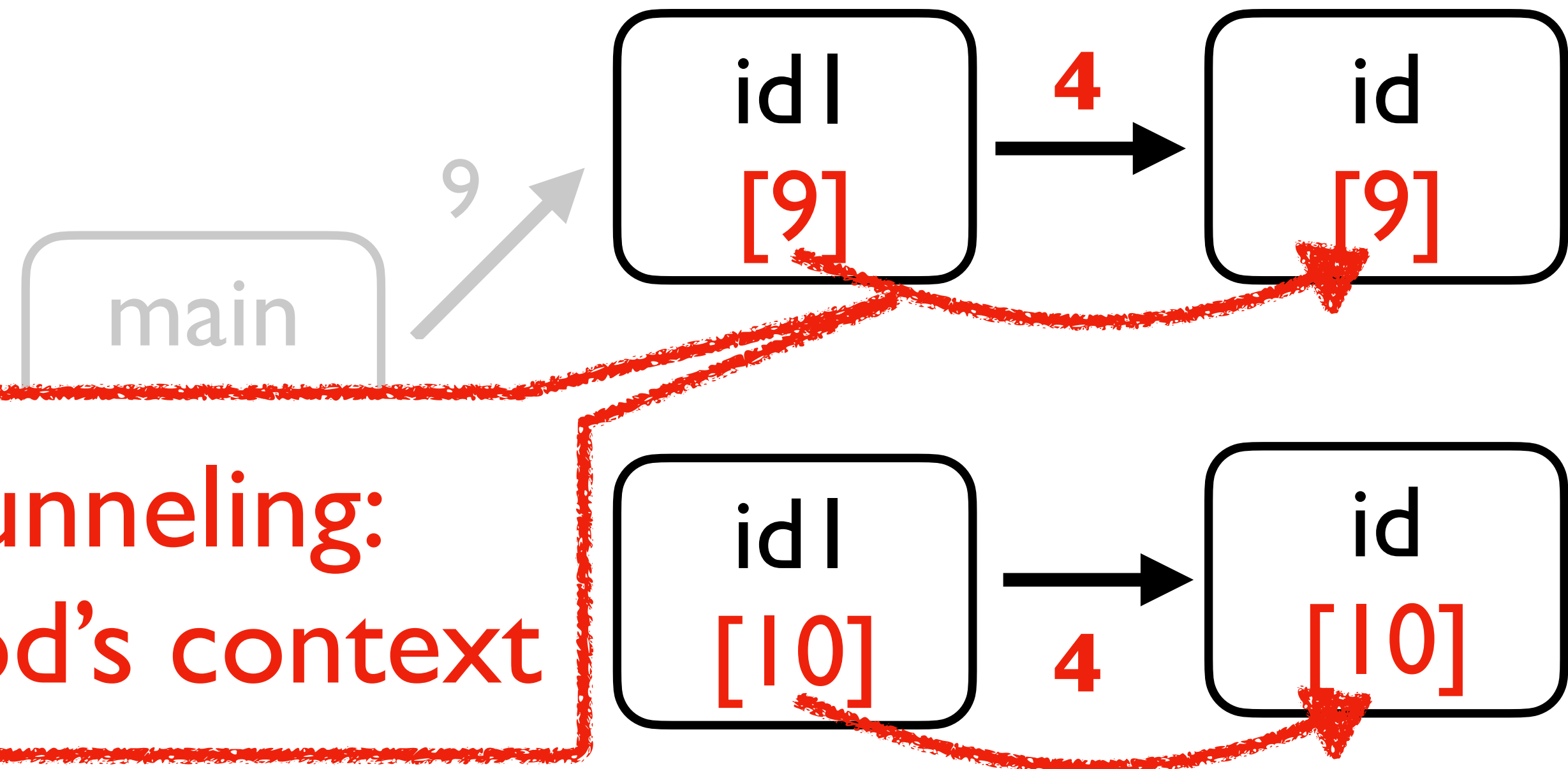
I-CFA with context tunneling
($T = \{4\}$)

Unimportant call-sites that should not be used as context elements

Call-site Sensitivity vs Object Sensitivity

- Context tunneling can remove the limitation of call-site sensitivity

```
0: class C{
1:   id(v){
2:     return v;}
3:   idI(v){
4:     return id0(v);}
5: }
6: main
7: c1 = (A) new C();//query1
8: c2 = (B) new C();//query2
9: a = (A) c1.idI(new A());//query1
10: b = (B) c2.idI(new B());//query2
11: }
```



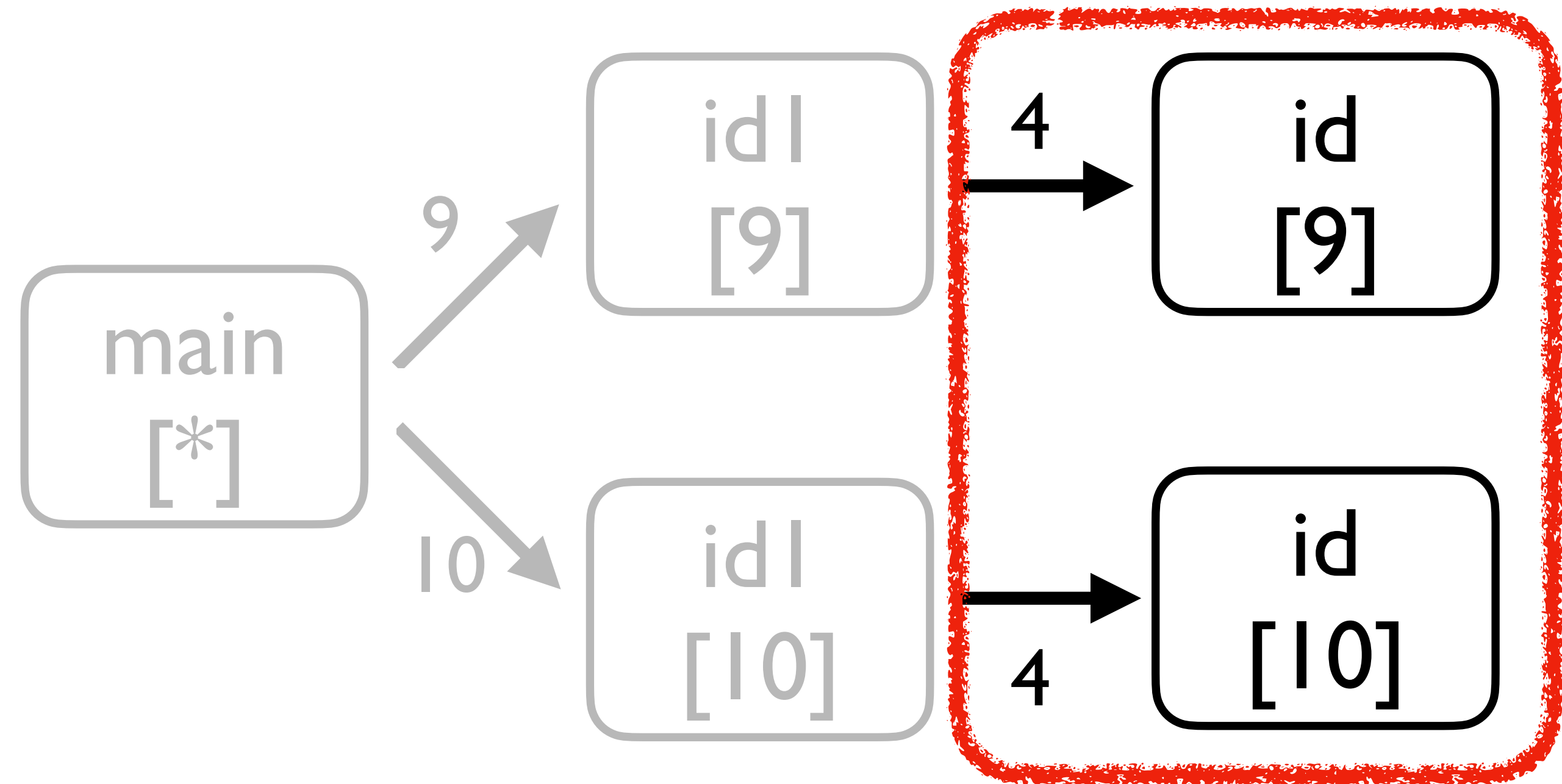
Apply context tunneling:
Inherit caller method's context

I-CFA with context tunneling
(T = {4})

Call-site Sensitivity vs Object Sensitivity

- **Context tunneling** can remove the limitation of call-site sensitivity

```
0: class C{
1:   id(v){
2:     return v;}
3:   idl(v){
4:     return id0(v);}
5: }
6: main(){
7:   c1 = new C();//C1
8:   c2 = new C();//C2
9:   a = (A) c1.idl(new A());//query1
10:  b = (B) c2.idl(new B());//query2
11: }
```



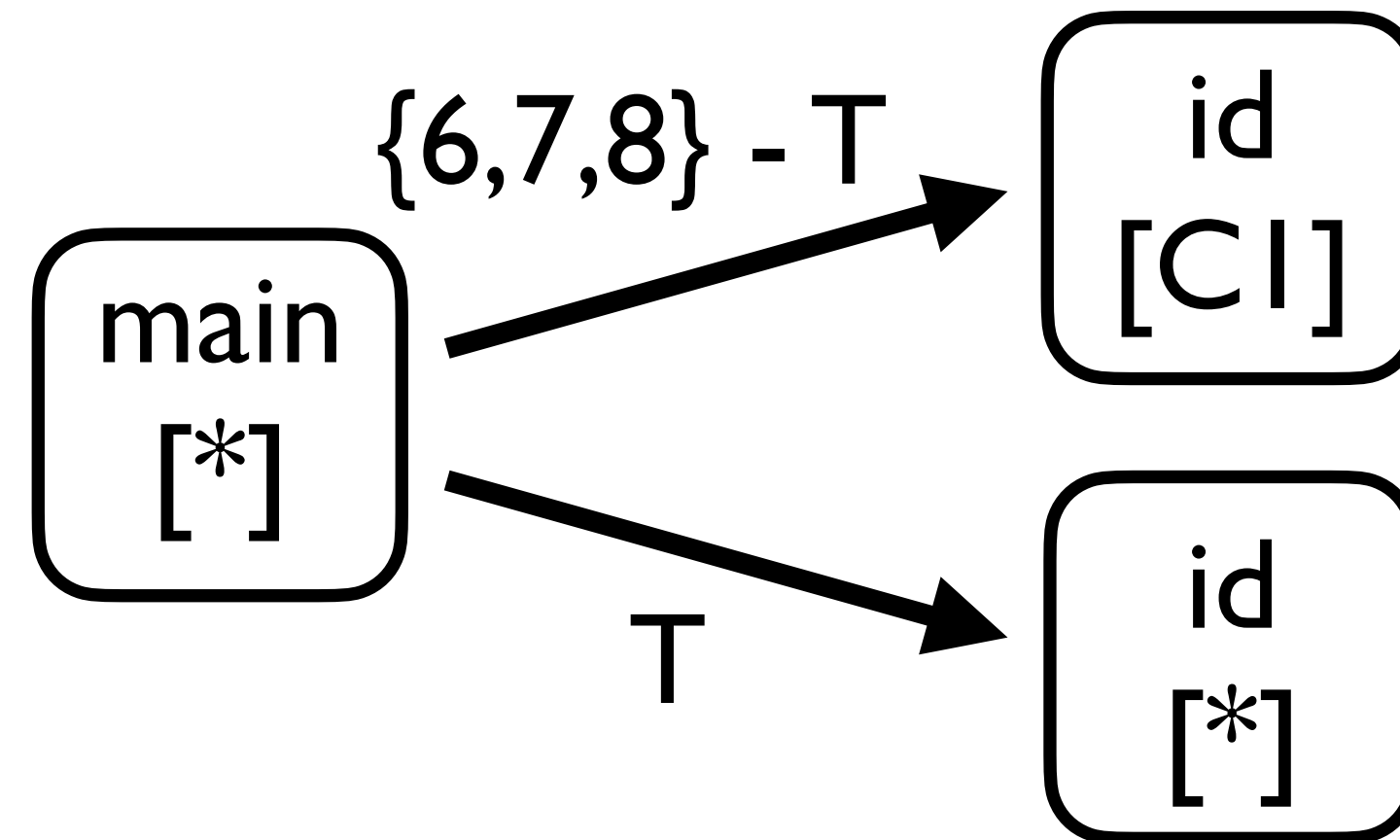
I-CFA with context tunneling
($T = \{4\}$)

With tunneling, I-CFA separates the nested method calls

Call-site Sensitivity vs Object Sensitivity

- Object sensitivity still suffers from its **limitation**

```
0: class C{
1:   id(v){
2:     return v;}
3: }
4: main(){
5:   c1 = new C();//C1
6:   a = (A) c1.id(new A());
7:   b = (B) c1.id(new B());
8:   c = (C) c1.id(new C());
9: }
```



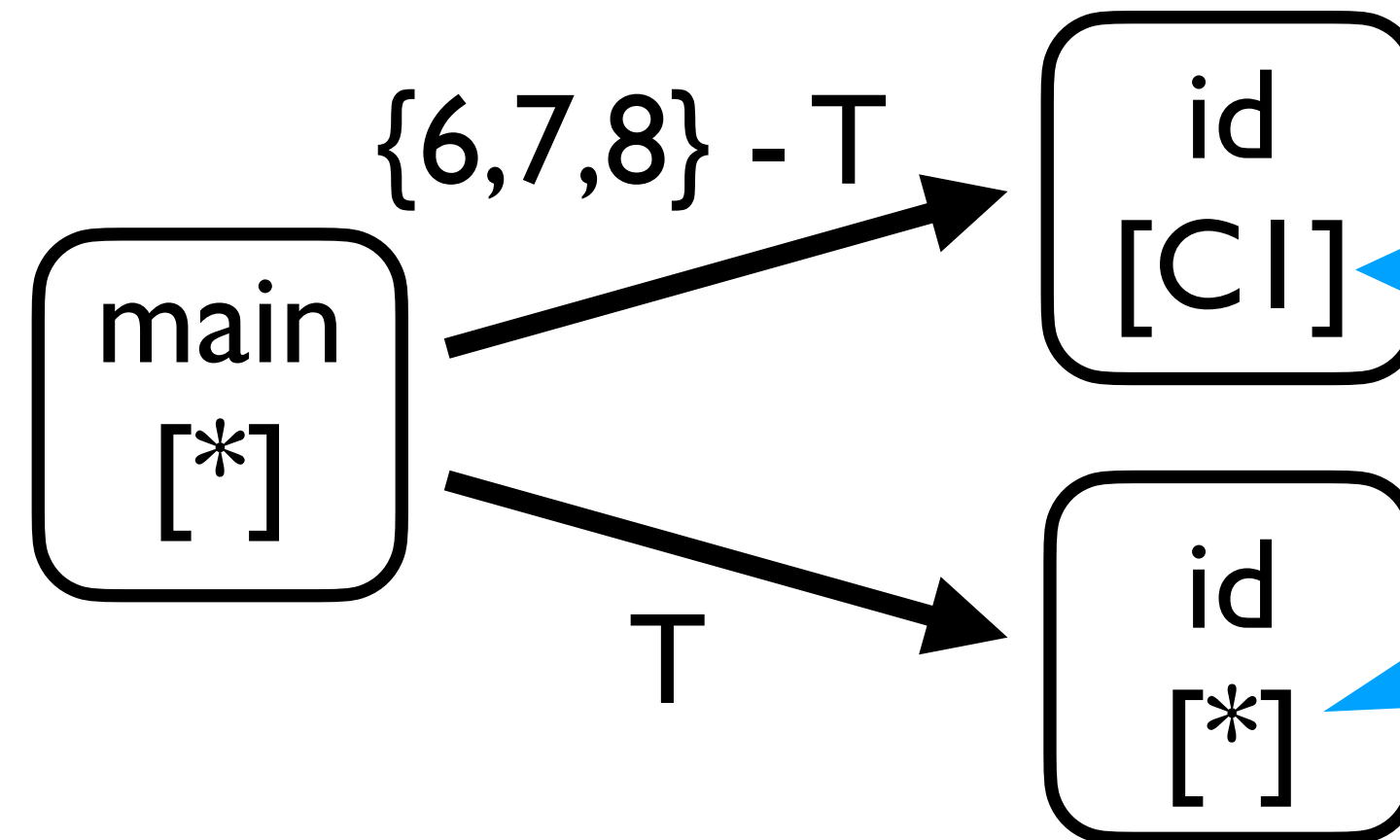
Call-graph of I-Obj with tunneling T

I-Obj + Tunneling
(T = ?)

Call-site Sensitivity vs Object Sensitivity

- Object sensitivity still suffers from its **limitation**

```
0: class C{
1:   id(v){
2:     return v;}
3: }
4: main(){
5:   cI = new C();//CI
6:   a = (A) cI.id(new A());
7:   b = (B) cI.id(new B());
8:   c = (C) cI.id(new C());
9: }
```



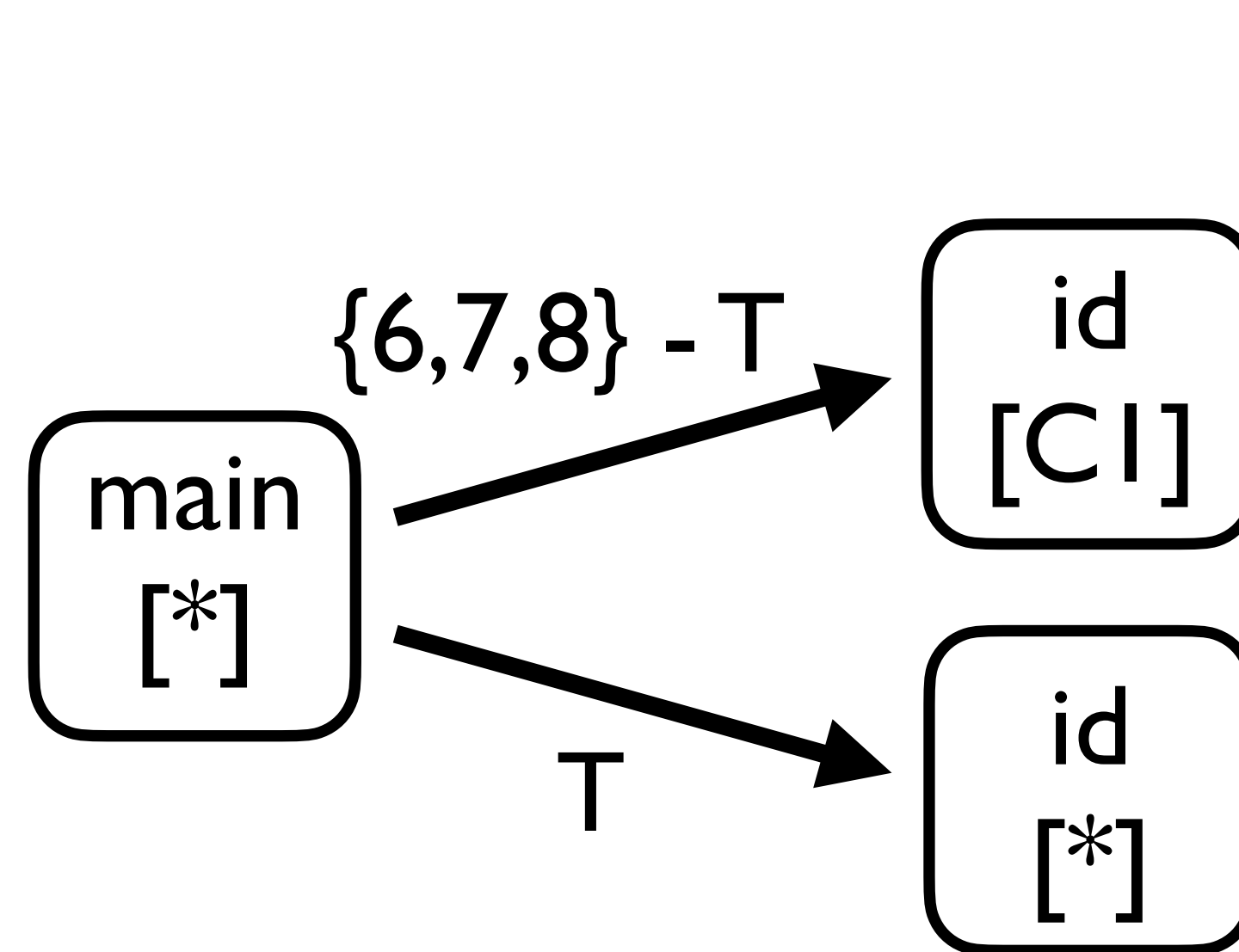
Unable to separate the three method calls with the two contexts

I-Obj + Tunneling
(T = ?)

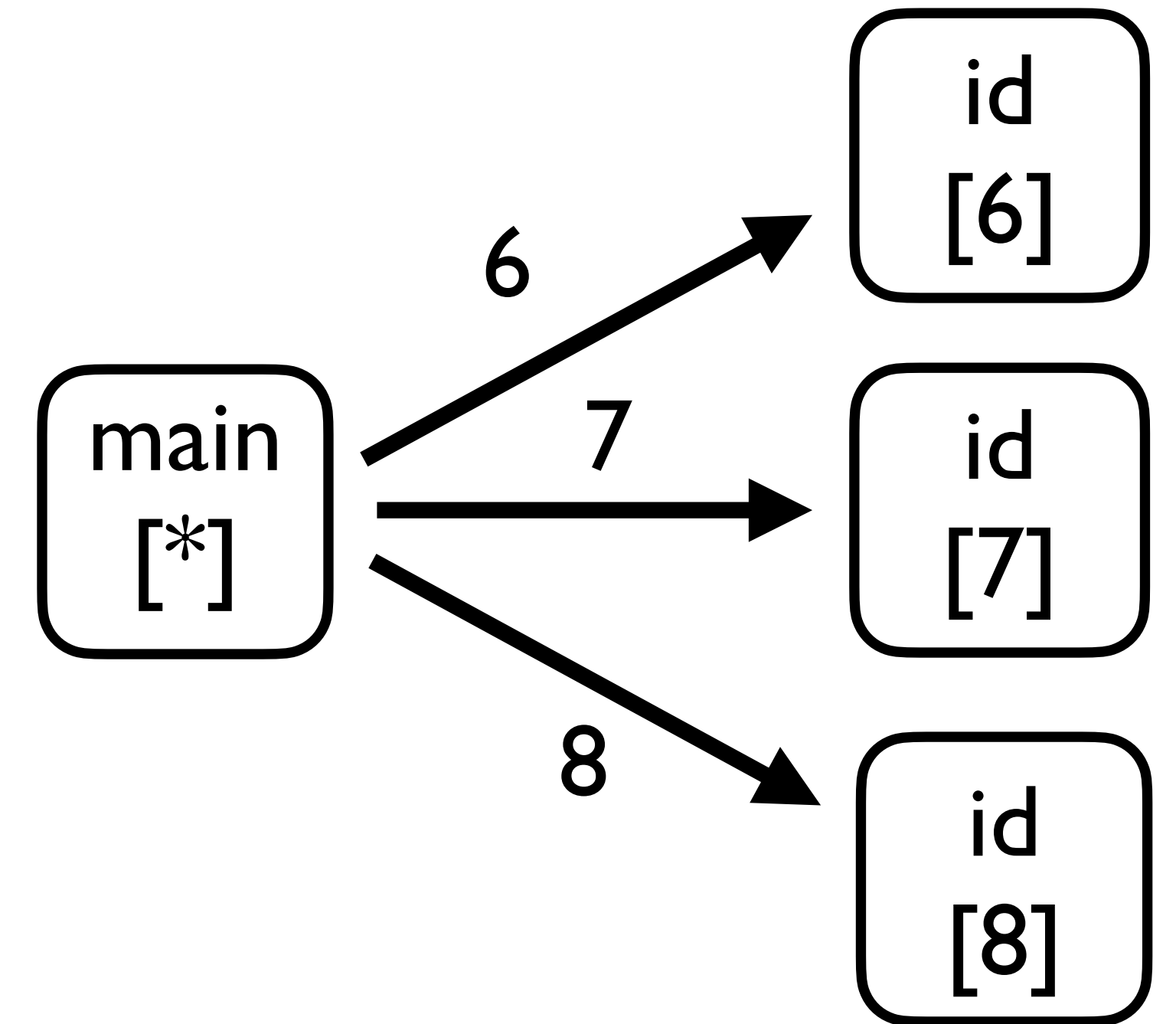
Call-site Sensitivity vs Object Sensitivity

- Object sensitivity still suffers from its **limitation**

```
0: class C{
1:   id(v){
2:     return v;}
3: }
4: main(){
5:   c1 = new C();//C1
6:   a = (A) c1.id(new A());
7:   b = (B) c1.id(new B());
8:   c = (C) c1.id(new C());
9: }
```



I-Obj + Tunneling
(T = ?)



I-CFA

Call-site sensitivity easily separates the three method calls

Call-site Sensitivity vs Object Sensitivity

- Object sensitivity still suffers from its limitation

Observation

When context tunneling is included

- Limitation of call-site sensitivity is **removed**
- Limitation of object sensitivity is **not removed**

```
0: c
1:
2:
3: }
4: n
5:
6: a
7: b = (B) cl.id(new B());
8: c = (B) cl.id(new C());
9: }
```

Call-site Sensitivity vs Object Sensitivity

- Object sensitivity still suffers from its limitation

Observation

When context tunneling is included

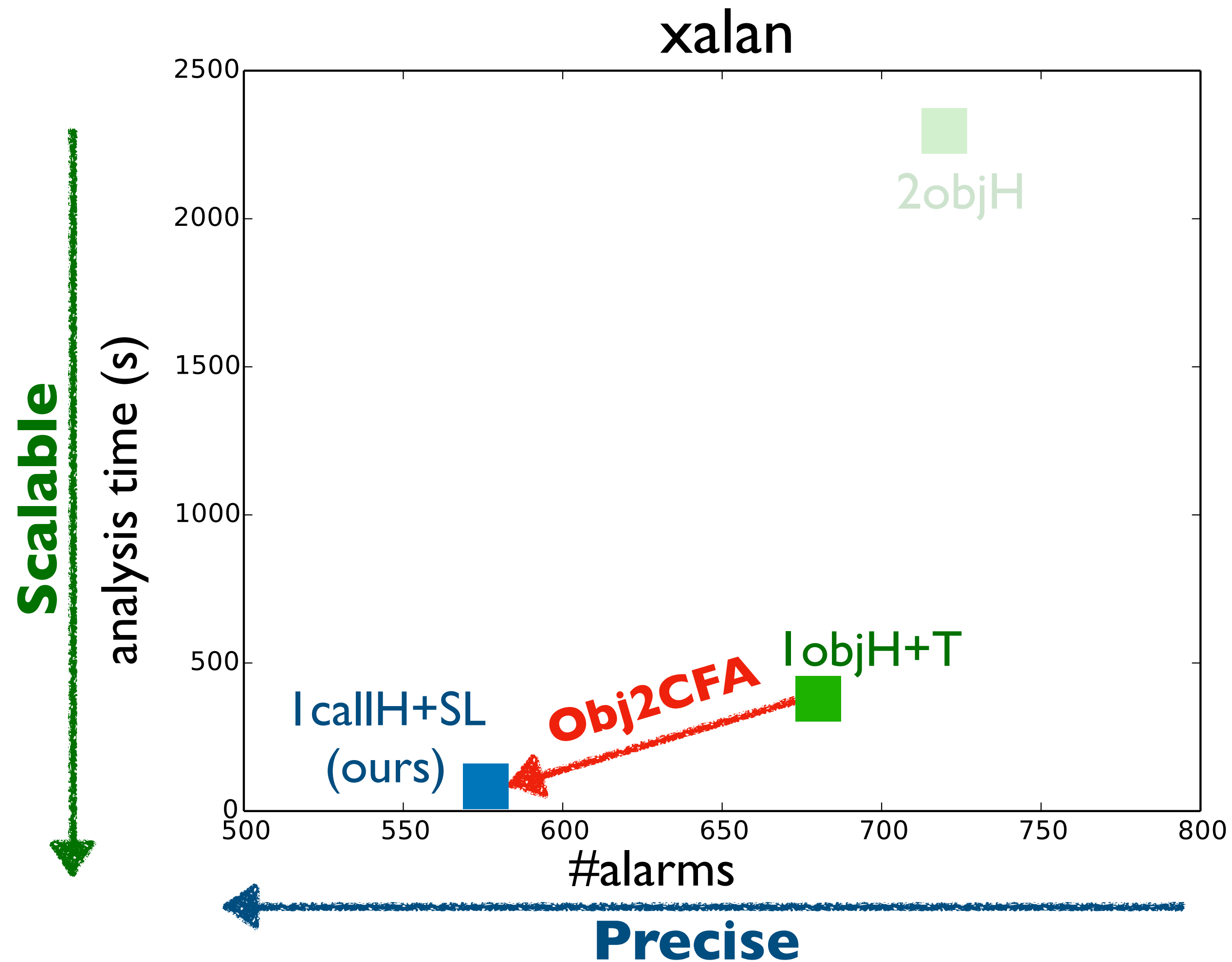
- Limitation of call-site sensitivity is **removed**
- Limitation of object sensitivity is **not removed**

Our claim

If context tunneling is included,
call-site sensitivity is more precise than object sensitivity

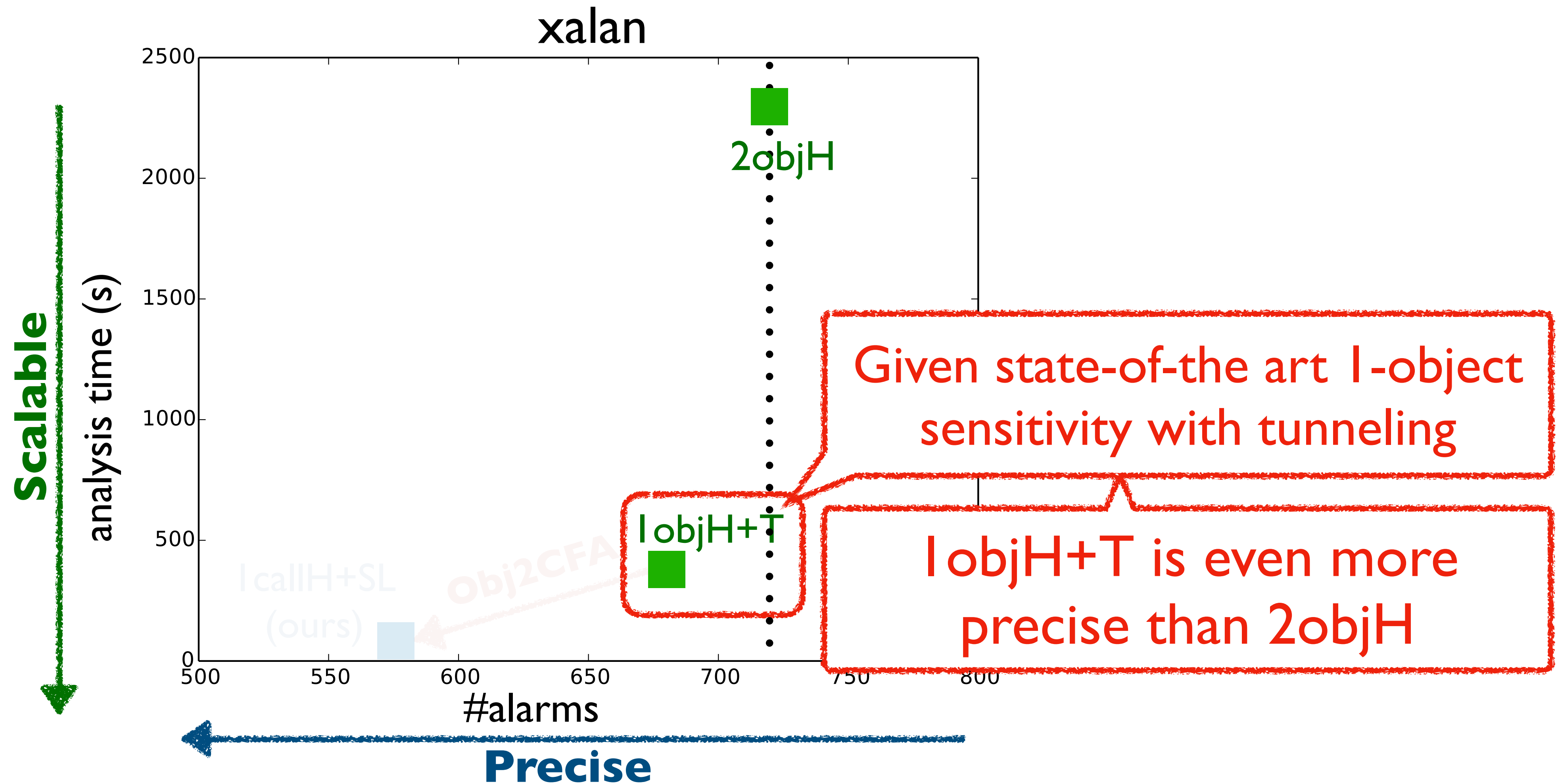
Our Technique : **Obj2CFA**

- **Obj2CFA** transforms a given **object sensitivity** into a more precise **CFA**



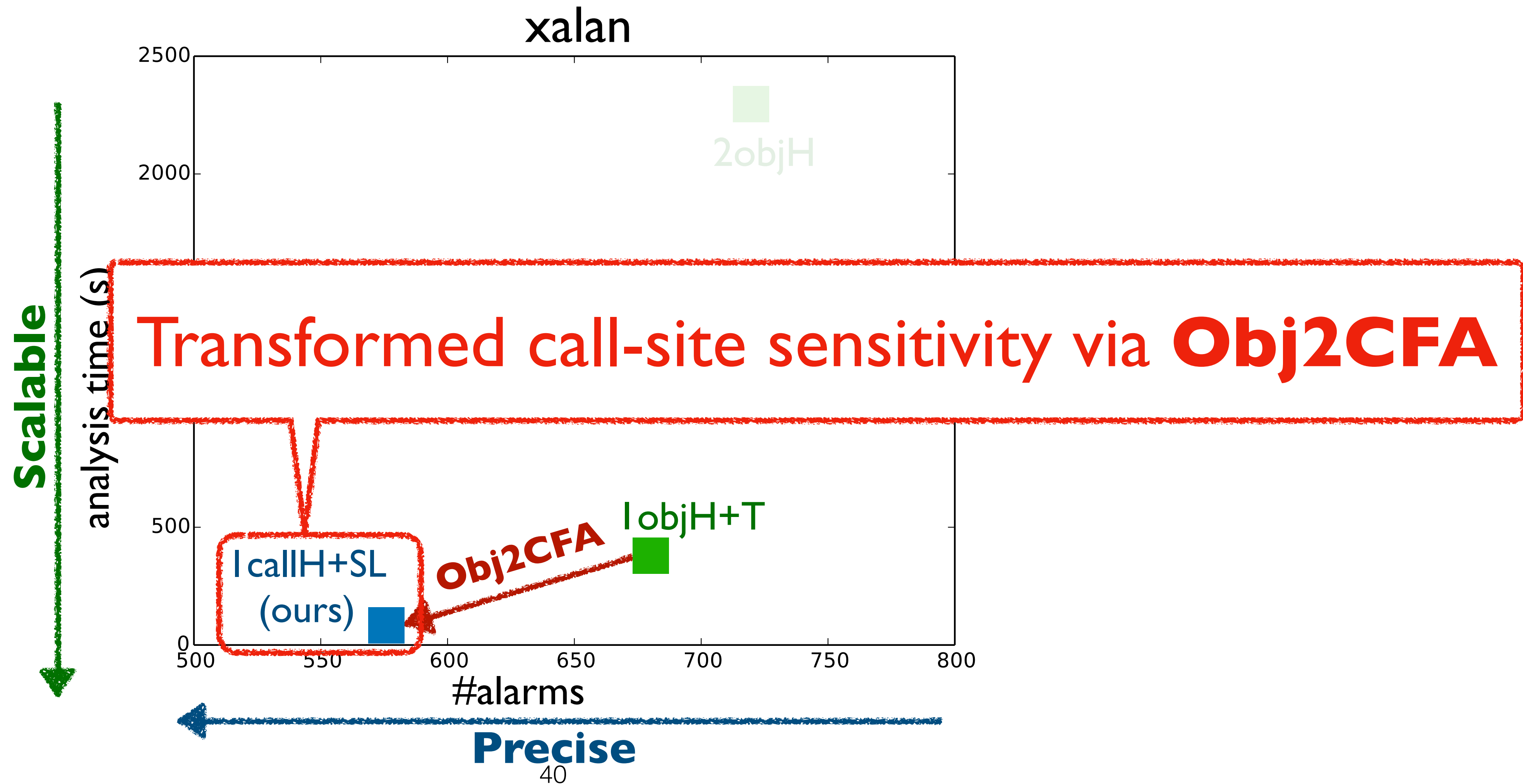
Our Technique : **Obj2CFA**

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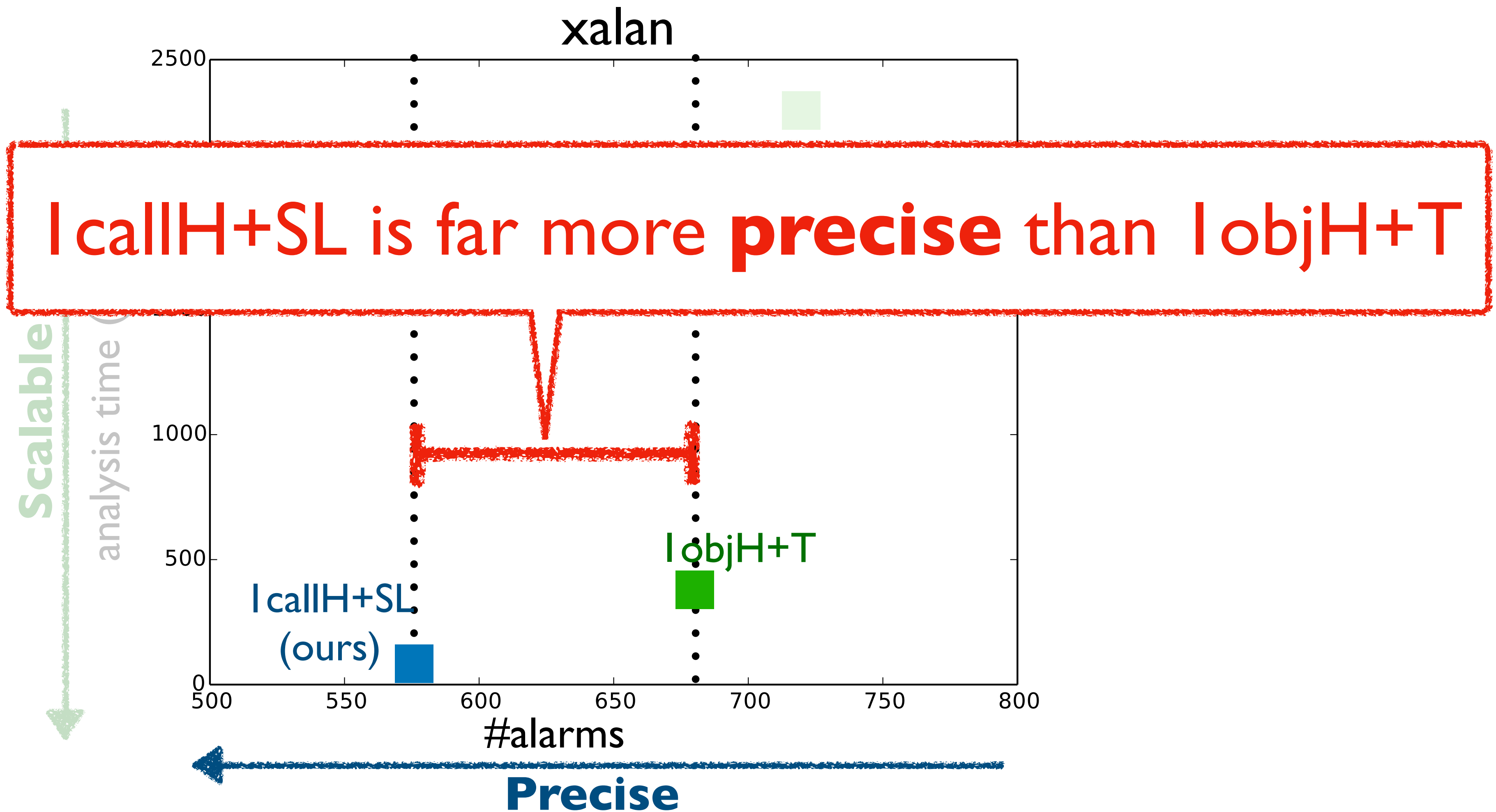
Our Technique : **Obj2CFA**

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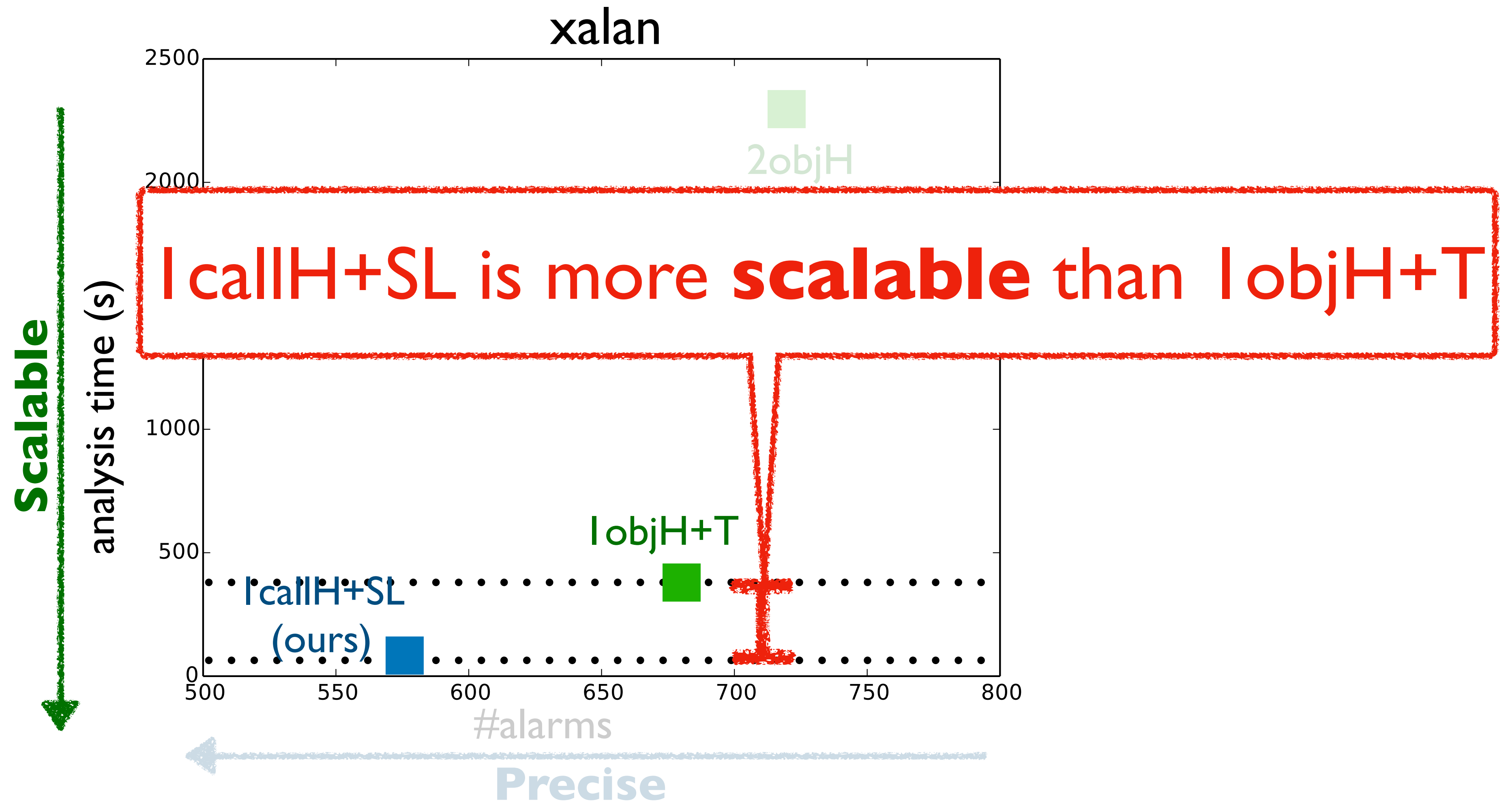
Our Technique : **Obj2CFA**

- **Obj2CFA** transforms a given **object sensitivity** into a more precise **CFA**



Our Technique : **Obj2CFA**

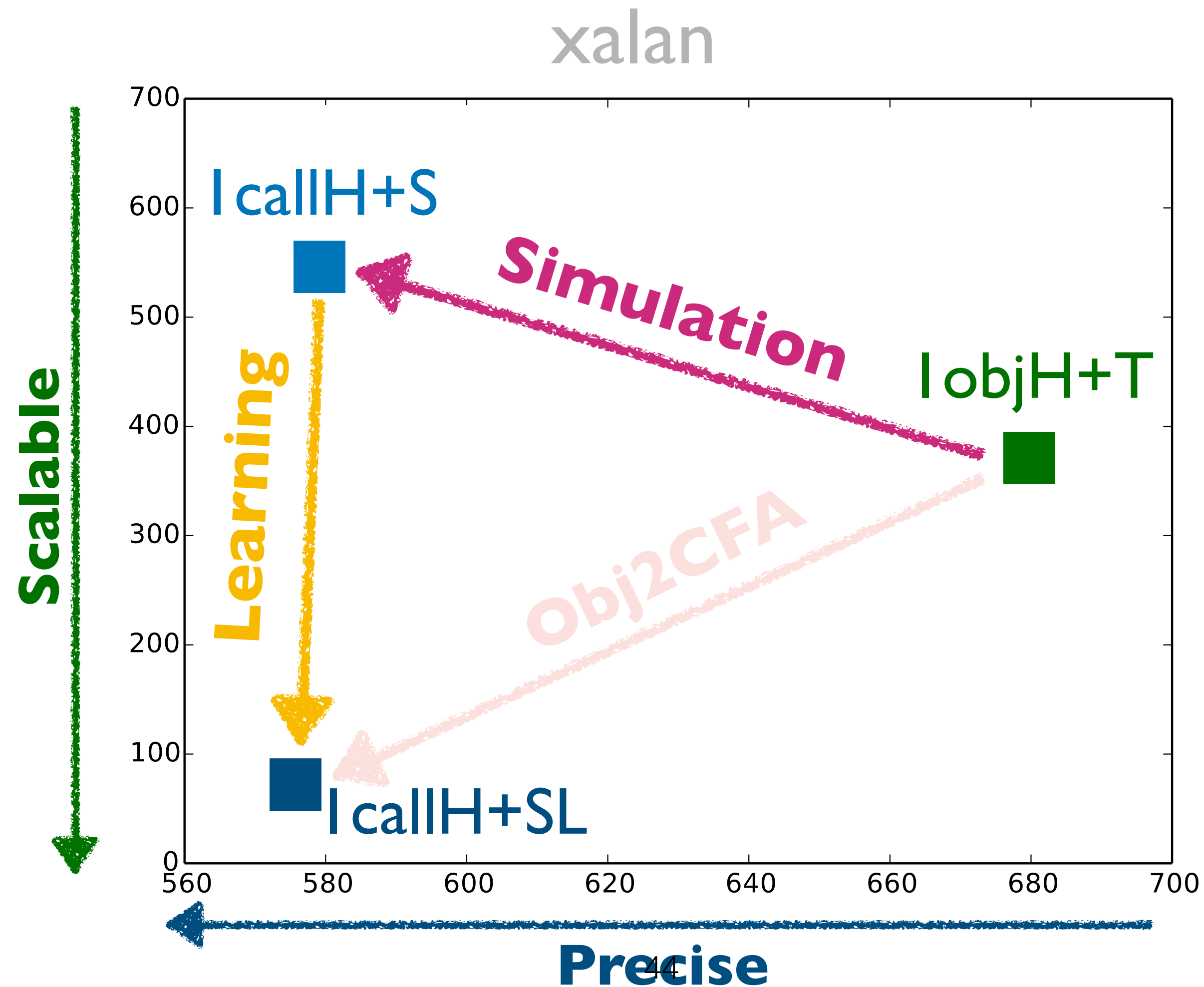
- **Obj2CFA** transforms a given **object sensitivity** into a more precise **CFA**



Detail of Obj2CFA

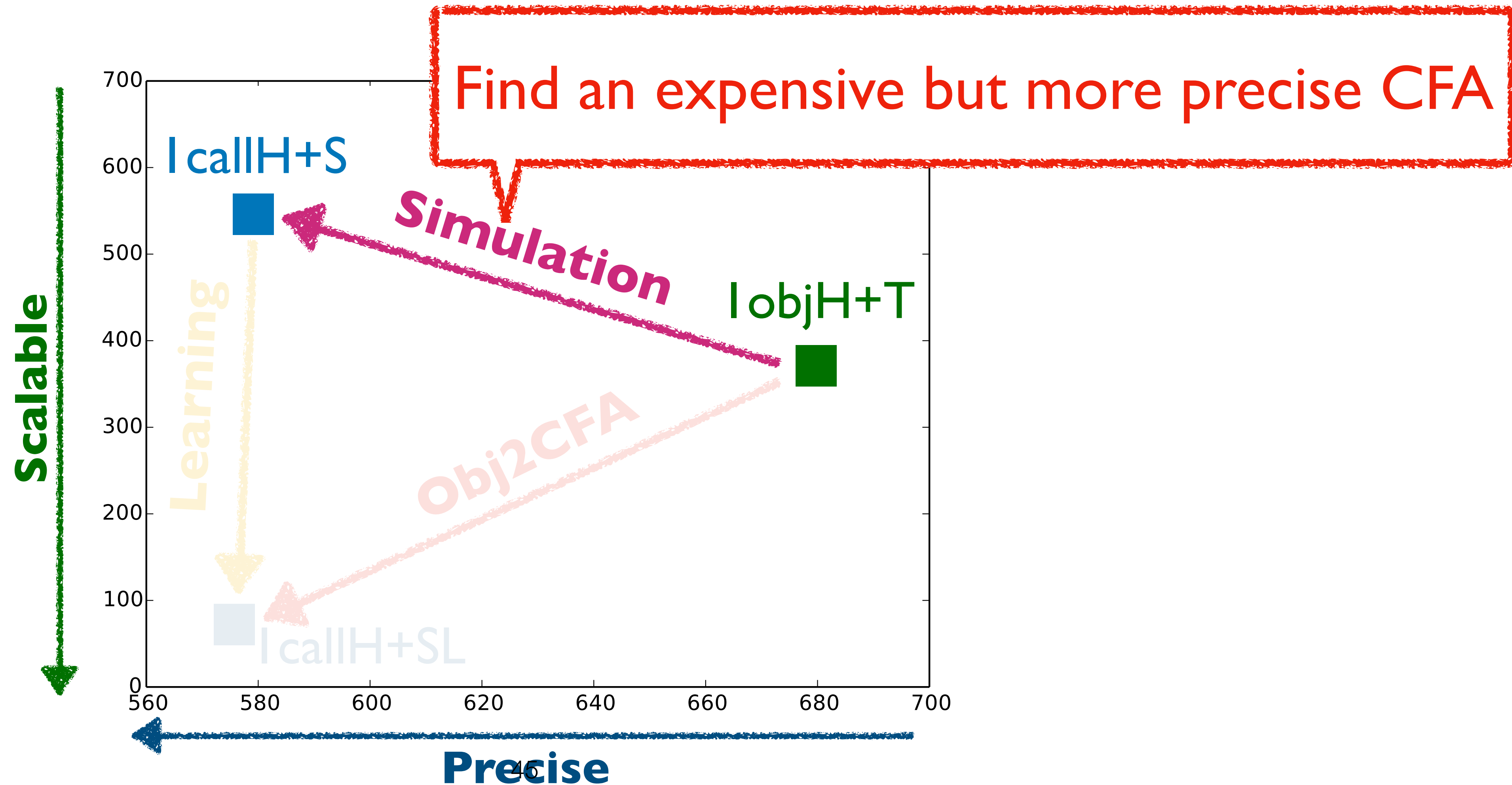
Our Technique : **Obj2CFA**

- **Obj2CFA** consists of **simulation** and simulation-guided **learning**



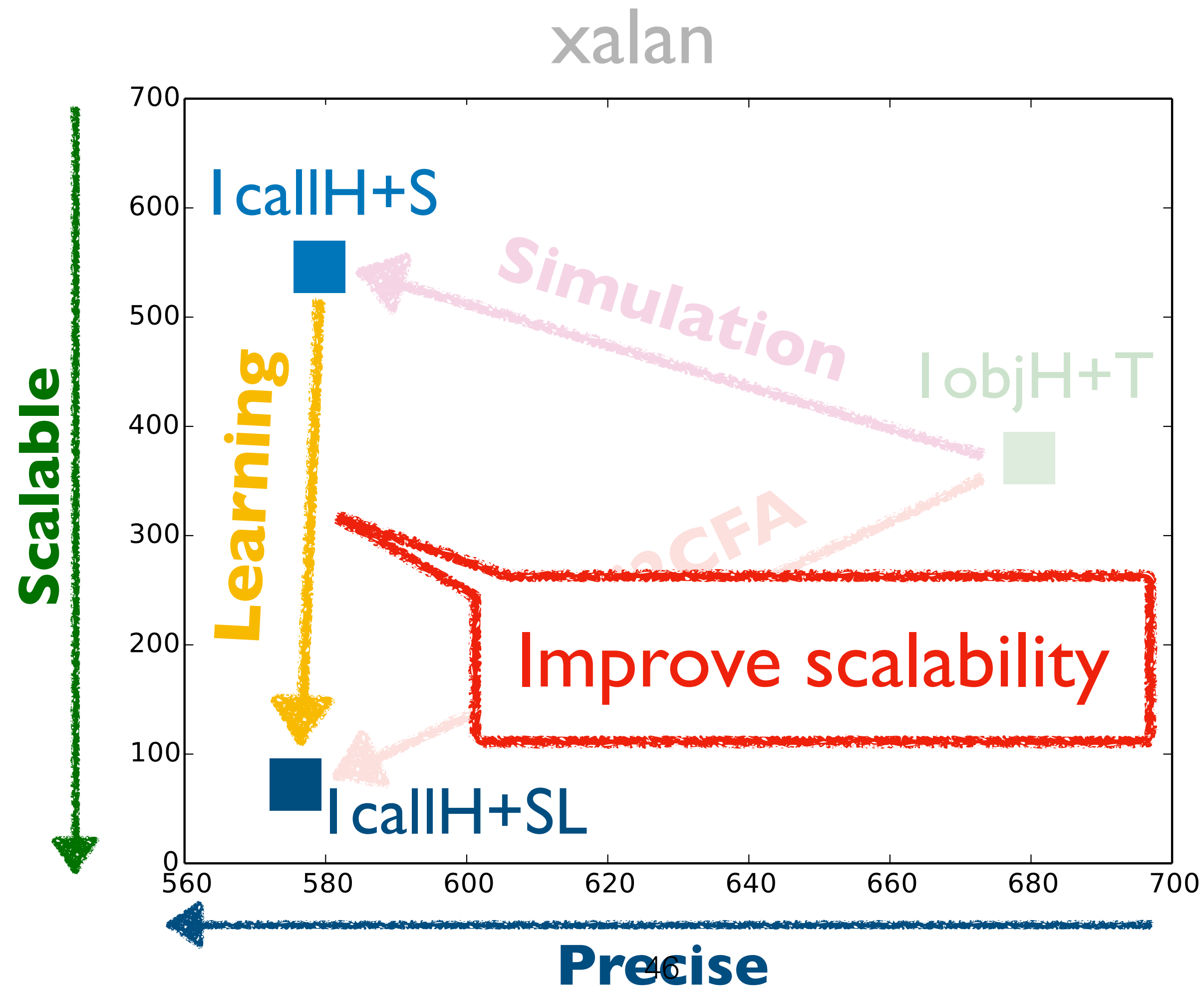
Our Technique : **Obj2CFA**

- **Obj2CFA** consists of **simulation** and simulation-guided **learning**



Our Technique : **Obj2CFA**

- **Obj2CFA** consists of **simulation** and simulation-guided **learning**



Technique I: Simulation

- Running example to illustrate Simulation

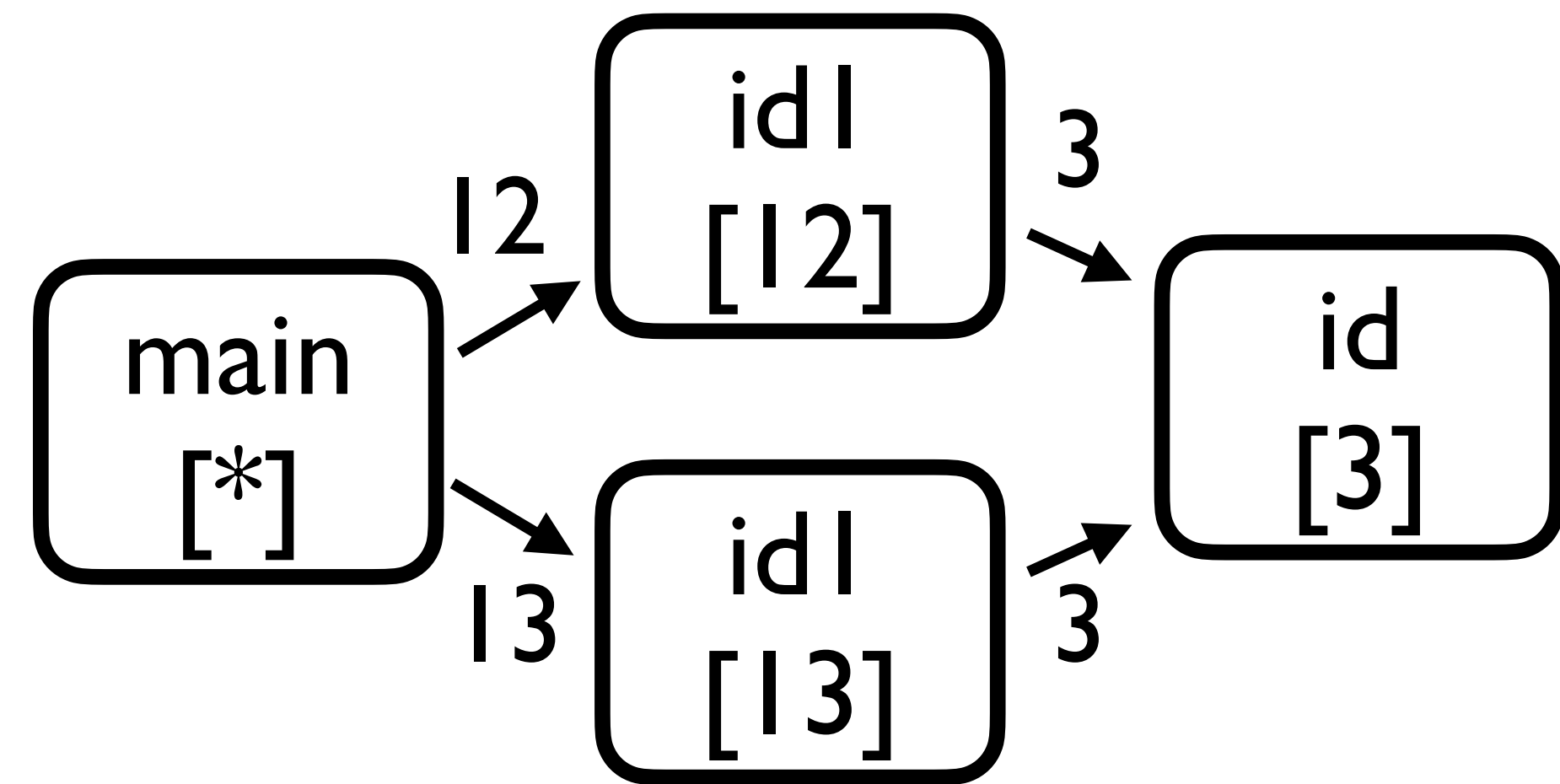
```
1: class C{
2:   id(v){return v;}
3:   idl(v){return id(v);}
4:   foo(){
5:     A a = (A) this.id(new A());}//query1
6:     B b = (B) this.id(new B());}//query2
7: }
8: main(){
9:   c1 = new C(); //C1
10:  c2 = new C(); //C2
11:  c3 = new C(); //C3
12:  A a = (A) c1.idl(new A()); //query3
13:  B b = (B) c2.idl(new B()); //query4
14:  c3.foo();
15: }
```

Technique I: Simulation

- Running example to illustrate Simulation

```
1: class C{
2:   id(v){return v;}
3:   idl(v){return id(v);}
4:   foo(){
5:     A a = (A) this.id(new A());} //query1
6:     B b = (B) this.id(new B());} //query2
7: }
8: main(){
9:   c1 = new C(); //C1
10:  c2 = new C(); //C2
11:  c3 = new C(); //C3
12:  A a = (A) c1.idl(new A()); //query3
13:  B b = (B) c2.idl(new B()); //query4
14:  c3.foo();
15: }
```

Limitation of conventional I-CFA



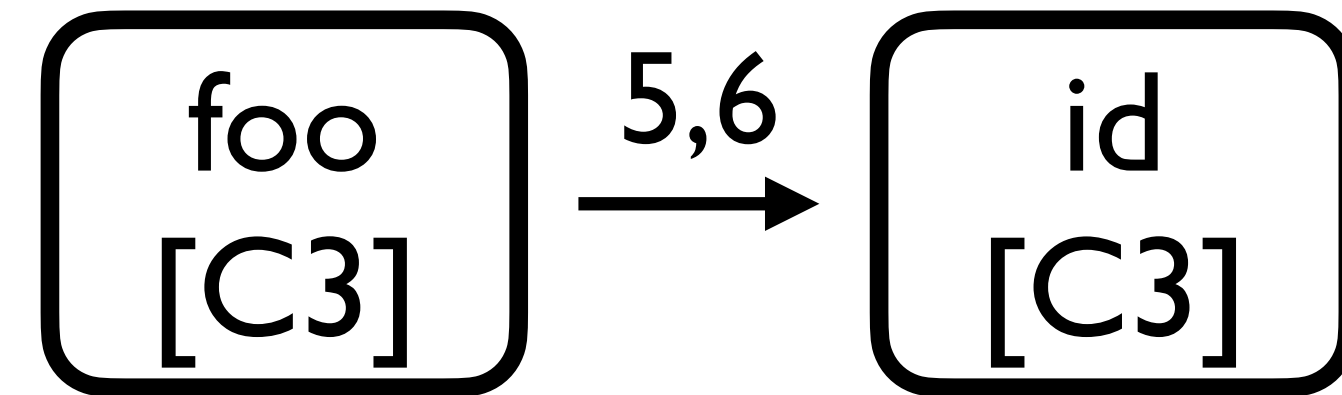
Technique I: Simulation

- Running example to illustrate Simulation

```
1: class C{
2:   id(v){return v;}
3:   idl(v){return id(v);}
4:   foo(){
5:     A a = (A) this.id(new A()); //query1
6:     B b = (B) this.id(new B()); //query2
7:   }
8: main(){
9:   c1 = new C(); //C1
10:  c2 = new C(); //C2
11:  c3 = new C(); //C3
12:  A a = (A) c1.idl(new A()); //query3
13:  B b = (B) c2.idl(new B()); //query4
14:  c3.foo();
15: }
```



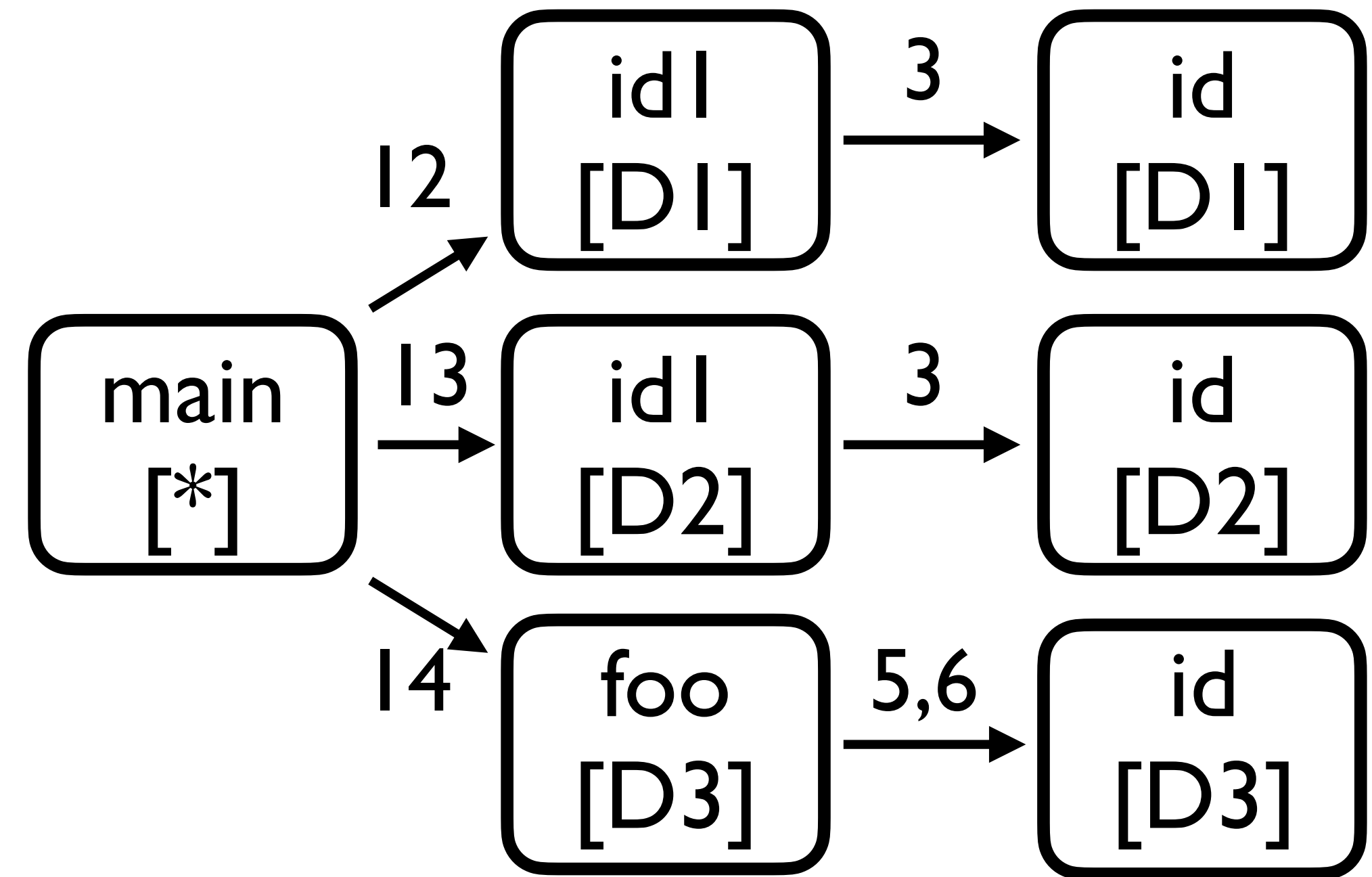
Limitation of object sensitivity



Technique 1: Simulation

- Given **object sensitivity** is conventional **l-object sensitivity** (e.g., $T = \emptyset$)

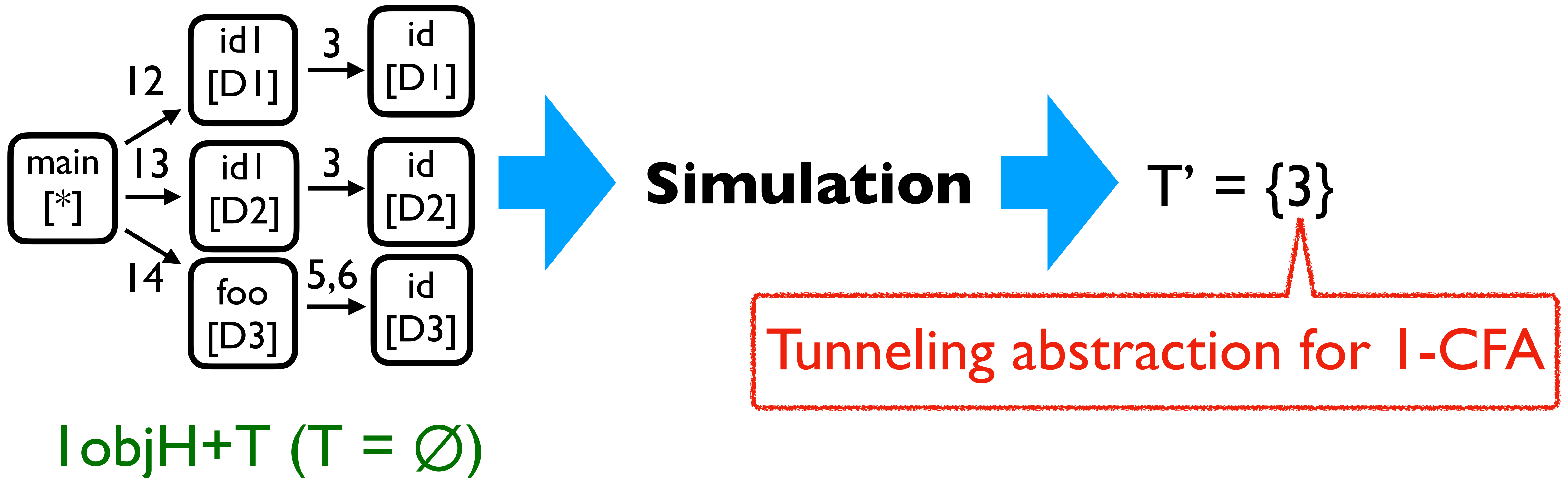
```
1: class C{
2:   id(v){return v;}
3:   idl(v){return id(v);}
4:   foo(){
5:     A a = (A) this.id(new A());} //query1
6:     B b = (B) this.id(new B());} //query2
7: }
8: main(){
9:   c1 = new C(); //C1
10:  c2 = new C(); //C2
11:  c3 = new C(); //C3
12:  A a = (A) c1.idl(new A()); //query3
13:  B b = (B) c2.idl(new B()); //query4
14:  c3.foo();
15: }
```



l objH+T (T = ∅)

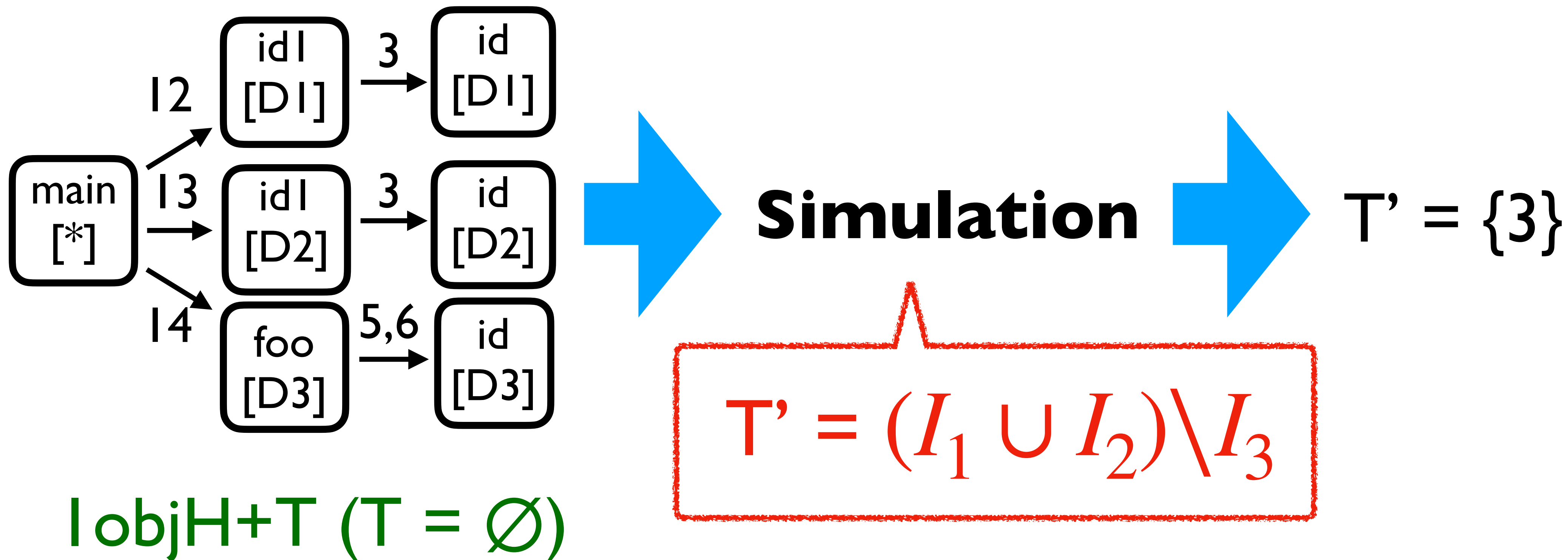
Technique 1: Simulation

- **Simulation** takes a call-graph and produce a tunneling abstraction for CFA



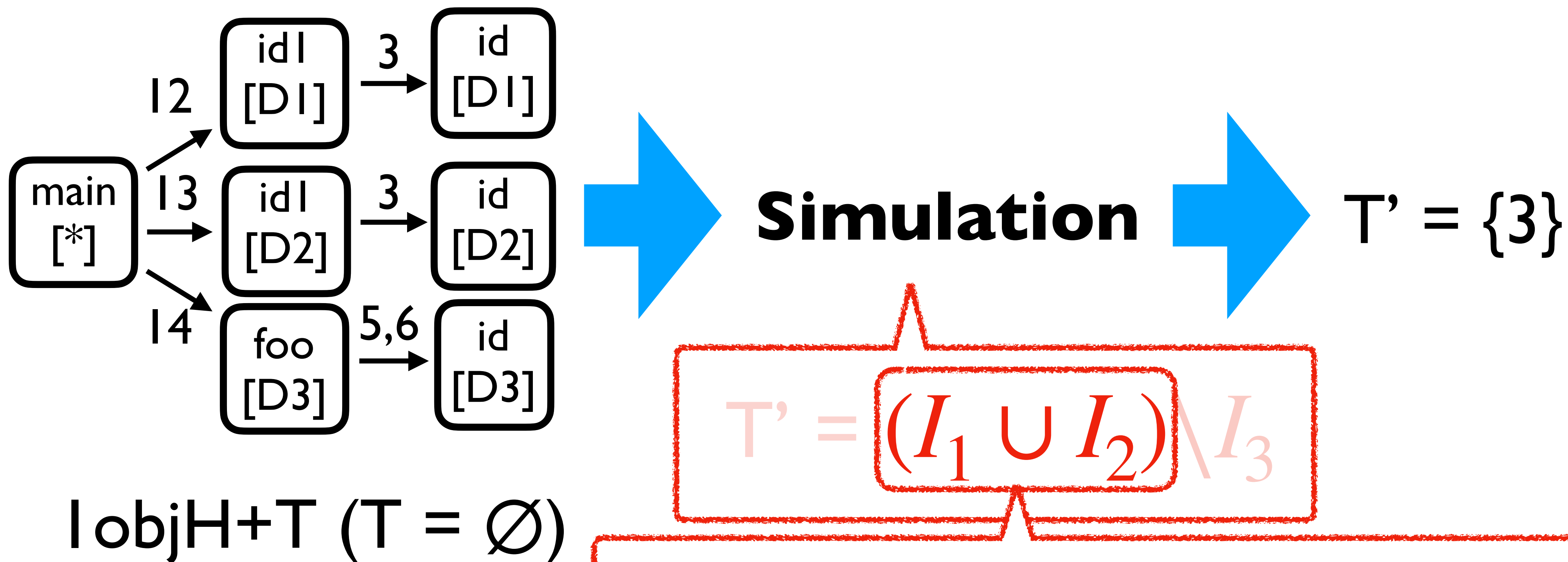
Technique 1: Simulation

- **Simulation** takes a call-graph and produce a tunneling abstraction for CFA



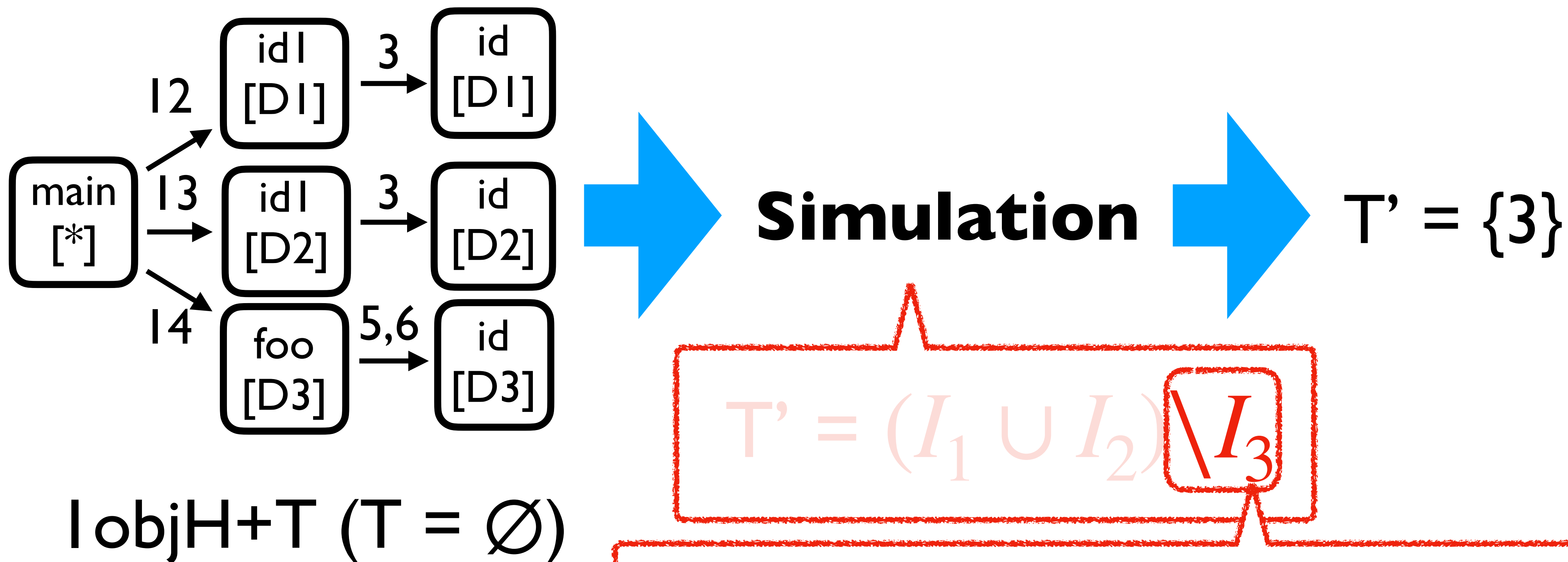
Technique 1: Simulation

- **Simulation** takes a call-graph and produce a tunneling abstraction for CFA



Technique 1: Simulation

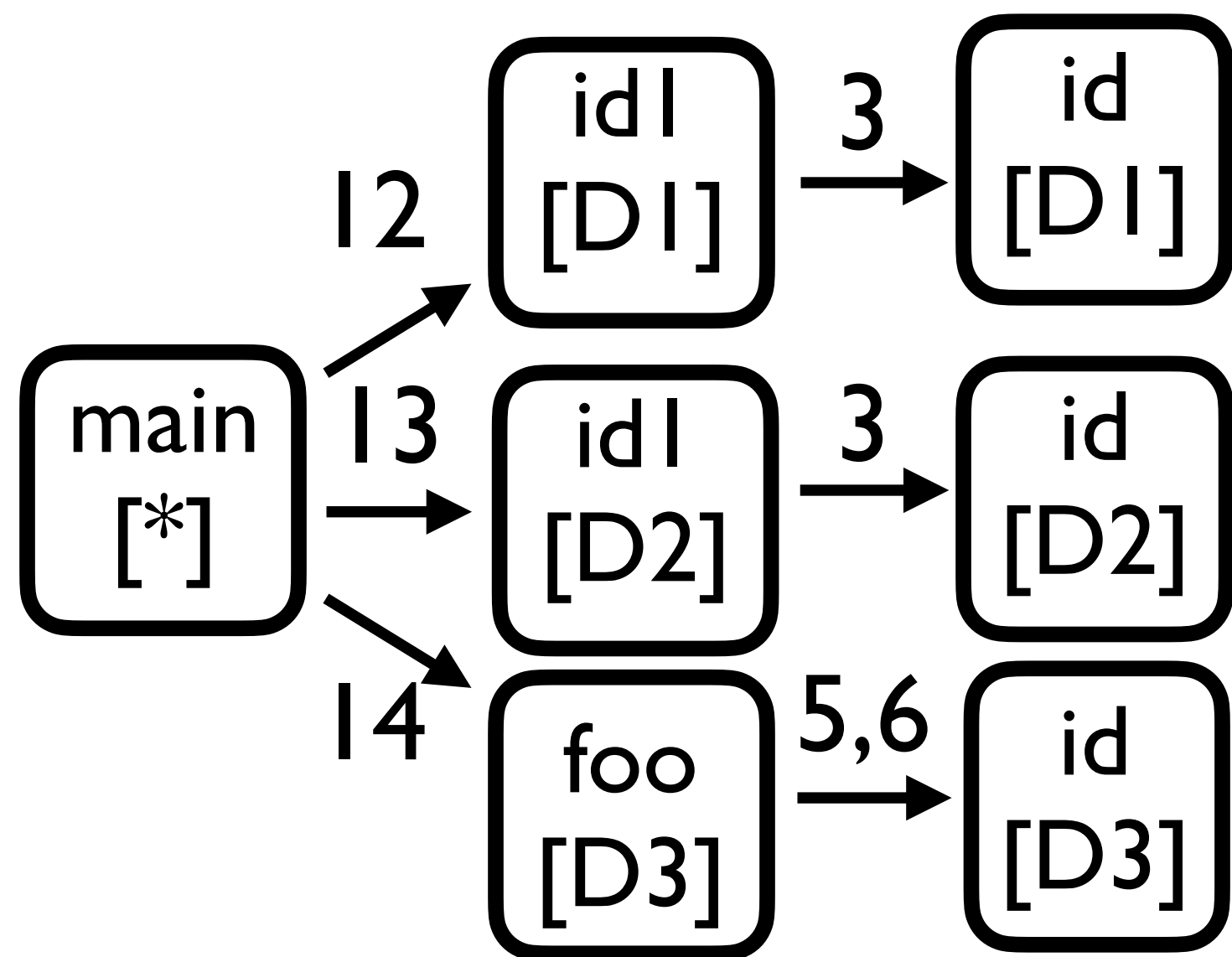
- **Simulation** takes a call-graph and produce a tunneling abstraction for CFA



Tunneling should be avoided for improving precision

Technique 1: Simulation

- **Simulation** takes a call-graph and produce a tunneling abstraction for CFA



Intuition of Simulation

Suppose the call-graph is produced from 1-CFA + T' and infer the T'

~~IobjH+T ($T = \emptyset$)~~

$\text{IcallH+T}'$

What is T'?

Intuition Behind Simulation ($I_1 \cup I_2$)

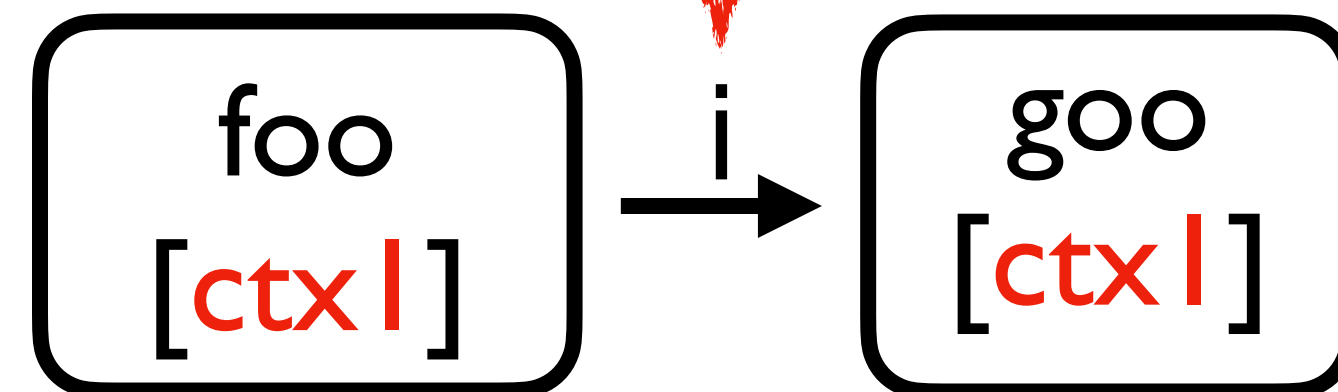
- If tunneling is applied to i , two properties inevitably appear at i

We track the two properties to find the T'

Intuition Behind Simulation ($I_1 \cup I_2$)

- If tunneling is applied to i , two properties inevitably appear at i

Tunneling is applied



Property of context tunneled call-sites

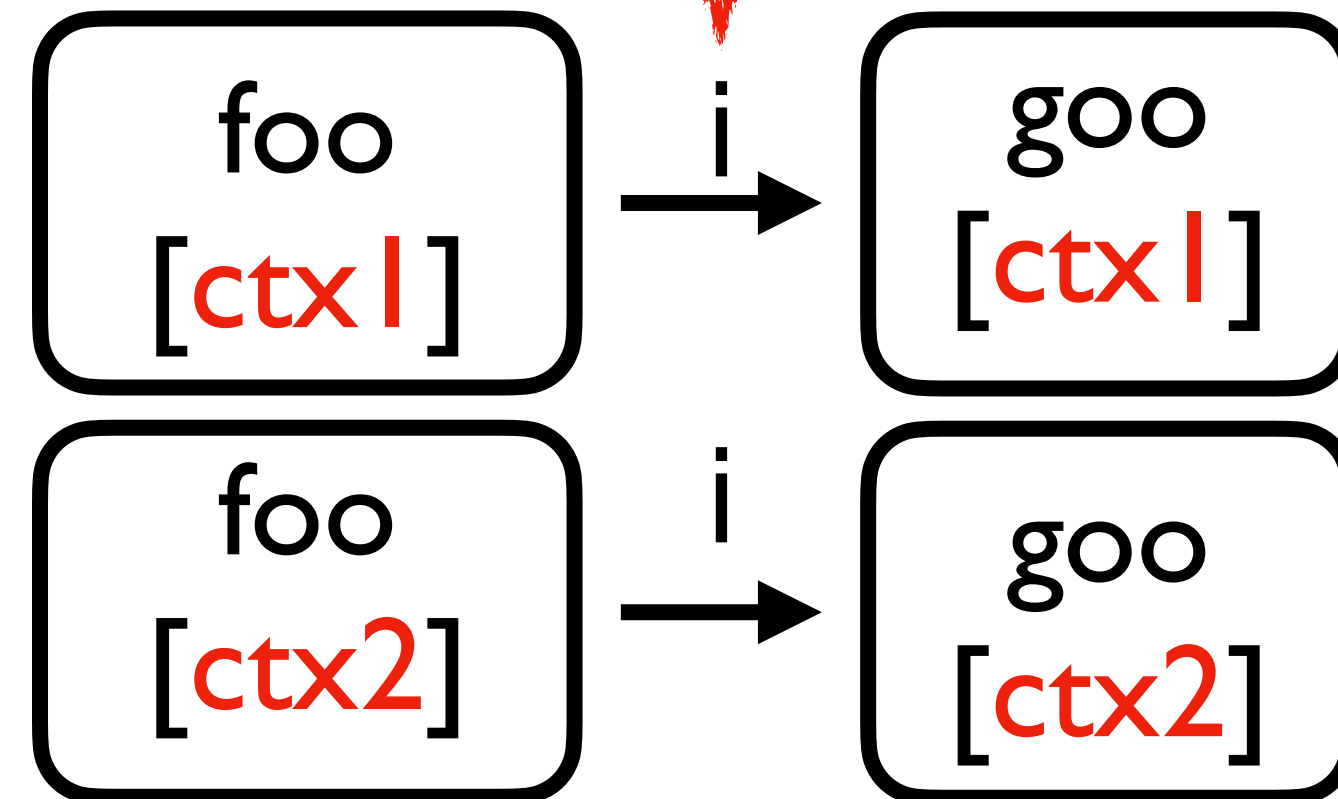
I_1

- Property 1: caller and callee methods have the **same context**

Intuition Behind Simulation ($I_1 \cup I_2$)

- If tunneling is applied to i , two properties inevitably appear at i

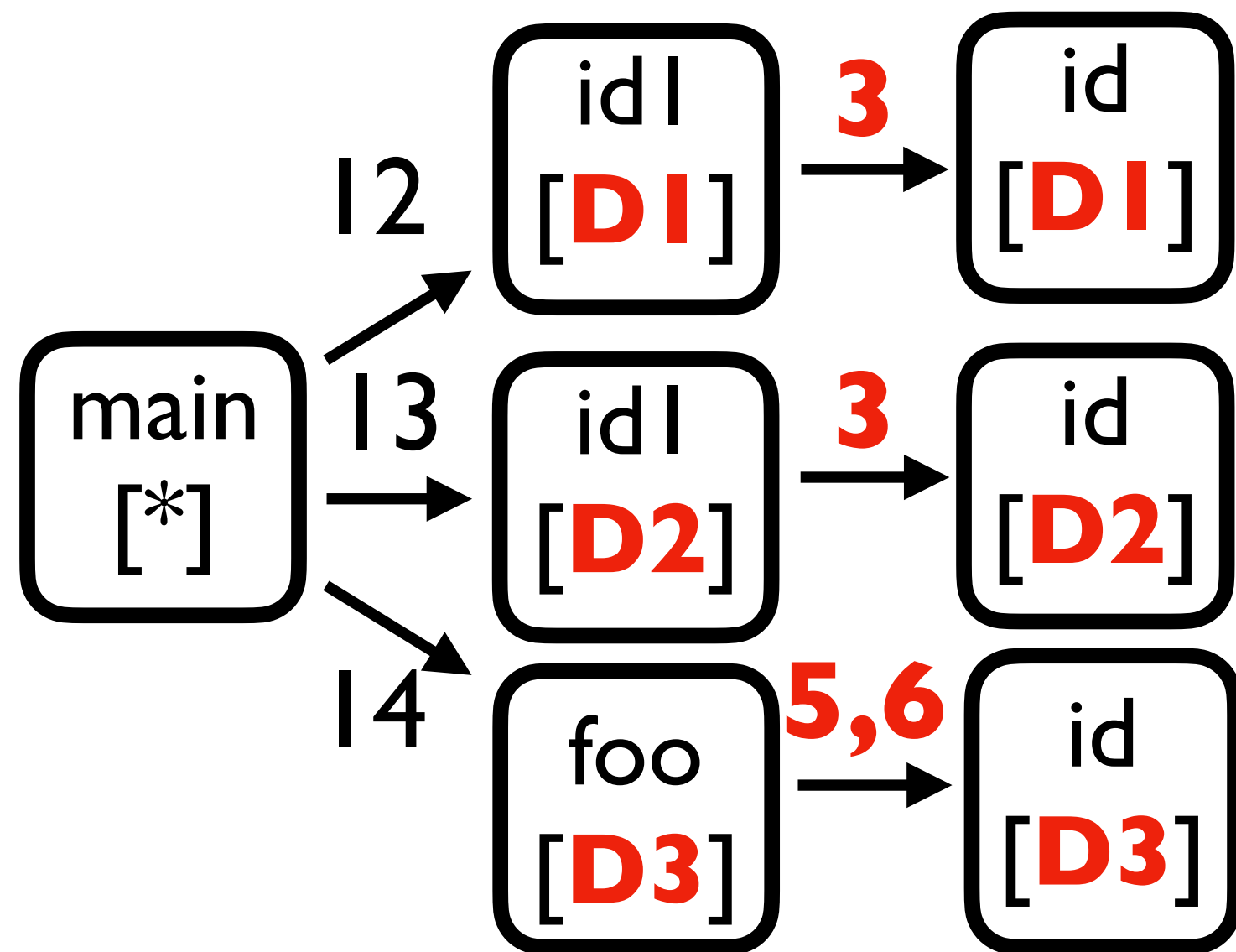
Tunneling is applied



- **Property of context tunneled invocations**
- **Property 2: different caller contexts imply different callee contexts**

Intuition Behind Simulation ($I_1 \cup I_2$)

- Suppose given call-graph is produced from $I \text{ callH}+T'$ and infer what T' is



- I_1 : caller and callee methods have the **same context**

$$I_1 = \{3, 5, 6\}$$

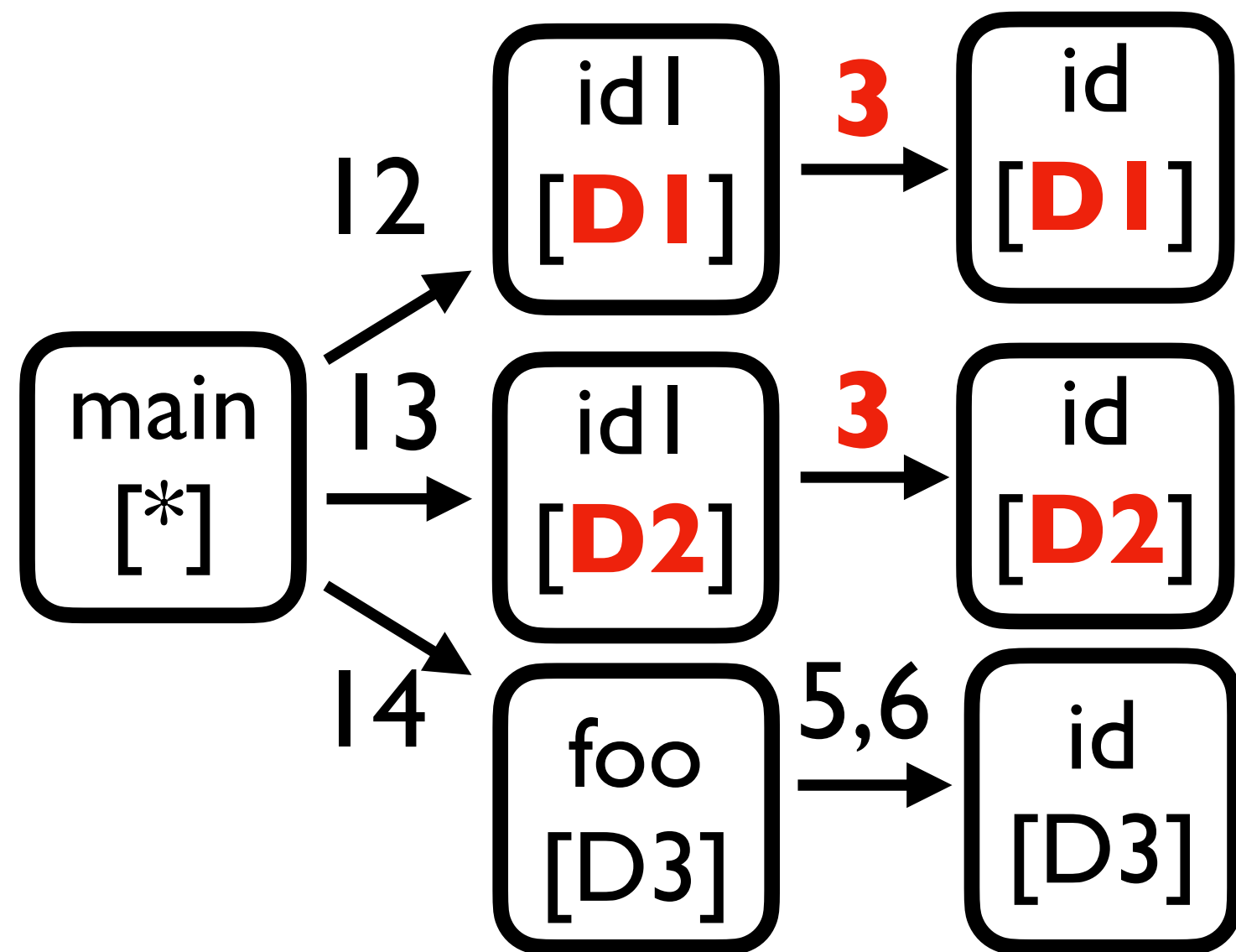
~~$I \text{ objH}+T$ ($T = \emptyset$)~~

$I \text{ callH}+T'$

What is T' ?

Intuition Behind Simulation ($I_1 \cup I_2$)

- Suppose given call-graph is produced from $I_{callH+T}$ and infer what T is



- I_1 : caller and callee methods have the **same context**

$$I_1 = \{3, 5, 6\}$$

- I_2 : different caller ctx imply different callee ctx

$$I_2 = \{3\}$$

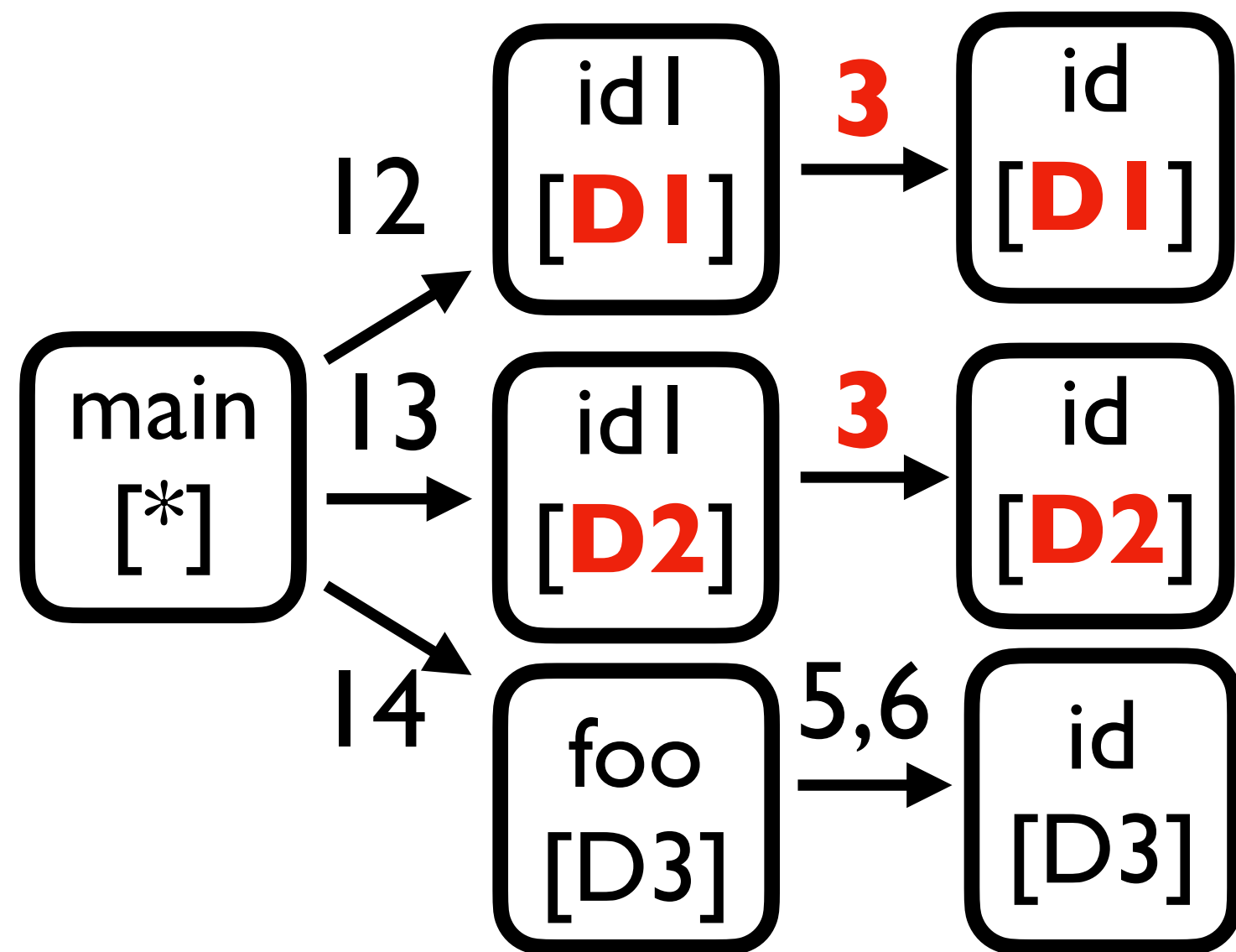
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- I_1 : caller and callee methods have the **same context**

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$$I_2 = \{3\}$$

$$T' = I_1 \cup I_2 = \{3, 5, 6\}$$

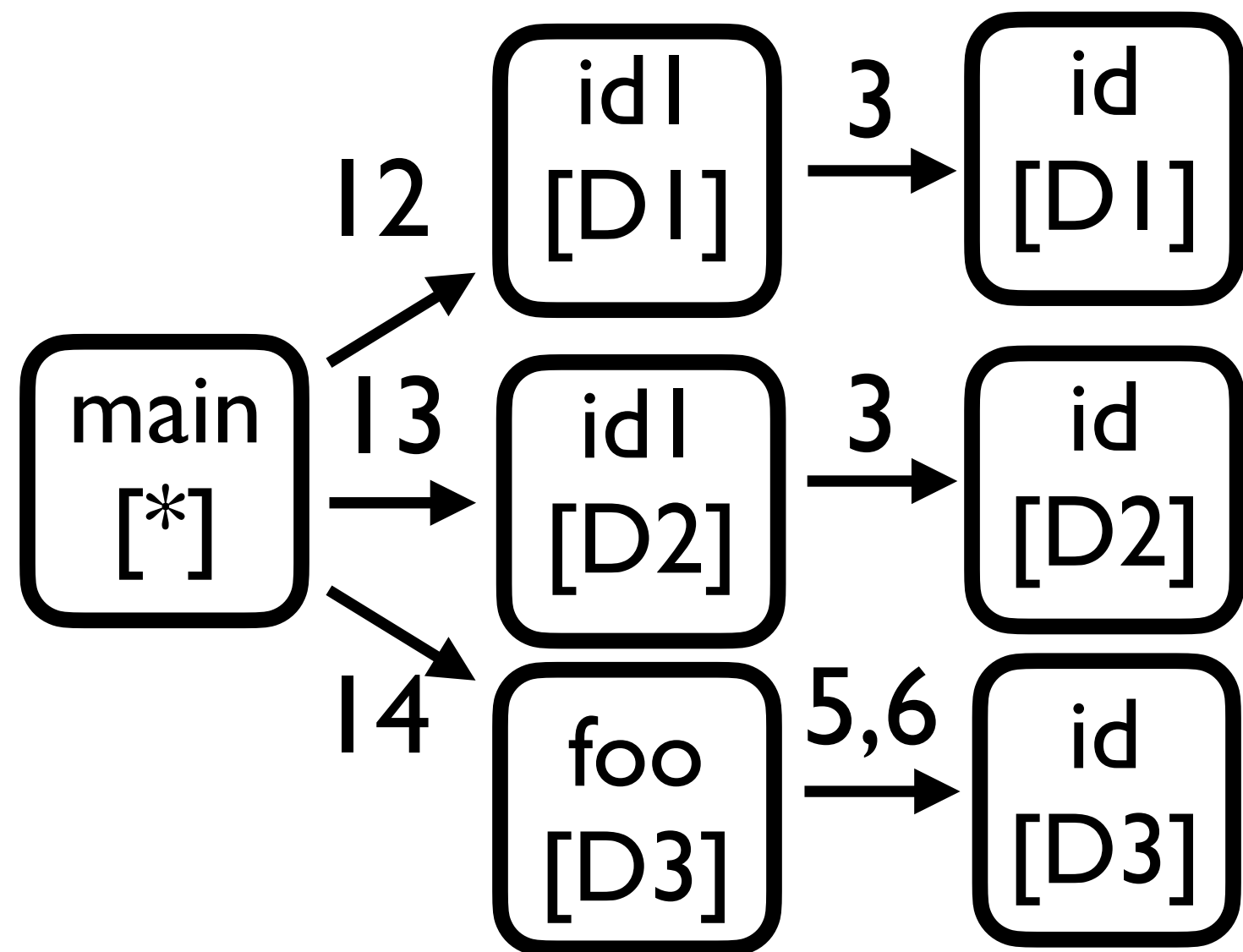
~~I_{objH+T} ($T = \emptyset$)~~

$I_{callH+T'}$

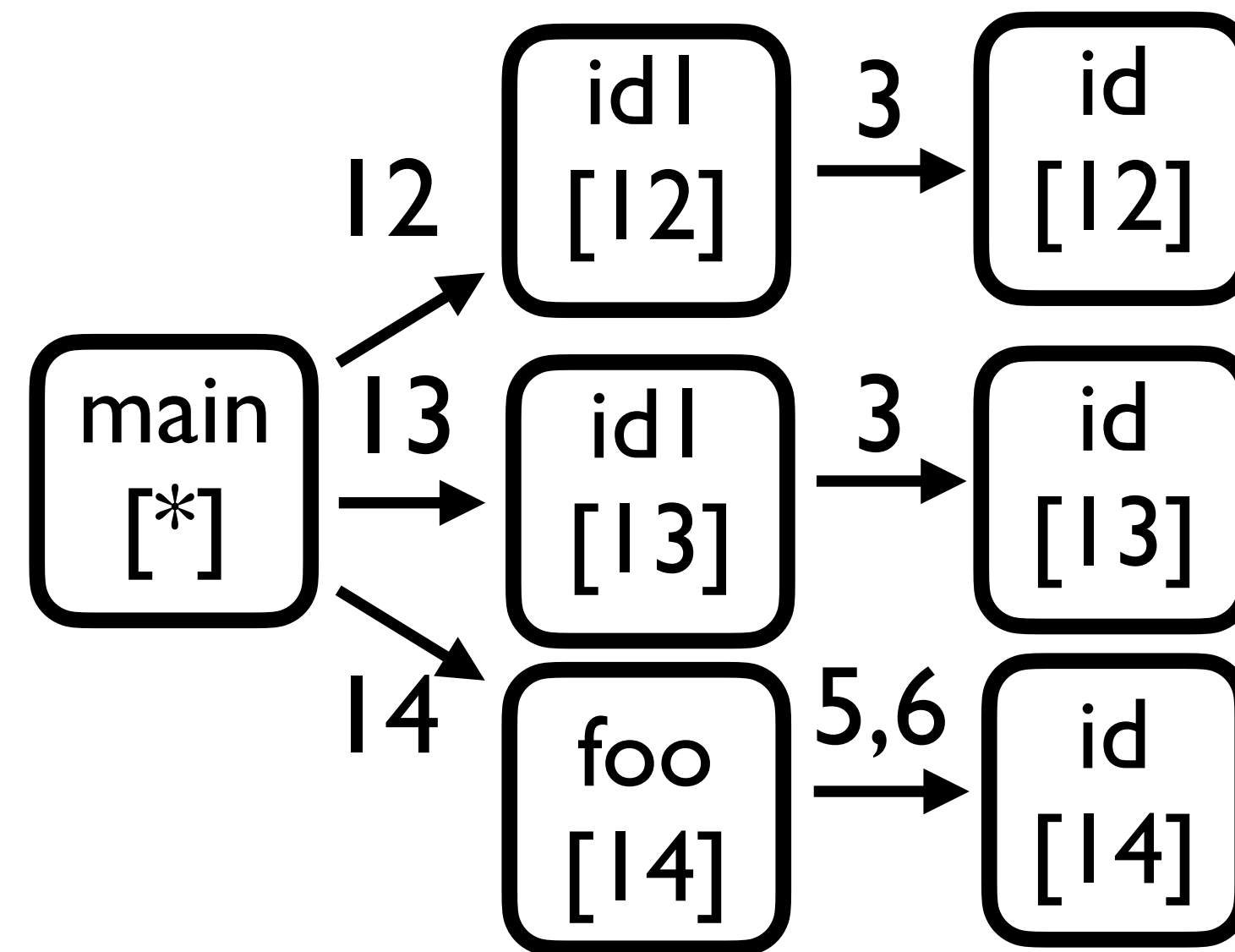
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Intuition Behind Simulation ($I_1 \cup I_2$)

- Suppose given call-graph is produced from $I \text{ callH}+T'$ and infer what T' is



$I \text{ objH}+T$ ($T = \emptyset$)

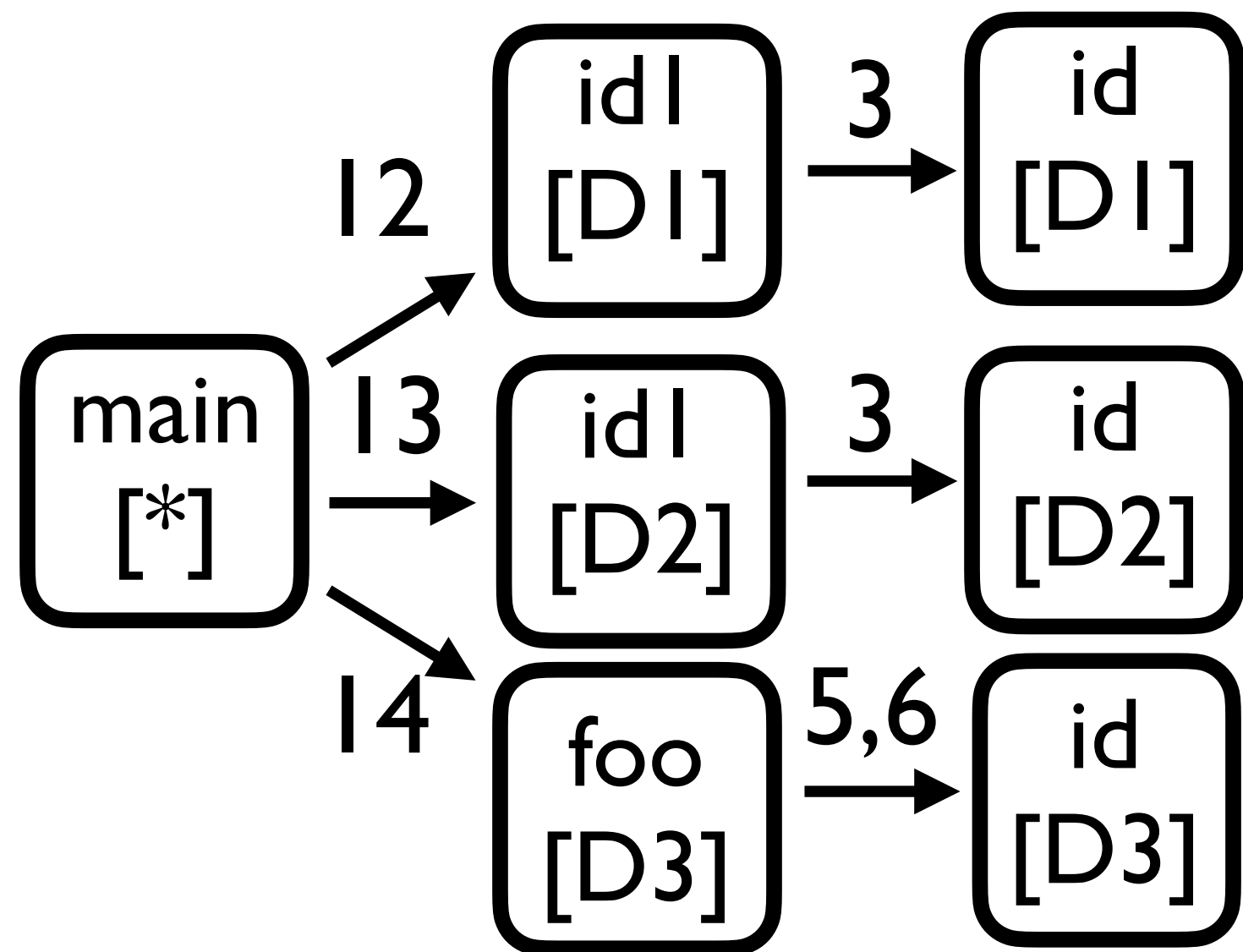


$I \text{ callH}+T'$ ($T' = \{3,5,6\}$)

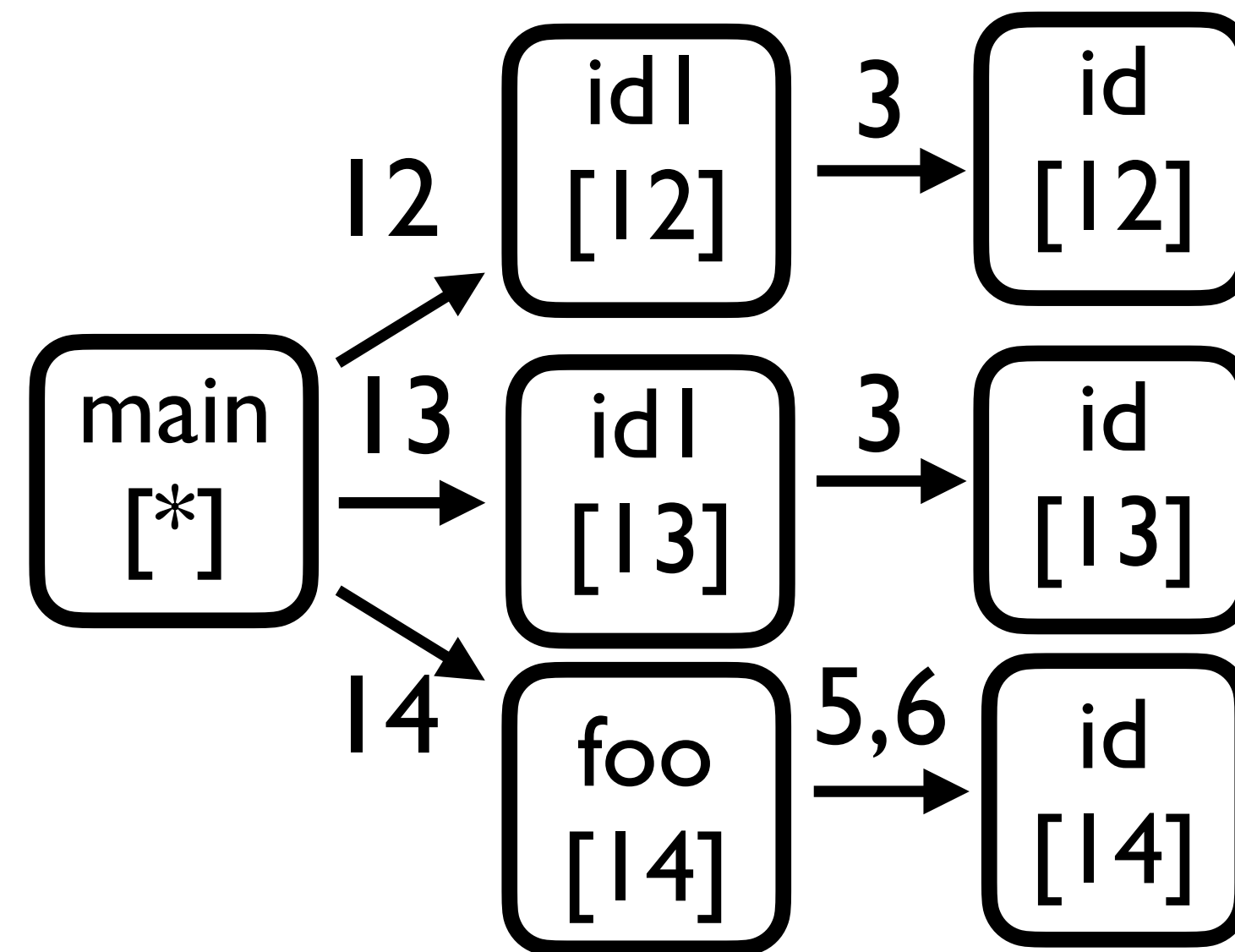
Intuition Behind Simulation ($I_1 \cup I_2$)

- Suppose given call-graph and infer what T' is

Exactly the same analyses



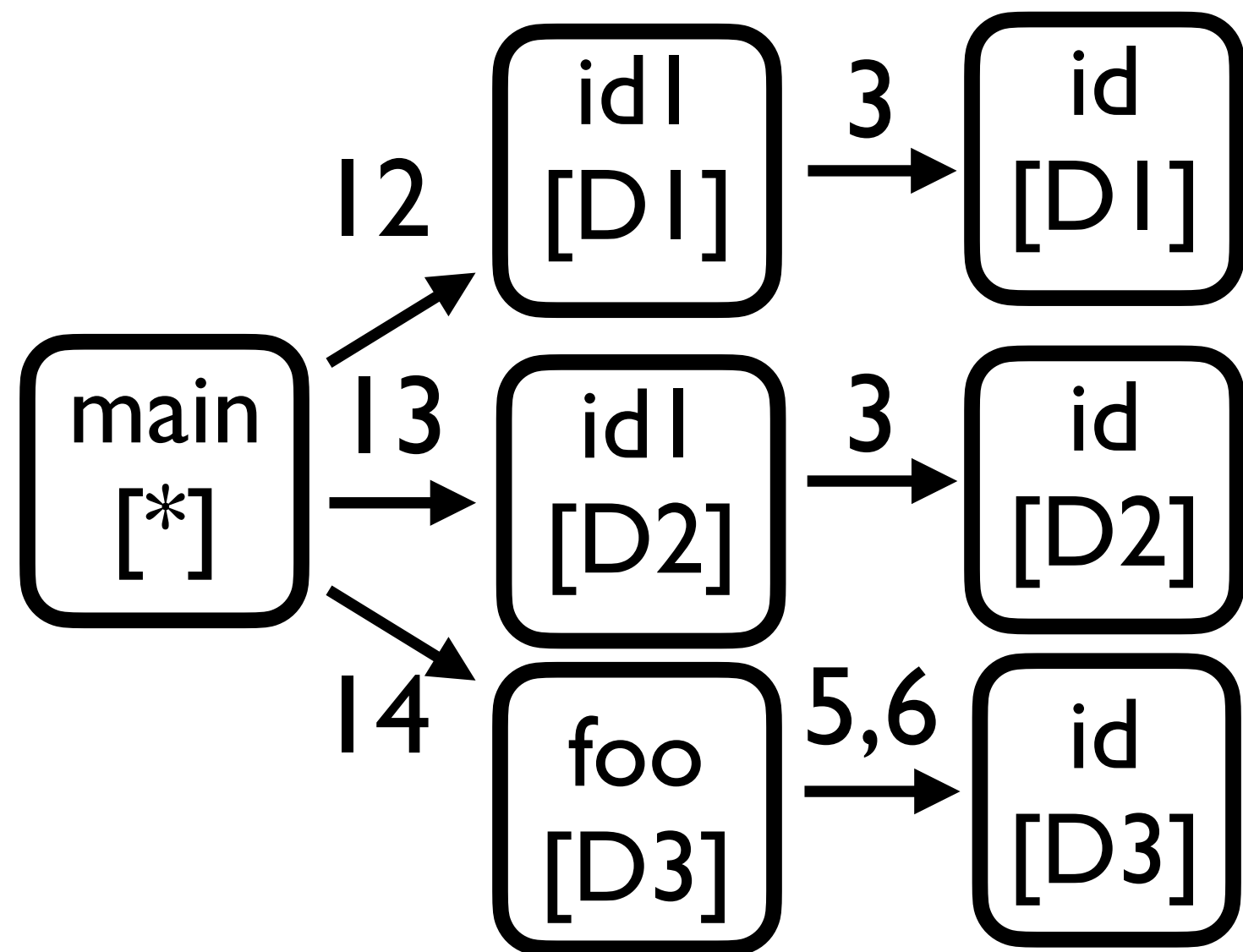
I_{objH+T} ($T = \emptyset$)



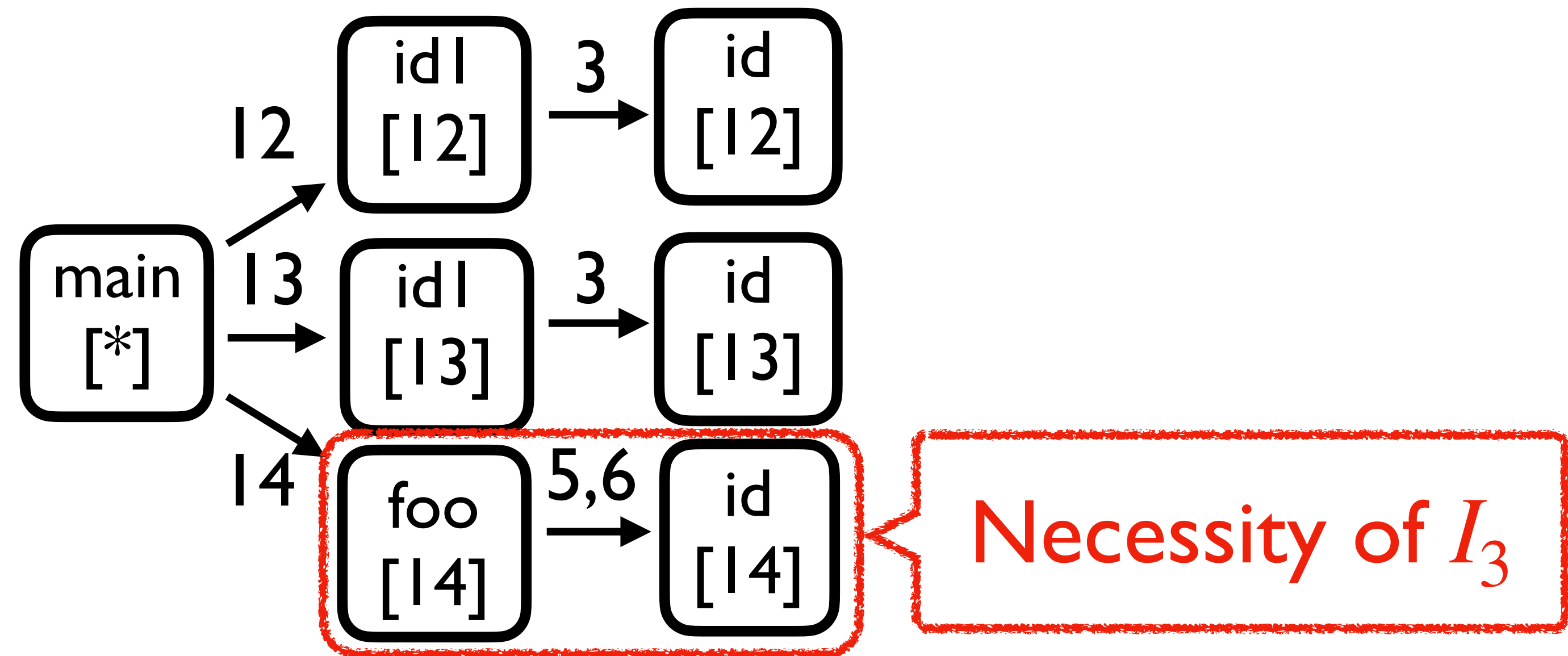
$I_{callH+T'}$ ($T' = \{3, 5, 6\}$)

Intuition Behind Simulation ($I_1 \cup I_2$)

- Suppose given call-graph is produced from $I_{callH+T}$ and infer what T is



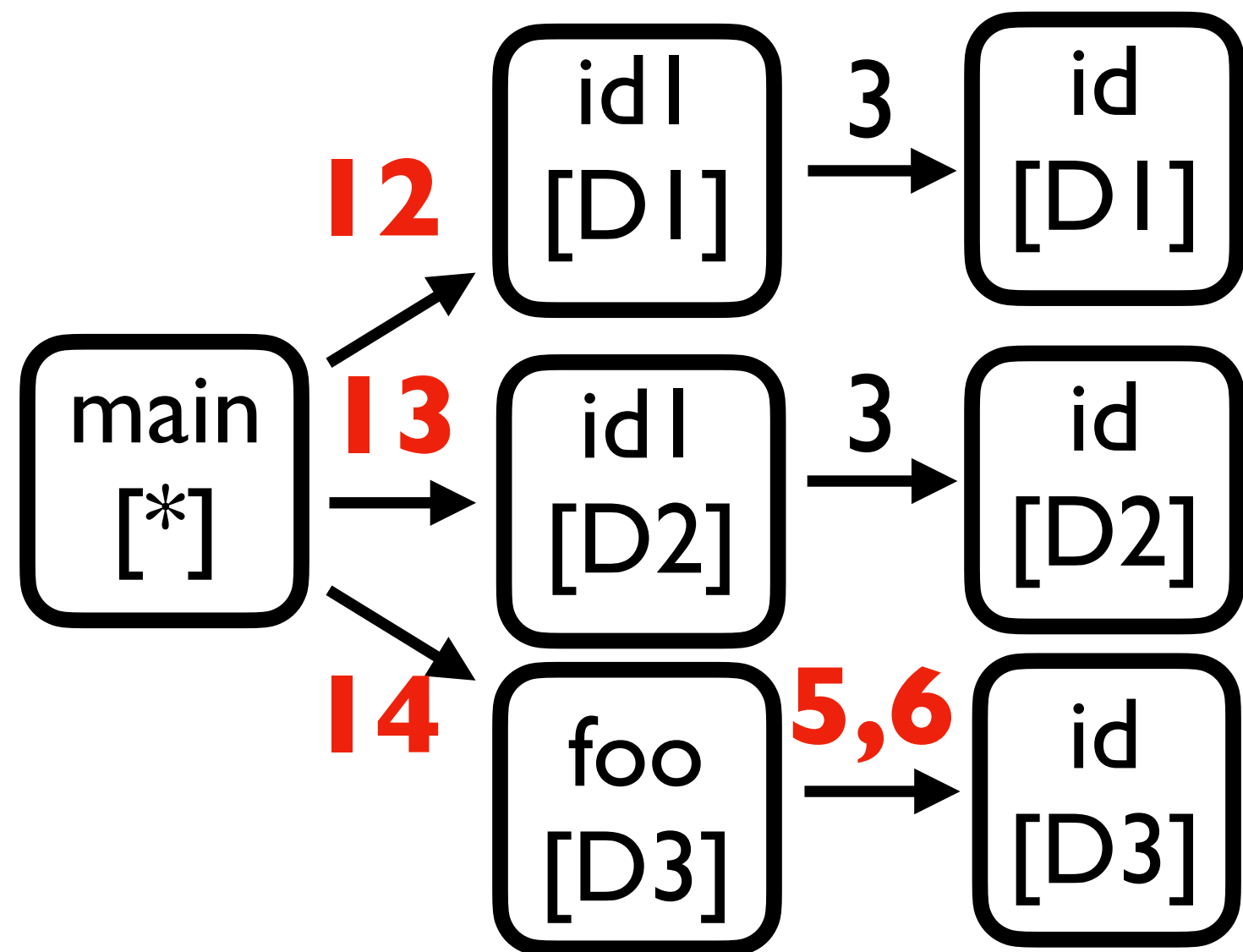
I_{objH+T} ($T = \emptyset$)



$I_{callH+T'}$ ($T' = \{3,5,6\}$)

Intuition Behind Simulation (I_3)

- I_3 : Tunneling should be avoided for improving precision



- I_1 : caller and callee methods have the same context

$$I_1 = \{3, 5, 6\}$$

- I_2 : different caller ctx imply different callee ctx

$$I_2 = \{3\}$$

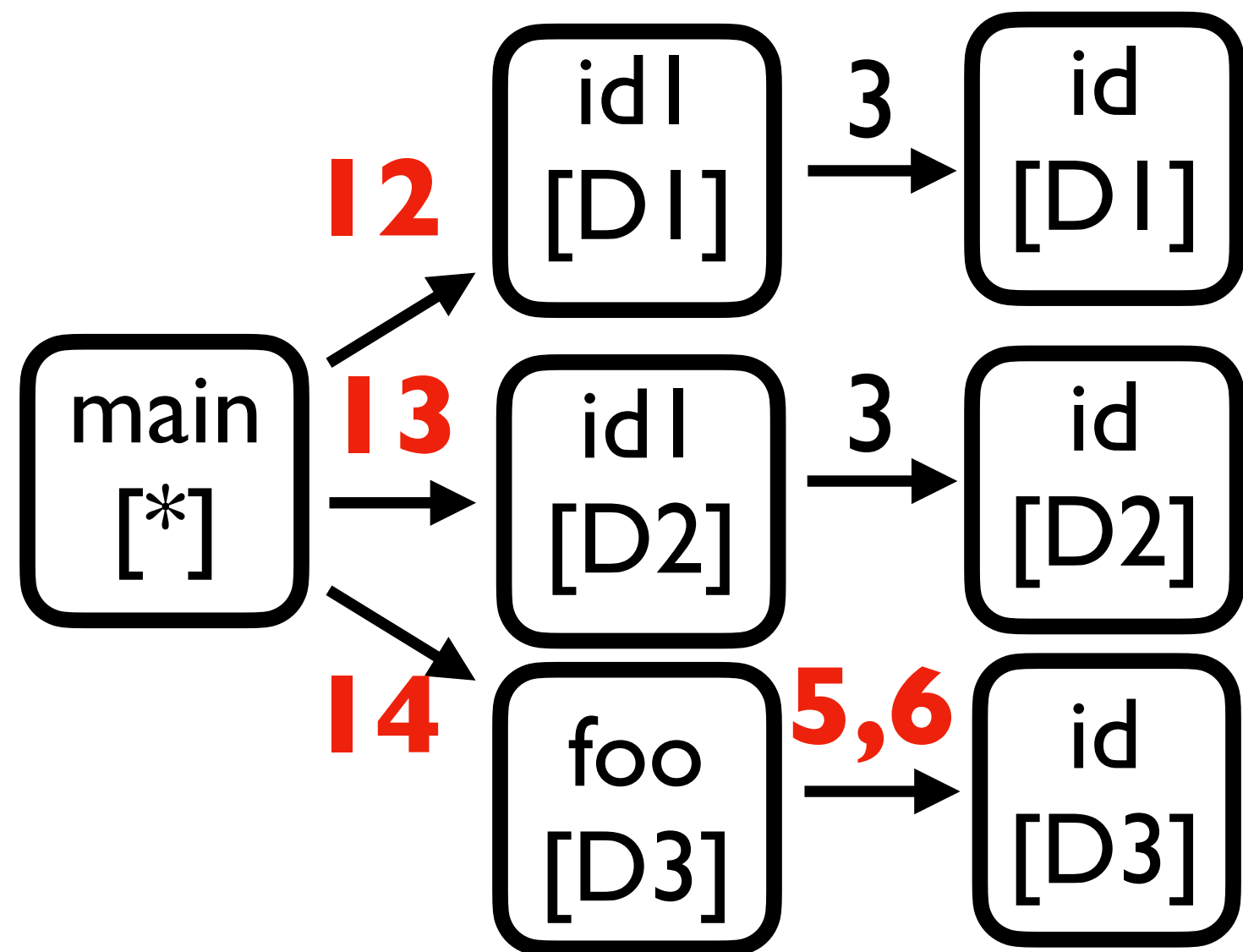
- I_3 : given object sensitivity produced only one context

$$I_{objH+T} (T = \emptyset)$$

$$I_3 = \{5, 6, 12, 13, 14\}$$

Intuition Behind Simulation

- The inferred tunneling abstraction T' is a singleton set $\{3\}$



- I_1 : caller and callee methods have the same context

$$I_1 = \{3, 5, 6\}$$

- I_2 : different caller ctx imply

$$I_2 = \{3\}$$

$$T' = (I_1 \cup I_2) \setminus I_3 = \{3\}$$

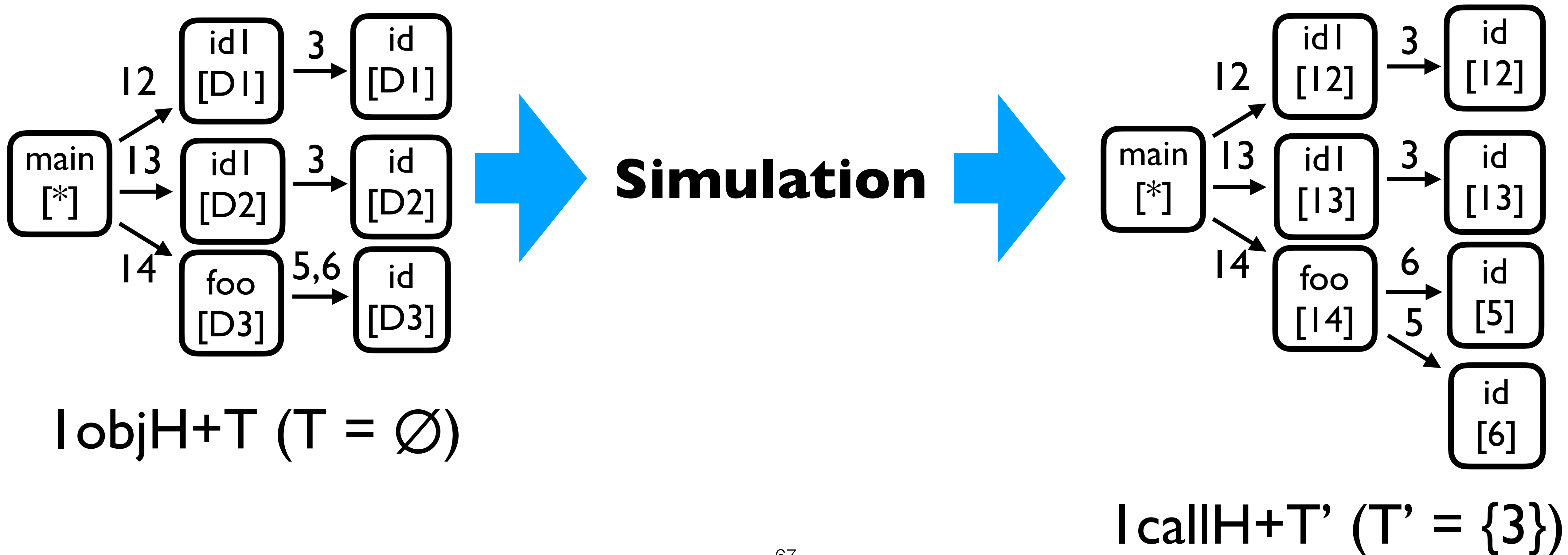
- I_3 : given object sensitivity produced only one context

$$I_3 = \{5, 6, 12, 13, 14\}$$

$\text{I objH+T} \quad (T = \emptyset)$

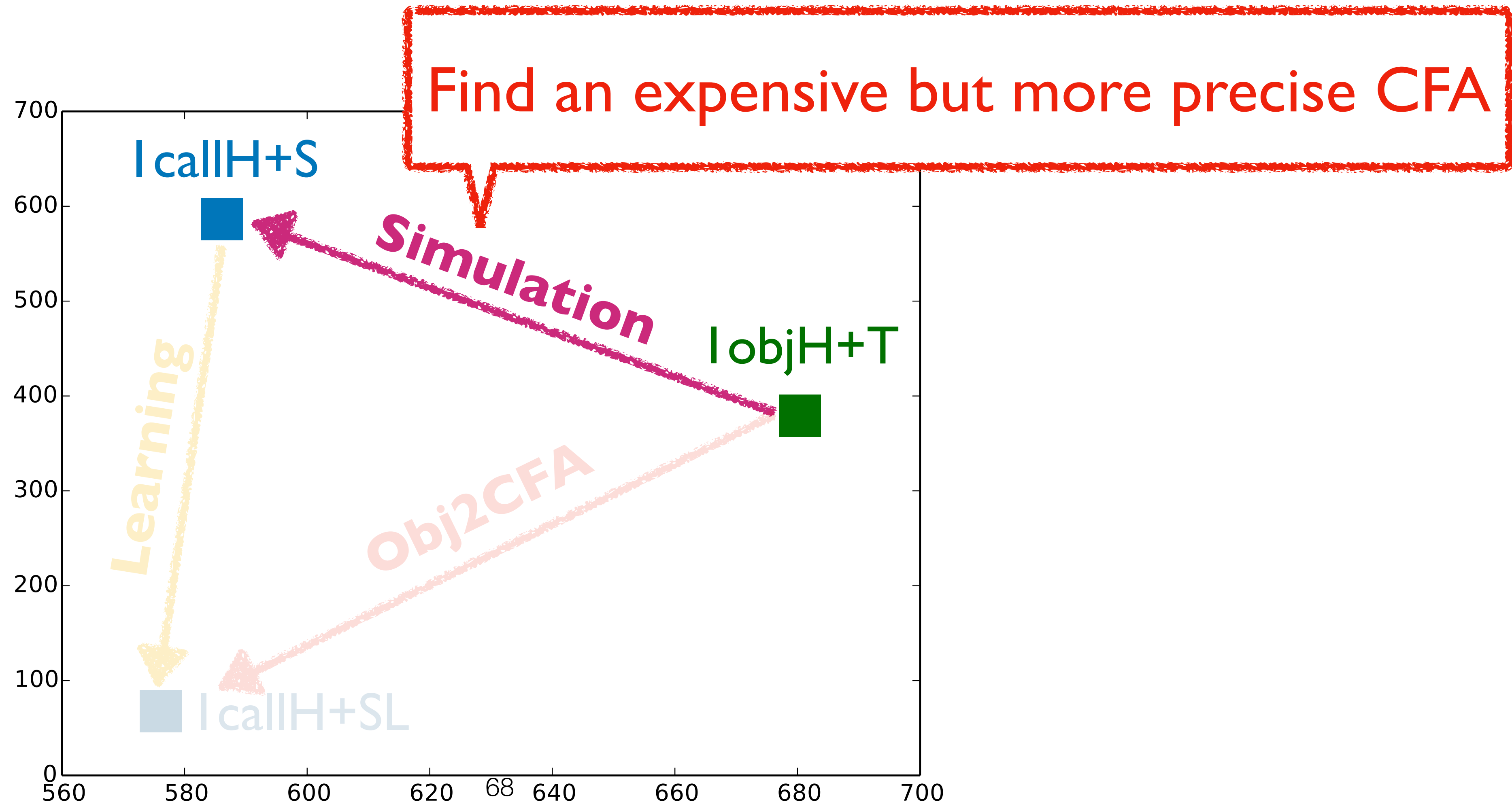
Technique 1: Simulation

- With T' , CFA becomes more precise than the given object sensitivity



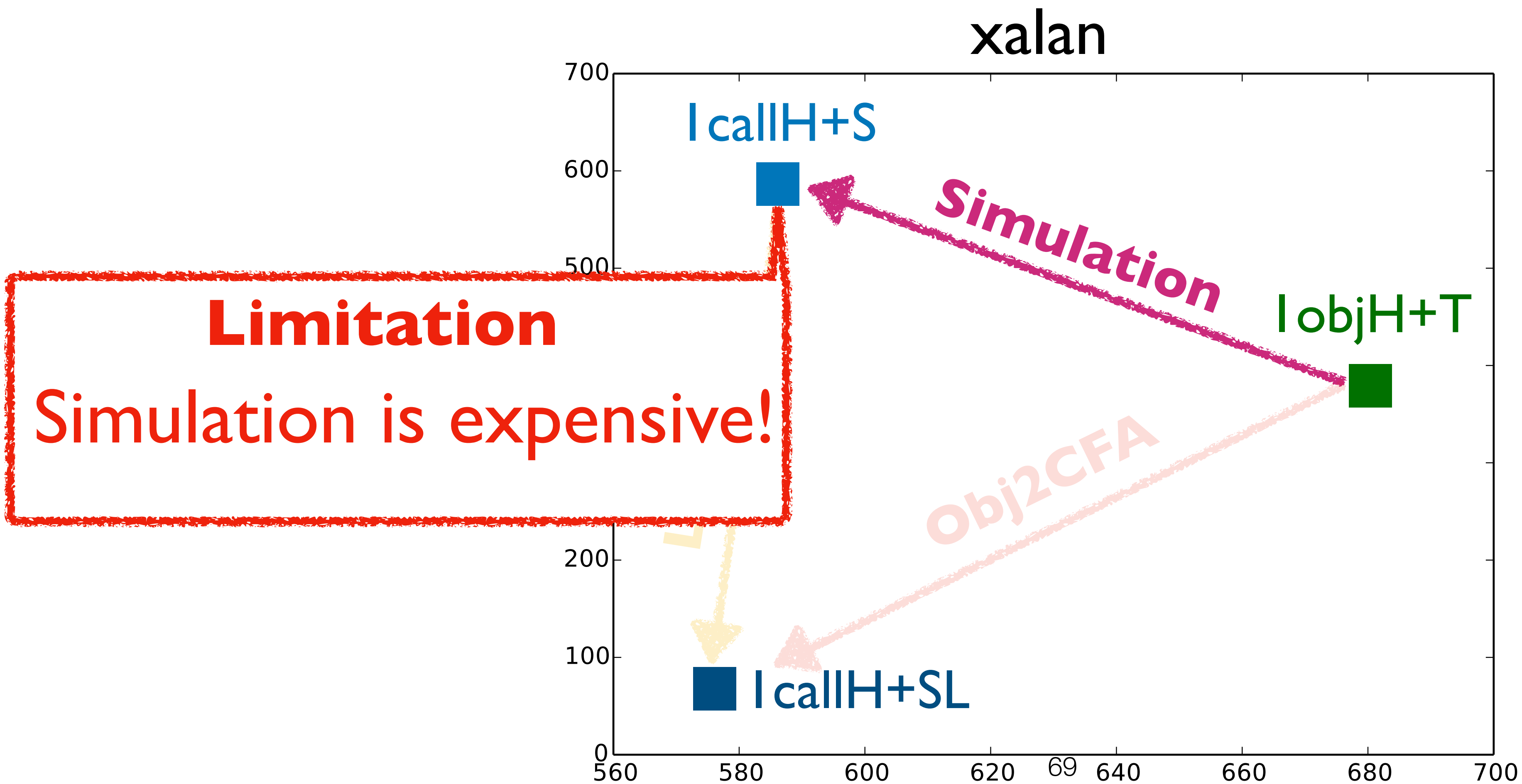
Our Technique : **Obj2CFA**

- **Obj2CFA** consists of **simulation** and simulation-guided **learning**



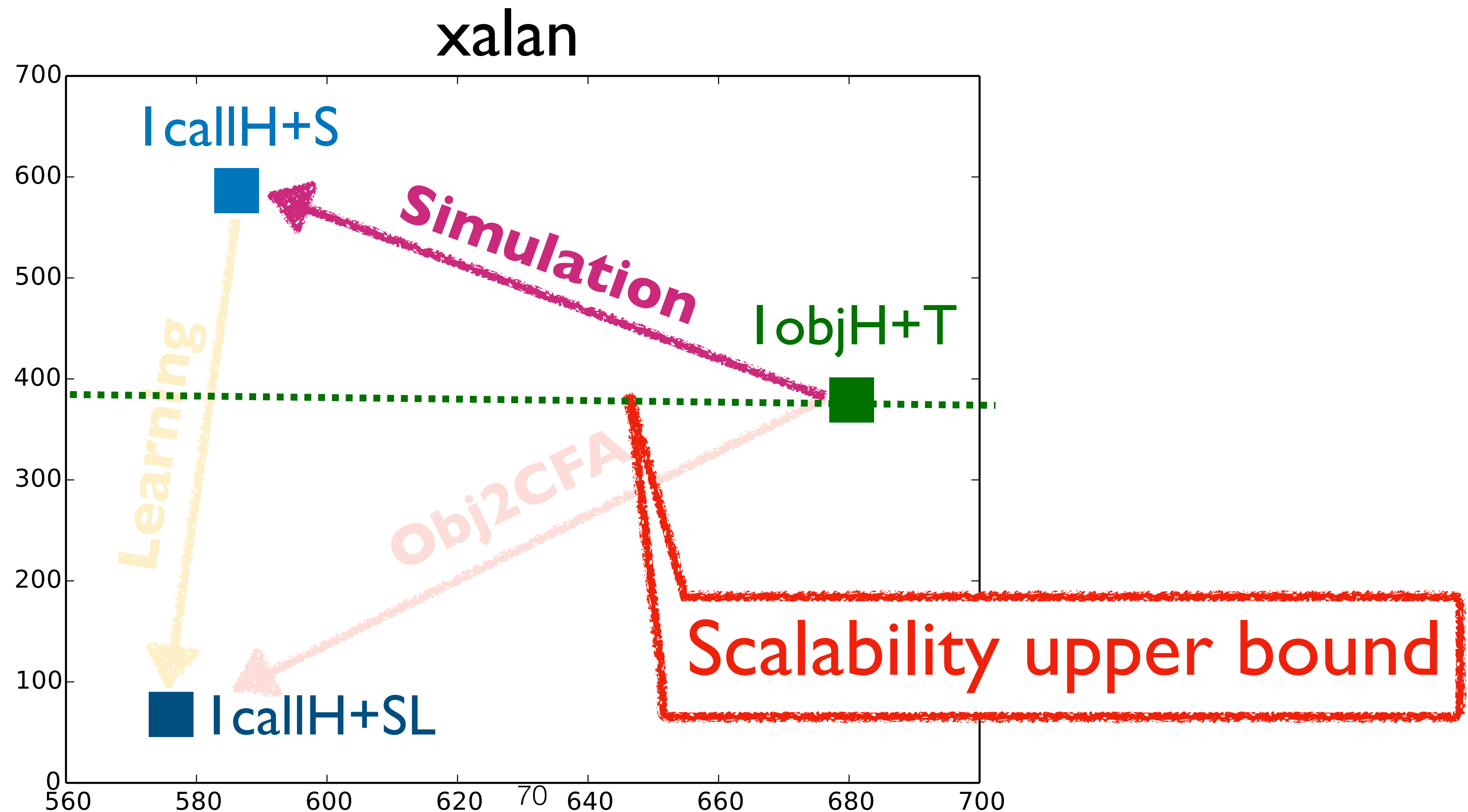
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Our Technique : **Obj2CFA**

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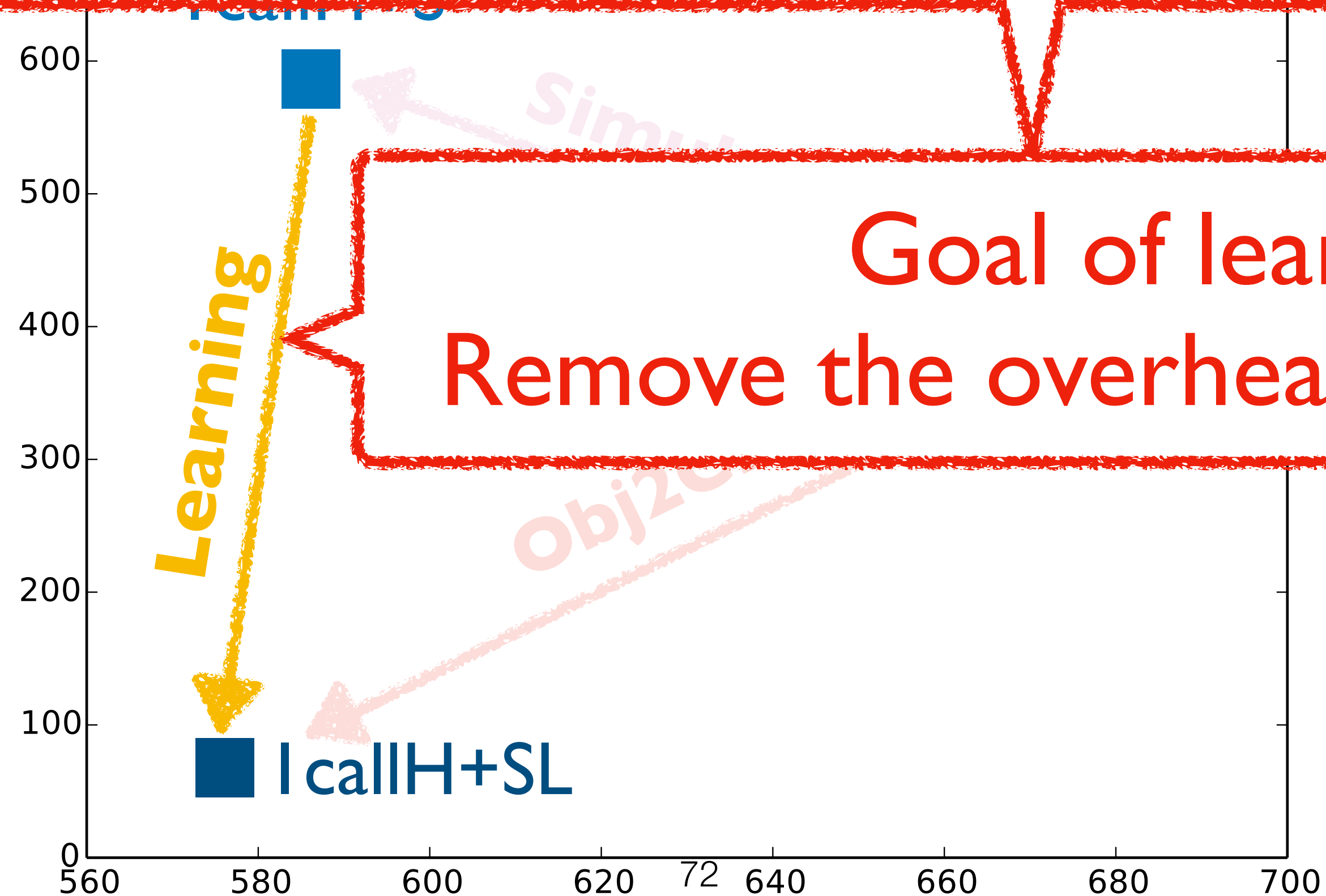
Our Technique : **Obj2CFA**

- **Obj2CFA** consists of **simulation** and simulation-guided **learning**



Our Technique • Obj2CEA

Given training programs and simulated tunneling abstractions, learning aims to find a model that produces similar tunneling abstractions without running the given object sensitivity



Goal of learning:
Remove the overhead of simulation

Our Technique • **Obi2CFA**

Given training programs and simulated tunneling abstractions, learning aims to find a model that produces similar tunneling

The learned model will produce tunneling abstractions without running object sensitivity

Details in paper



Evaluation

Setting

- Doop
- Pointer analysis framework for Java
- Research Question: which one is better?

Call-site sensitivity vs Object sensitivity

Context tunneling is included

Setting

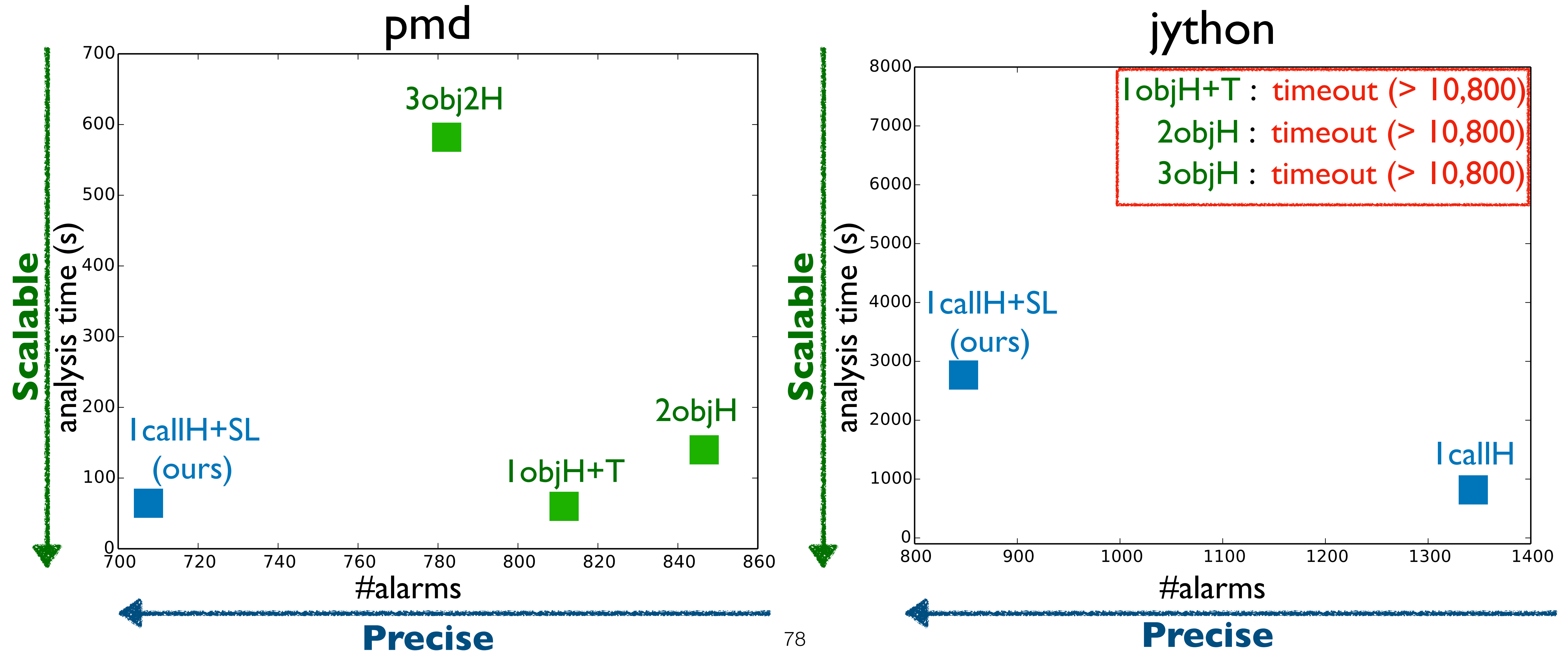
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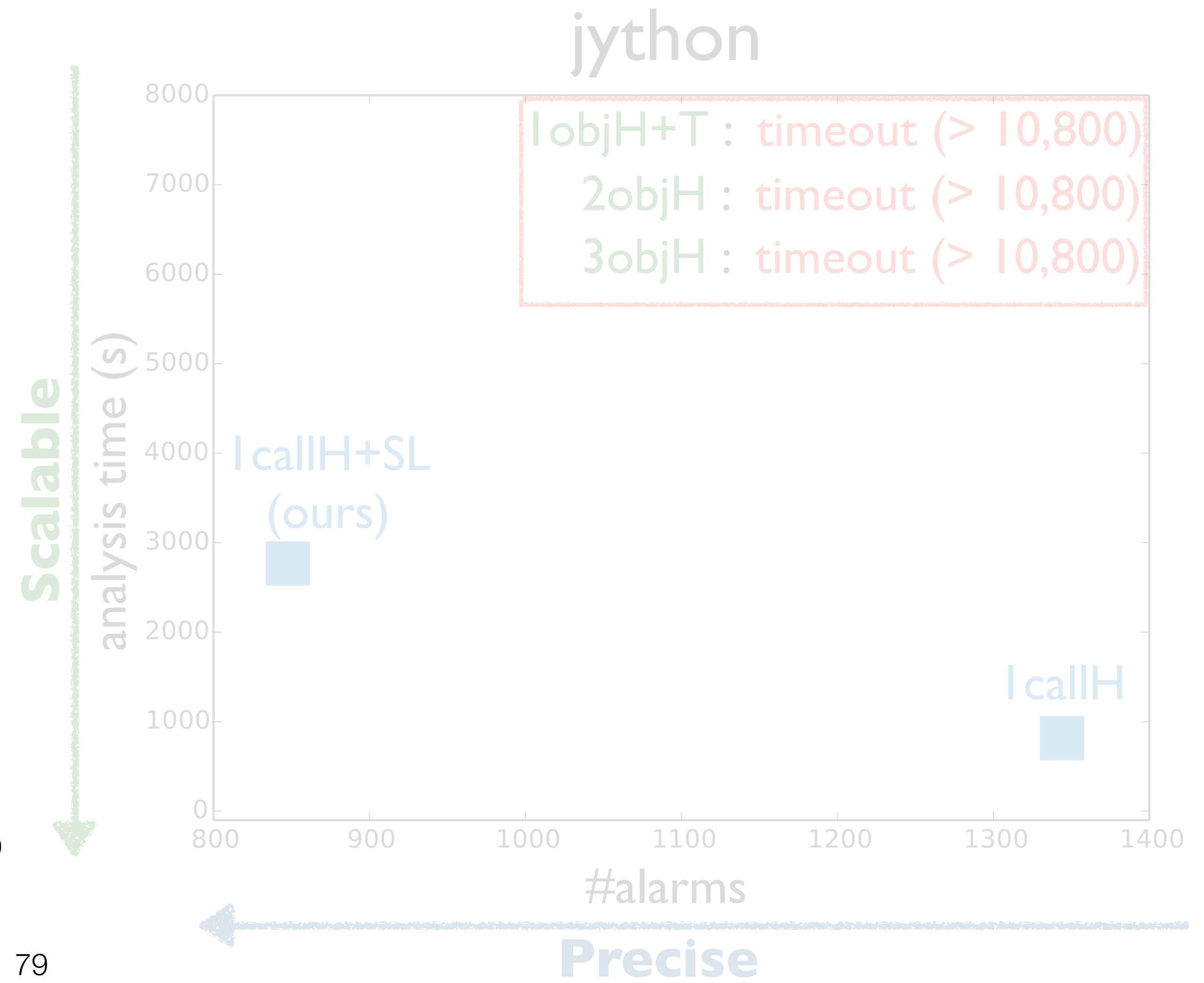
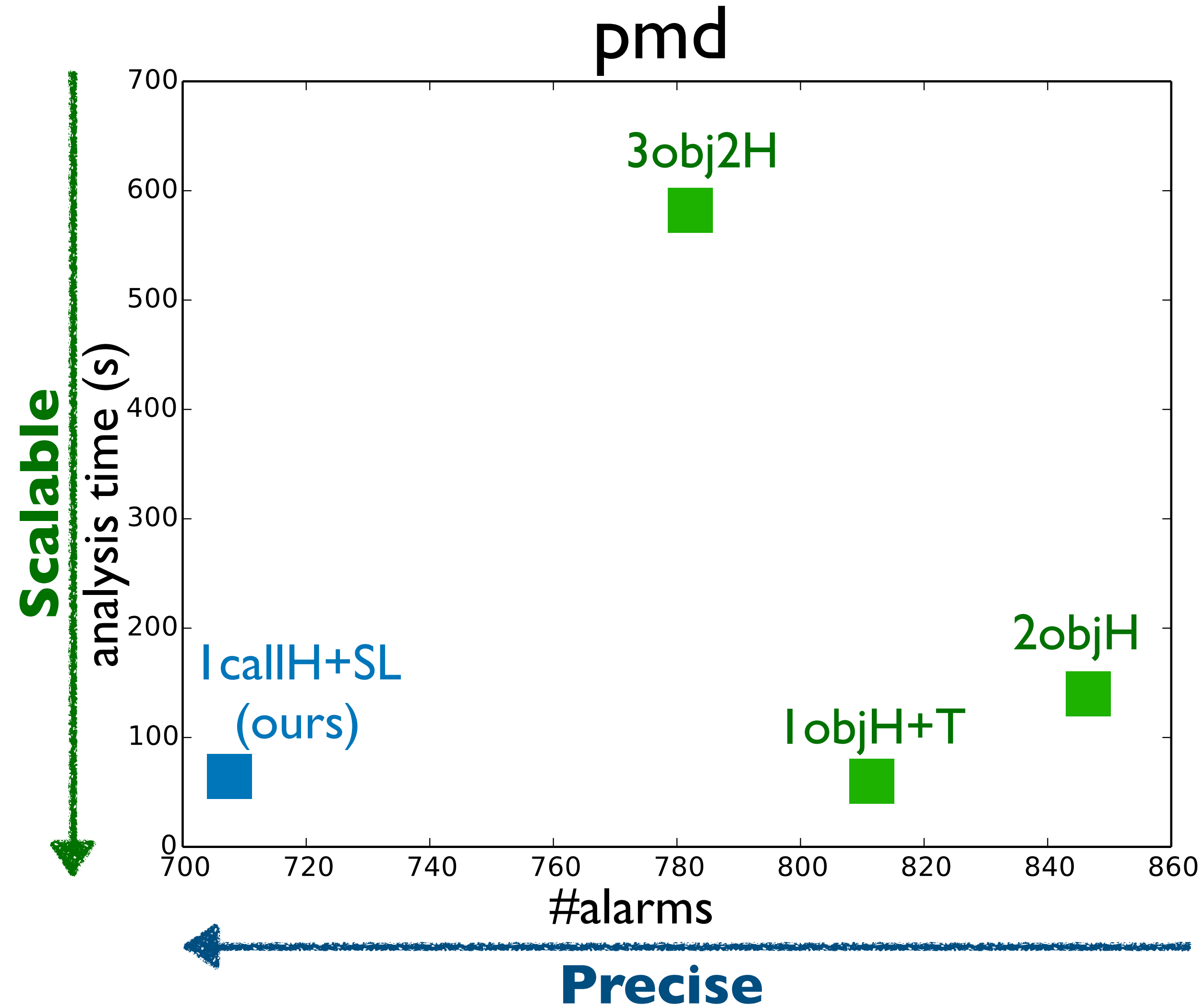
Call-site Sensitivity vs Object Sensitivity

- **lcallH+SL (ours)** is **more precise and scalable** than the **existing object sensitivities**



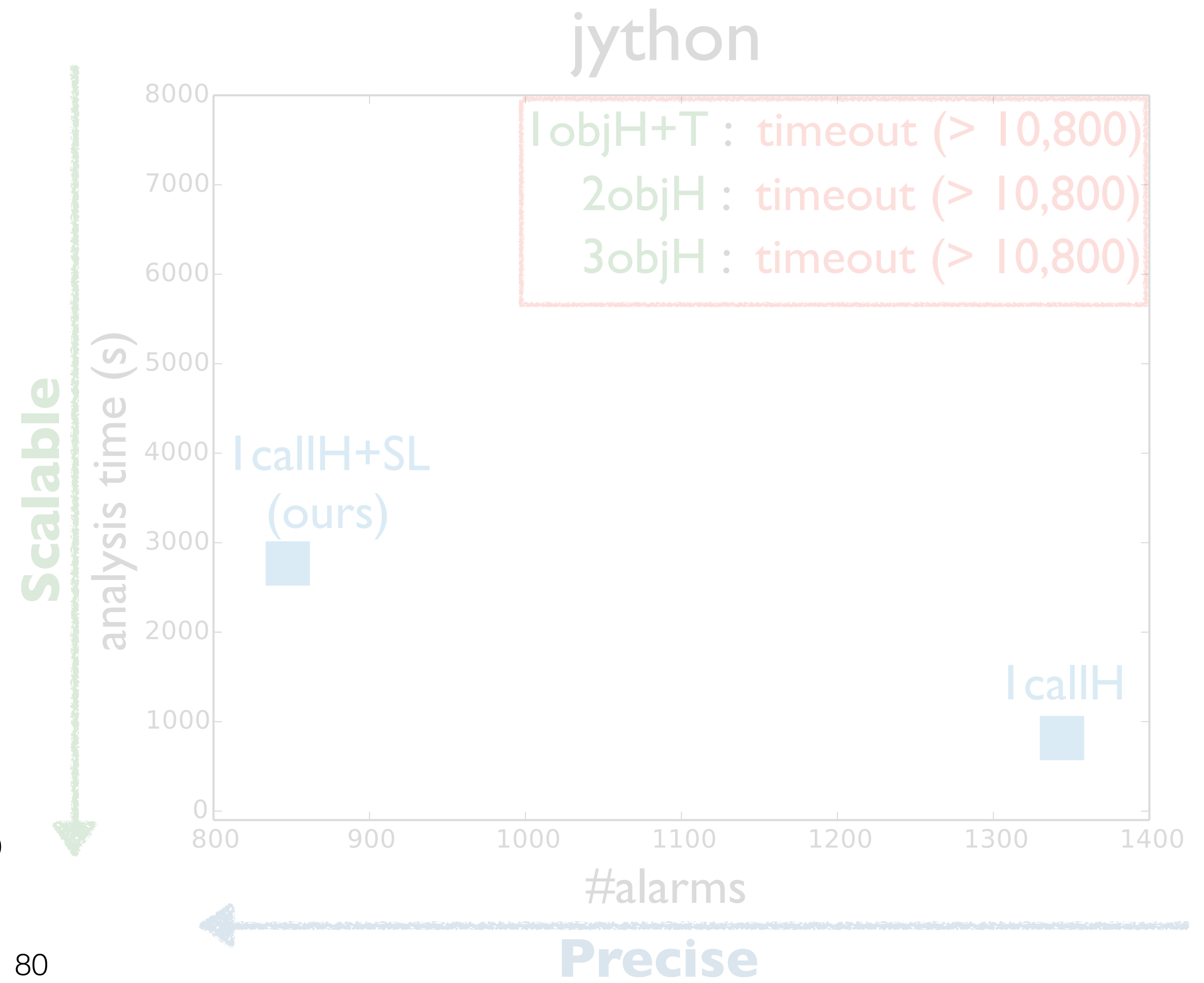
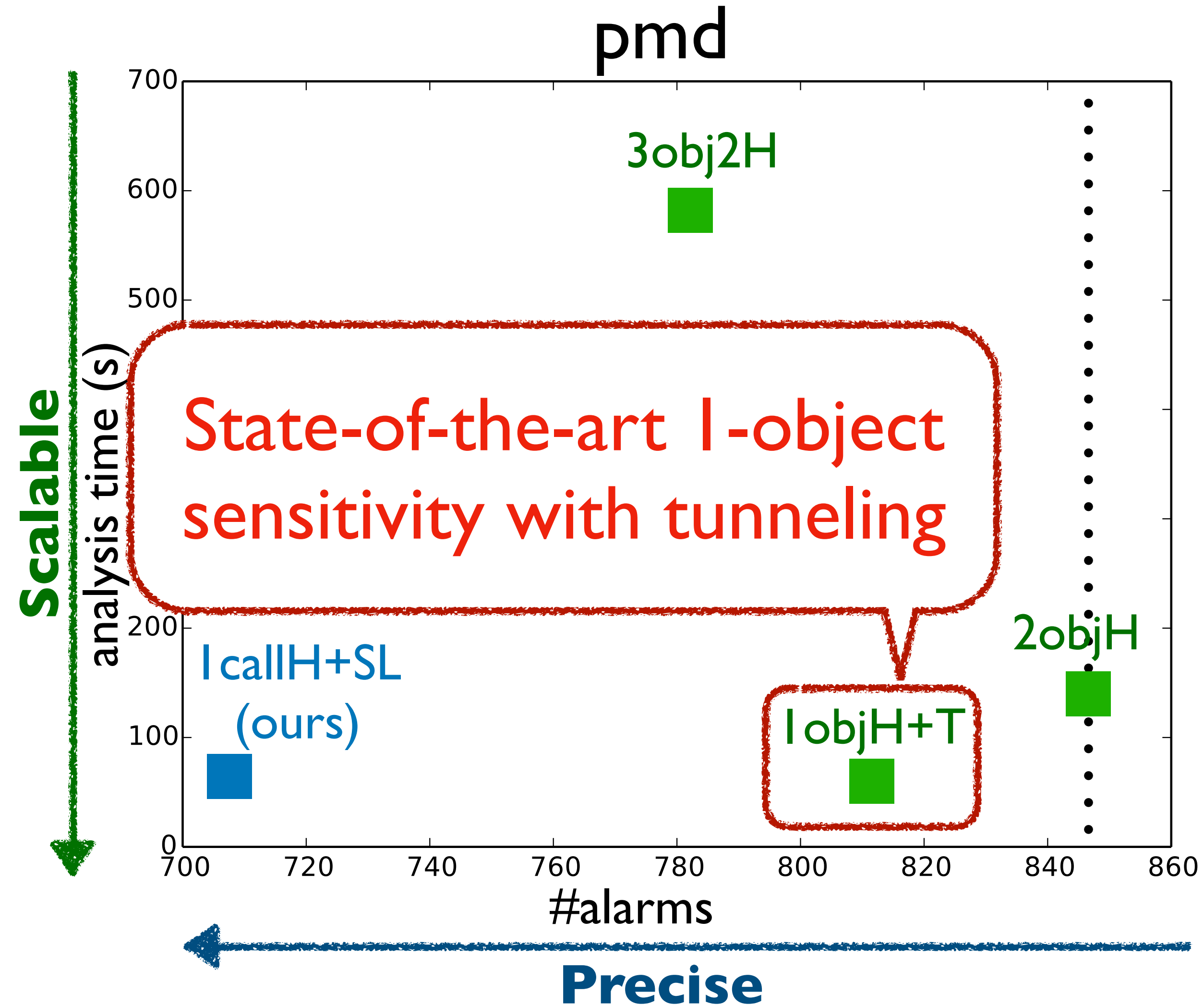
Call-site Sensitivity vs Object Sensitivity

- **1callH+SL (ours)** is **more precise** and **scalable** than the existing object sensitivities



Call-site Sensitivity vs Object Sensitivity

- `lcallH+SL` (ours) is **more precise** and **scalable** than the existing object sensitivities



Call-site Sensitivity vs Object Sensitivity

- `lcallH+SL` (ours) is **more precise** and **scalable** than the existing object sensitivities

pmd



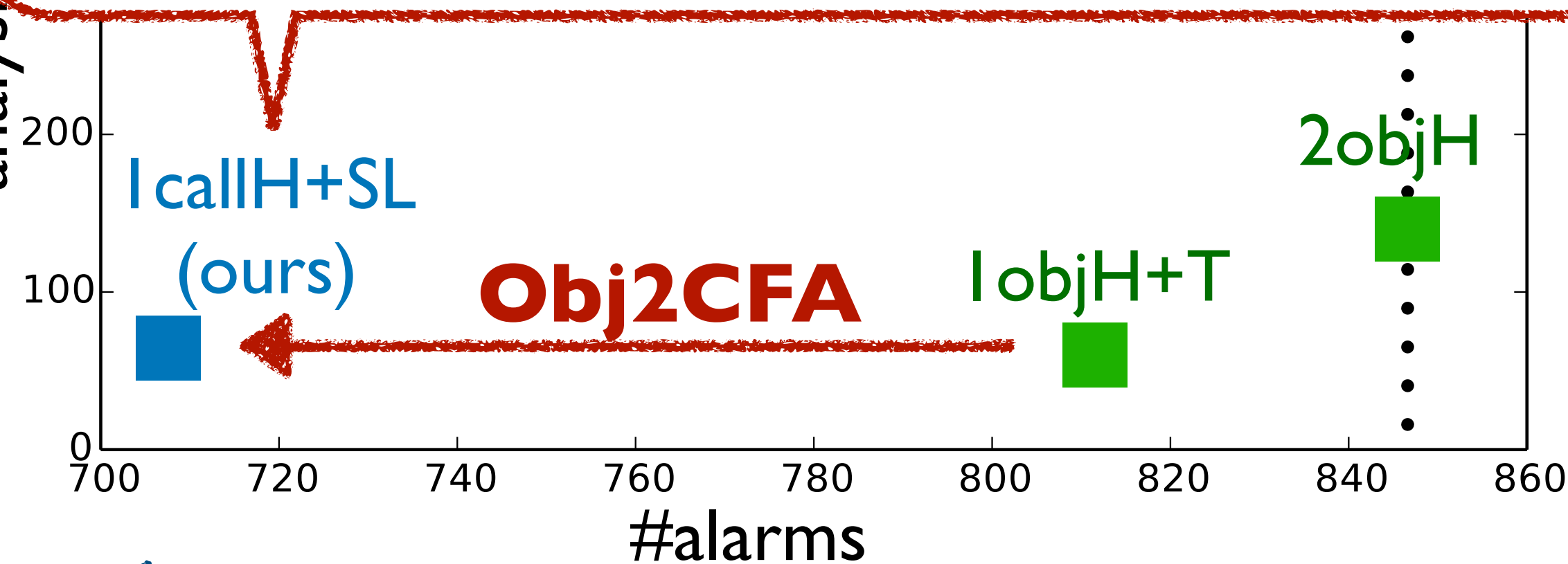
jython



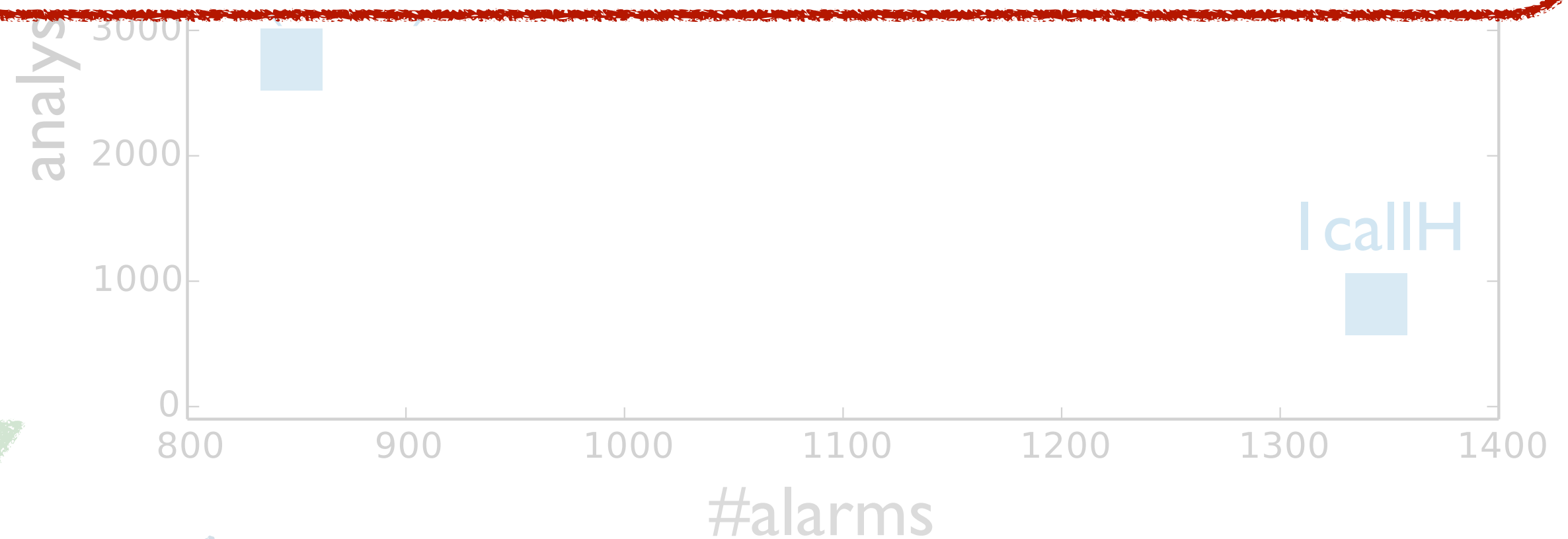
Transformed I-CFA with tunneling via **Obj2CFA** from `lobjH+T`

Scalable

analysis



Scalable



Precise

Precise

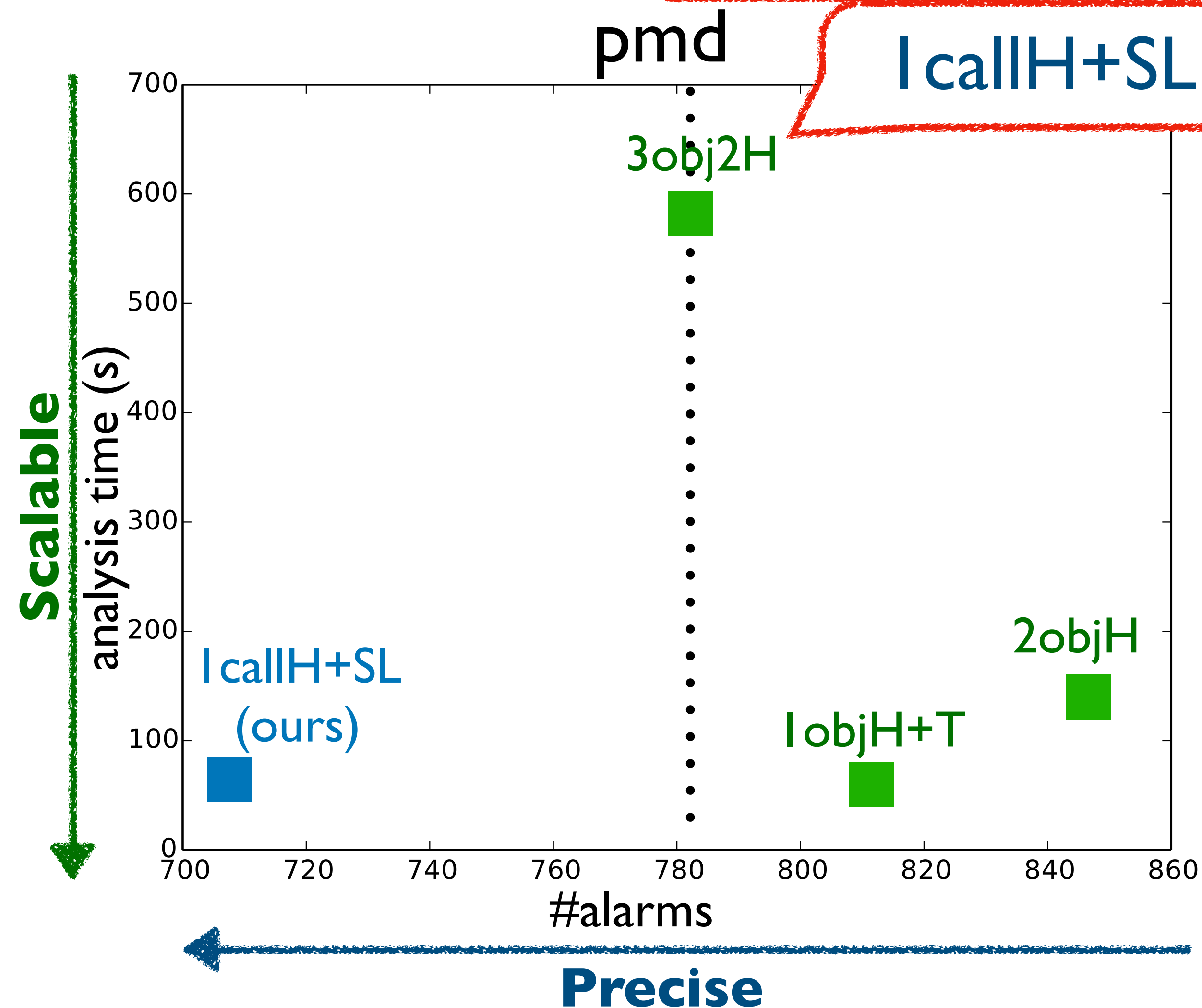
Precise

Precise

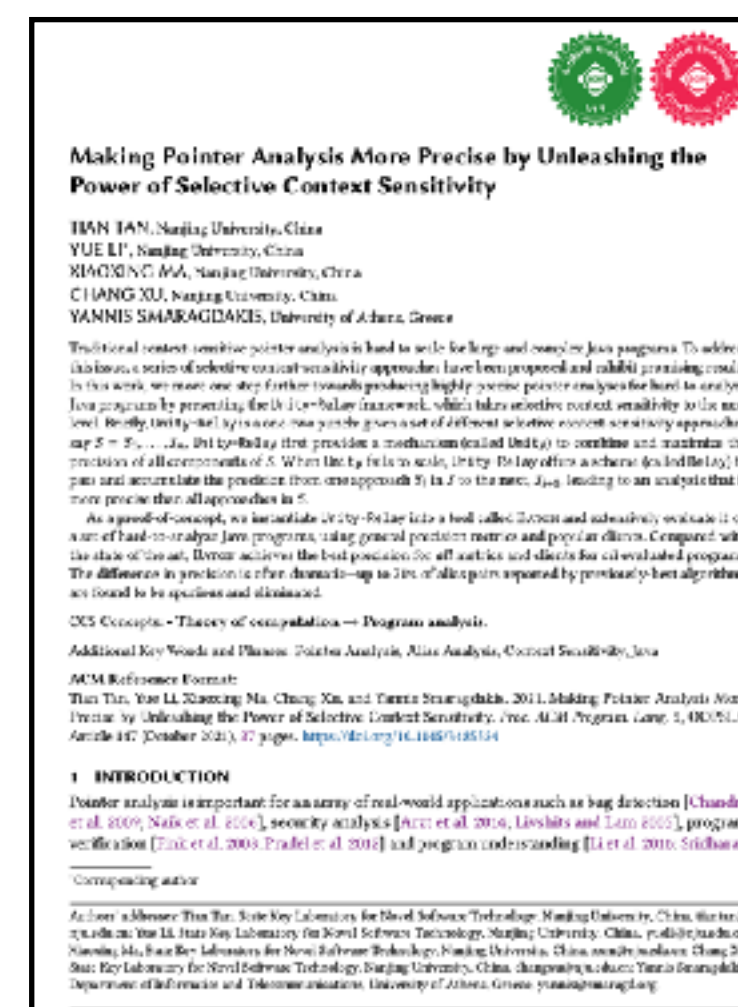
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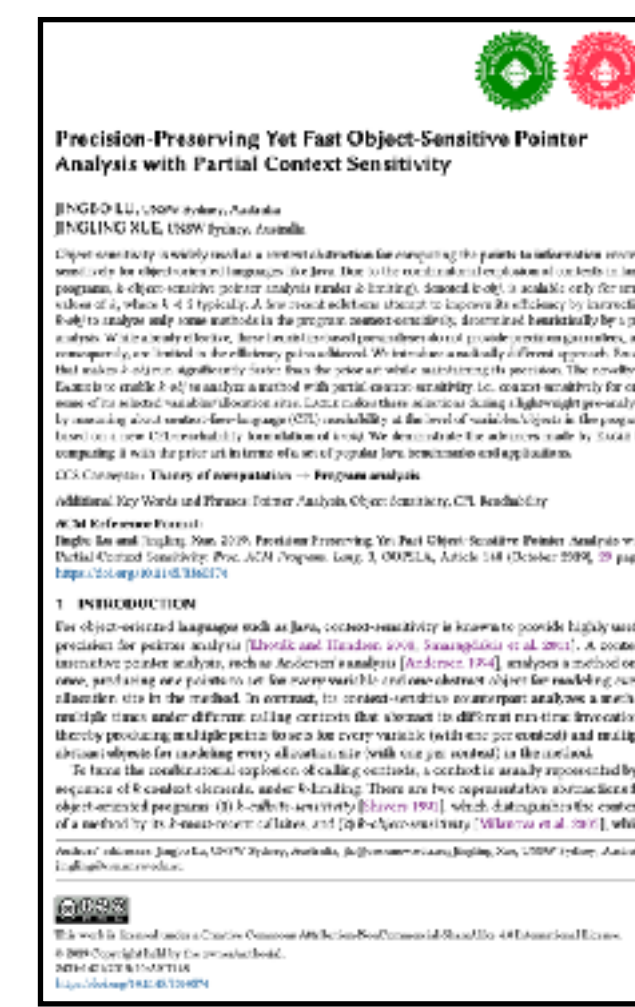
IcallH+SL is **even more precise** than 3obj2H



Precision upper bound of recent researches on **object sensitivity**



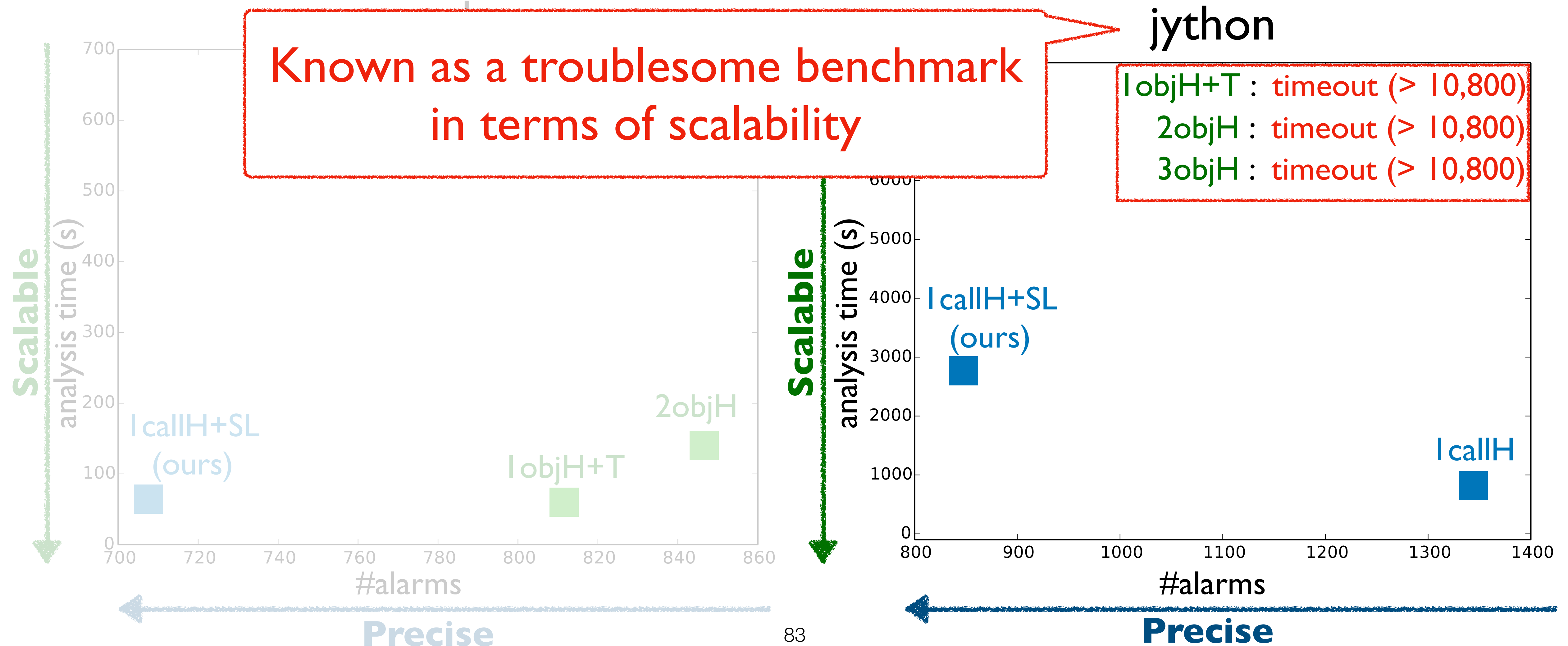
OOPSLA 2021



OOPSLA 2019

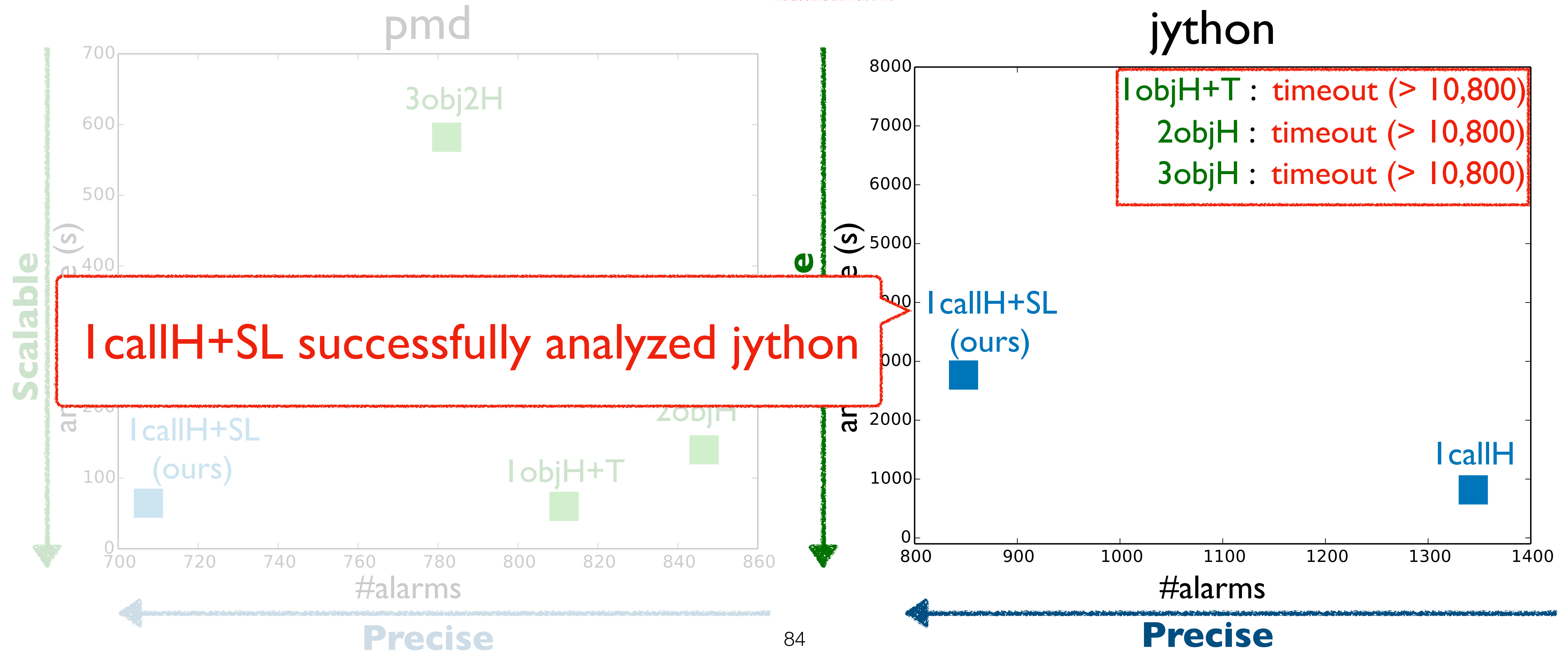
Call-site Sensitivity vs Object Sensitivity

- `lcallH+SL` (ours) is more precise and **scalable** than the existing object sensitivities



Call-site Sensitivity vs Object Sensitivity

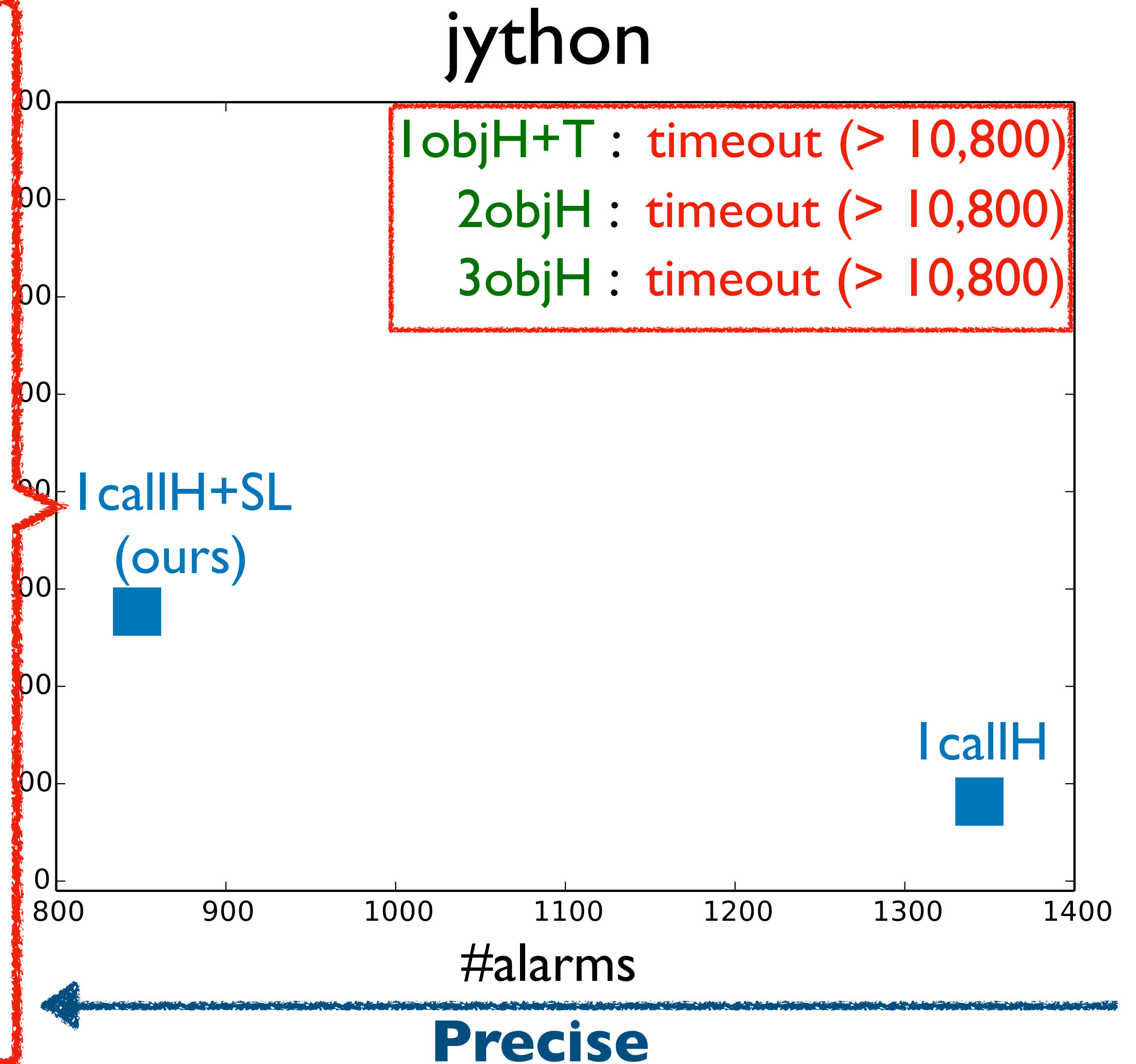
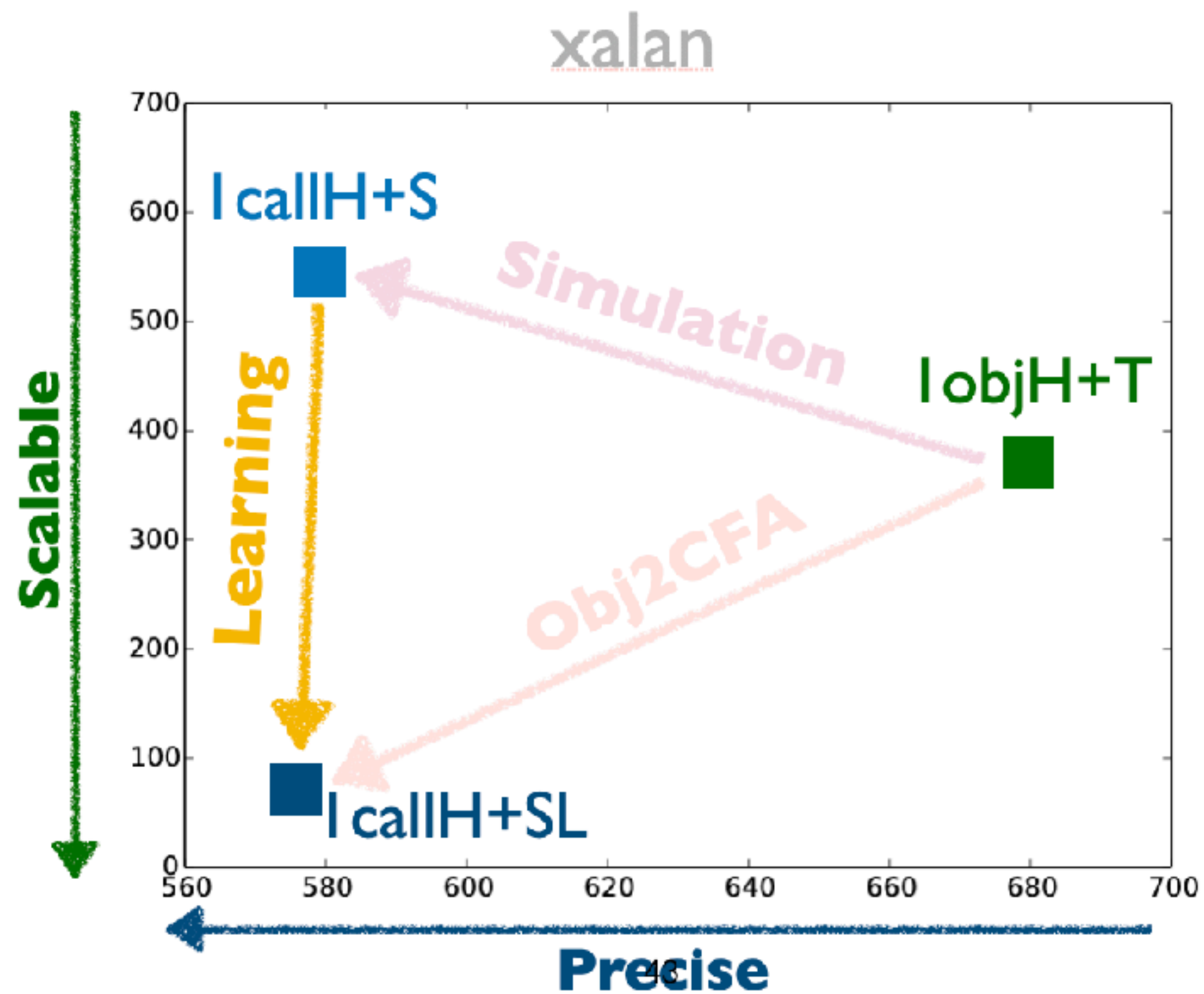
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Call-site Sensitivity vs Object Sensitivity

- `IcallH+SL` (ours) is more precise and **scalable** than the existing object sensitivities

- Necessity of learning
- `IcallH+S` is unable to analyze `xython`



Summary

- Currently, CFA is known as a bad context
- However, if context tunneling is included, CFA is not a bad context anymore
- We need to reconsider CFA from now on

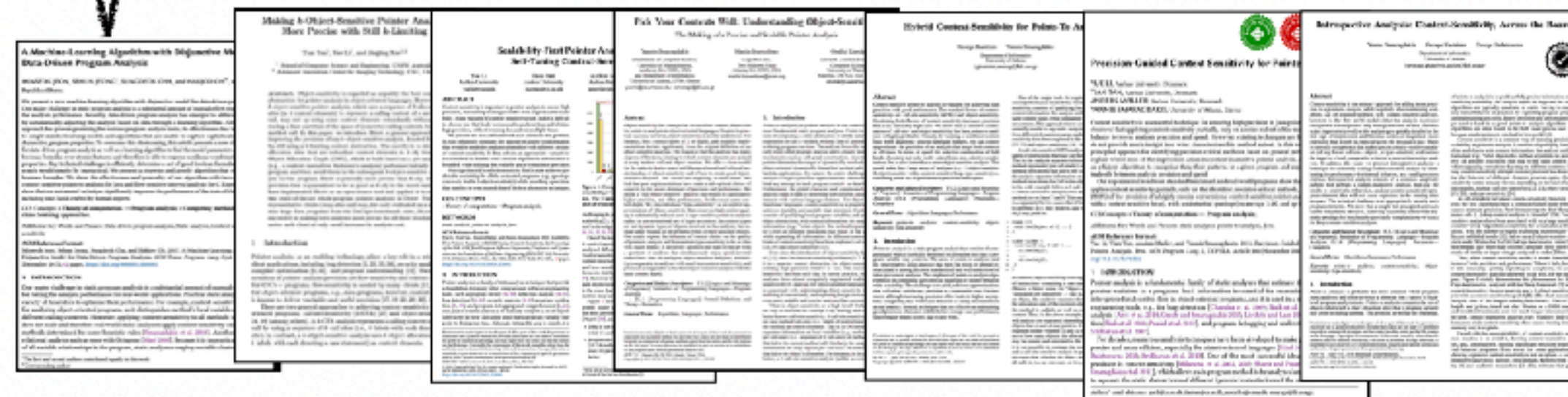
Thank you

Summary

- Currently, CFA is known as a bad context

- Call-site Sensitivity has been ignored

“... call-site-sensitivity is less important than others ...”
- Jeon et al. [2019]



1981

2002

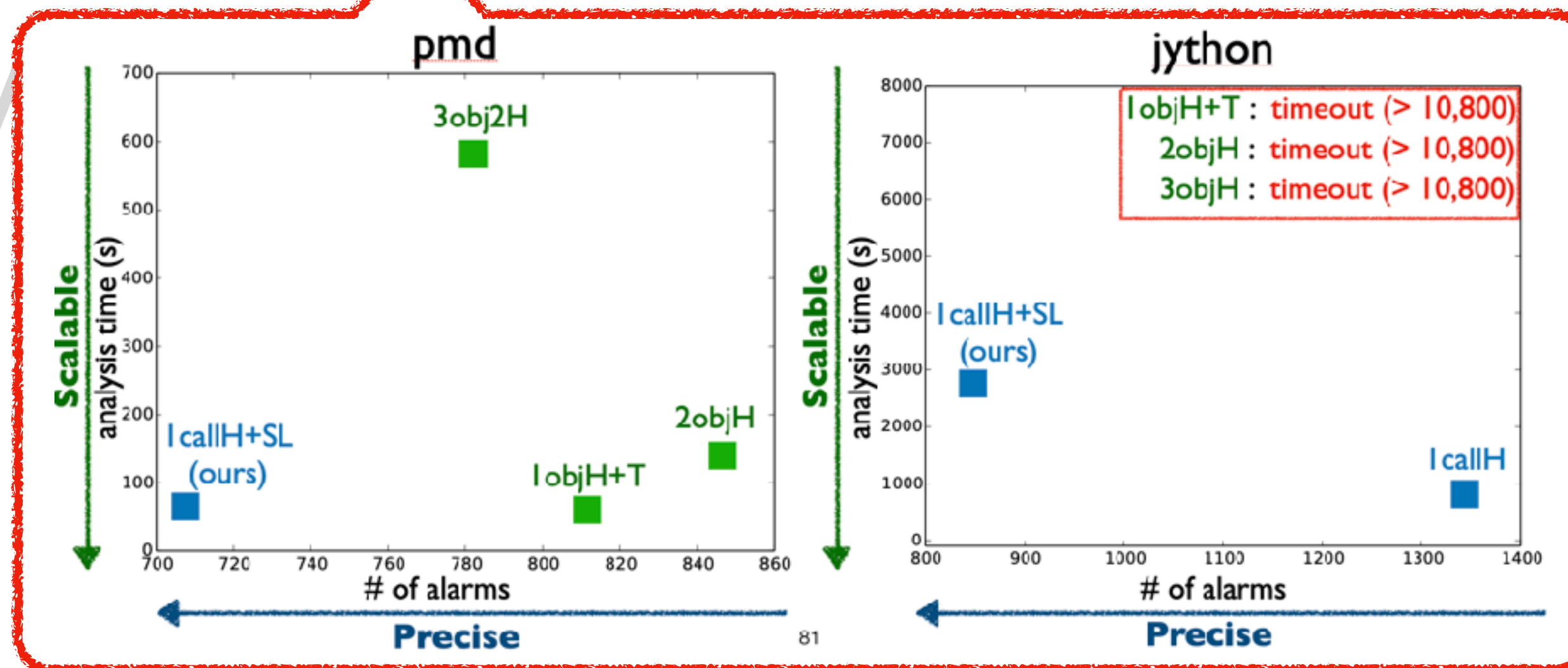
2010

2022

Summary

- Currently, CFA is known as a bad context

- However, if context tunneling is included, CFA is not a bad context anymore



- W

om now on

Return of CFA: Call-Site Sensitivity Can Be Superior to Object Sensitivity Even for Object-Oriented Programs

MINSEOK JEON and HAKJOO OH*, Korea University, Republic of Korea

In this paper, we challenge the commonly accepted wisdom in static analysis that object sensitivity is superior to call-site sensitivity for object-oriented programs. In static analysis of object-oriented programs, object sensitivity has been established as the dominant flavor of context sensitivity thanks to its outstanding precision. On the other hand, call-site sensitivity has been regarded as unsuitable and its use in practice has been constantly discouraged for object-oriented programs. In this paper, however, we claim that call-site sensitivity is generally a superior context abstraction because it is practically possible to transform object sensitivity into more precise call-site sensitivity. Our key insight is that the previously known superiority of object sensitivity holds only in the traditional k -limited setting, where the analysis is enforced to keep the most recent k context elements. However, it no longer holds in a recently-proposed, more general setting with context tunneling. With context tunneling, where the analysis is free to choose an arbitrary k -length subsequence of context strings, we show that call-site sensitivity can simulate object sensitivity almost completely, but not vice versa. To support the claim, we present a technique, called Obj2CFA, for transforming arbitrary context-tunneled object sensitivity into more precise, context-tunneled call-site sensitivity. We implemented Obj2CFA in Deep and used it to derive a new call-site-sensitive analysis from a state-of-the-art object-sensitive pointer analysis. Experimental results confirm that the resulting call-site sensitivity outperforms object sensitivity in precision and scalability for real-world Java programs. Remarkably, our results show that even 1-call-site sensitivity can be more precise than the conventional 3-object-sensitive analysis.

1 INTRODUCTION

"Since its introduction, object sensitivity has emerged as the dominant flavor of context sensitivity for object-oriented languages."

—Smaragdakis and Balatsouras [2015]

Context sensitivity is critically important for static program analysis of object-oriented programs. A context-sensitive analysis associates local variables and heap objects with context information of method calls, computing analysis results separately for different contexts. This way, context sensitivity prevents analysis information from being merged along different call chains. For object-

CFA wins!

uses the allocation-site of the receiver object (a) as the context of τ . The standard k -object-sensitive analysis [Milanov et al. 2002, 2005; Smaragdakis et al. 2011] maintains a sequence of



Obj



CFA

I return!

- We need to **reconsider** CFA from now on

Thank you